### QoS support in a switch

Analysis of Policed Traffic
Through a Switch with Shared Buffer Space

by

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## Background

- ISPs sign SLAs with customers
- Typically an SLA contains:
  - Priority or QoS class
  - Committed rate: rate the ISP guarantee to the customer
  - Promise for 'best-effort' for excess rate.
- The ISP police and mark the incoming traffic at the network edge, e.g., using leaky bucket.

# Background(2)

- Committed rates are used for routing, traffic engineering, etc.
- Excess traffic brings additional revenue
- To maximize profit, ISPs may use statistical multiplexing to overbook committed traffic.
- In case of congestion, committed traffic have precedence over excess traffic even with higher QoS class.

## **Implications**

- Every switching element must be able to differentiate between QoS classes and the commit bit.
- We assume no PFQ, hardware must be kept economically operational. Threholds are the way to go.

### Input port model

- VOQs
- Queue per VOQ and per class.
- Classes share buffer space. protection vs. utilization
- Committed and excess traffic must be kept in the same Q (or OOO)

Typically 4-8 QoS classes.

1 or 2 has strict priority; rest are served with WRR.

## Analysis Model

#### **Priorities**

- Priority queue 2 will simulate all the higher priority traffic in the system and will have strict priority over the other queues.
- Priority queue 1 has two traffic types, committed and excess.
- Priority queue 0 has two traffic types, committed and excess.

#### **Thresholds**

Thresholds limit the portion of buffer space traffic type can capture.

- Priority queue 2 rate decreases the service rate.
- Priority queue 1 committed and excess traffic have thresholds,  $\alpha_{1C}$  and  $\alpha_{1E}$ .
- Priority queue 0 committed and excess traffic have thresholds,  $\alpha_{0C}$  and  $\alpha_{0E}$

### The analyzed system

Two queues that share a buffer space of n packets (or cells).

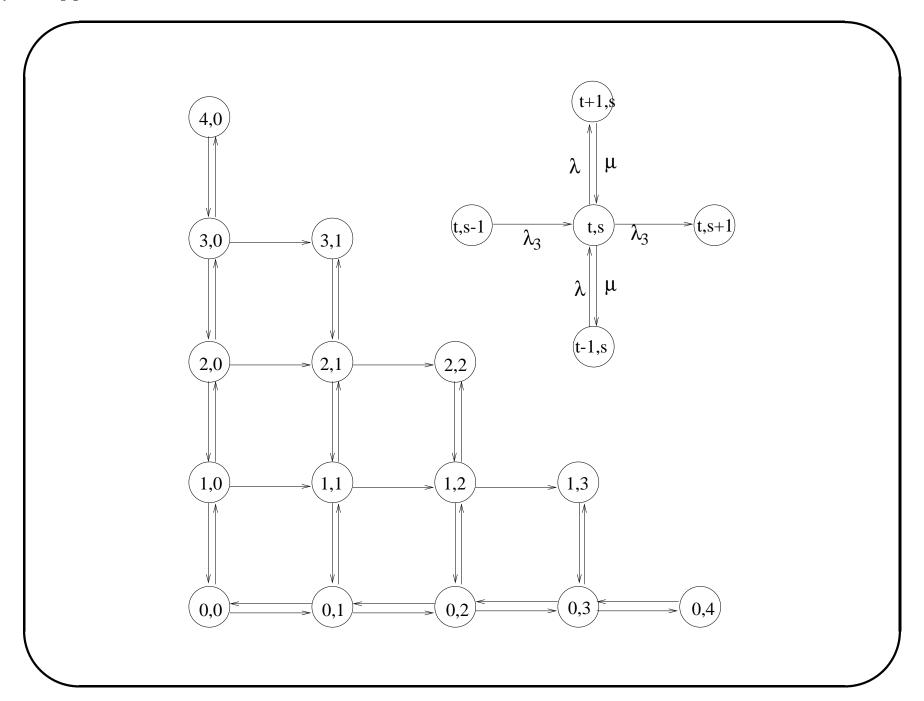
The high priority queue has committed traffic and excess traffic with rates  $\lambda_1$  and  $\lambda_2$ . The low priority queue has committed traffic with rate  $\lambda_3$ . Service rate is  $\mu$ .

Threshold is  $n_{th} = \alpha_{th} n$ , above threshold occupancy of the buffer excess high priority traffic is not accepted.

Strict priority.

The above system can be modeled by a continuous-time Markov chain with (n+1)(n+2)/2 states.

Each state is represented by the ordered pair (t, s), where t is the number of high priority packets in the buffer and s the number of low priority packets.



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The infinitesimal transition rates from state (t, s) to state (t', s'),  $q_{t,s,t',s'}$  are:

$$q_{t,s,t-1,s} = \mu$$

$$q_{0,s,0,s-1} = \mu$$

$$q_{t,s,t,s+1} = \lambda_{3}$$

$$q_{t,s,t+1,s} = \begin{cases} \lambda_{1} + \lambda_{2} & \text{if } t+s \leq n_{th} \\ \lambda_{1} & \text{if } t+s > n_{th} \end{cases}$$

$$-q_{t,s,t,s} = \begin{cases} \lambda_{1} + \lambda_{2} + \lambda_{3} & \text{if } t+s = 0 \\ \lambda_{1} + \lambda_{2} + \lambda_{3} + \mu & \text{if } 0 < t+s \leq n_{th} \\ \lambda_{1} + \lambda_{3} + \mu & \text{if } t+s > n_{th} \end{cases}$$

Note that  $-q_{t,s,t,s}$  is the transition rate out of state (t,s)

We wish to find the steady state probabilities,  $\pi_{t,s}$ .

$$\vec{\pi}Q = 0$$

$$\vec{\pi}Q = 0$$

$$\sum_{(t,s)} \pi_{t,s} = 1$$

This numerical solution requires  $O(n^{2(2+\alpha)}) \simeq O(n^5)$ .

In the following, we shall describe methods to make the problem more tractable, by presenting a recursive solution that requires only  $O(n^3)$  operations.

#### Calculating Drop probabilities

Let  $\lambda_1 + \lambda_3 < \mu$  or else our solution is meaningless.

$$\eta_1 = \sum_{i=0}^n \pi_{i,n-i} \tag{2}$$

$$\eta_2 = \sum_{i+j>n_{th}} \pi_{i,j} \tag{3}$$

$$\eta_3 = \sum_{i=0}^n \pi_{i,n-i} \tag{4}$$

By definition  $\eta_1 = \eta_3$  which shows that there is no preference between the two priority classes in the acceptance probability.

#### Average delay for the lower class:

Let  $\bar{N}_i$  be the average number of cells of type i in the system.

$$ar{N_3} = \sum_{i,j} \pi_{i,j} j$$

Using Little's Law we know that the average delay,  $T_3$ , is given by

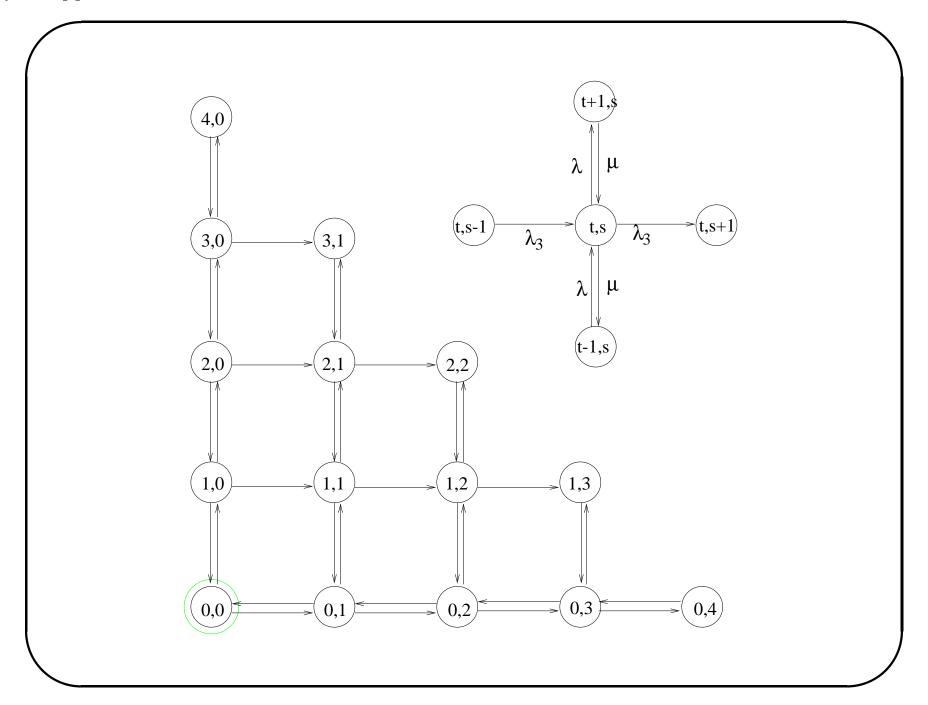
$$T_3 = \frac{\bar{N}_3}{(1 - \eta_3)\lambda_3} = \frac{\sum_{i,j} \pi_{i,j} j}{(1 - \sum_{i=0}^n \pi_{i,n-i})\lambda_3}$$
 (5)

### Reducing the Analysis Complexity Using Recurrence

#### Grand plan

Write the steady state probabilities of all the system states,  $\pi_{t,s}$ , as functions of  $\pi_{0,s}$ ,  $0 \le s \le n$ .

Then, we can write n equilibrium equations and together with the probability conservation equation we obtain n+1 linear equations that can be solved with complexity of  $O(n^3)$ .



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we can write the following n(n+1)/2 equilibrium equations

$$-(\lambda_{1} + \lambda_{2} + \lambda_{3} + \mu)\pi_{t,s} = \lambda_{3}\pi_{t,s-1} + \mu\pi_{t+1,s} + (\lambda_{1} + \lambda_{2})\pi_{t-1,s}, \quad t > 0, t + s \leq n, \quad (7)$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{t,s} = \lambda_{3}\pi_{t,s-1} + \mu\pi_{t+1,s} + (\lambda_{1} + \lambda_{2})\pi_{t-1,s}, \quad t > 0, t + s = n, \quad + 1$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{t,s} = \lambda_{3}\pi_{t,s-1} + \mu\pi_{t+1,s} + \lambda_{1}\pi_{t-1,s}, \quad t > 0, t + s > n_{th} + 1$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{t,0} = \mu\pi_{t+1,0} + (\lambda_{1} + \lambda_{2})\pi_{t-1,0}, \quad 1 < t \leq n_{th}$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{n_{th}+1,0} = \mu\pi_{n_{th}+2,0} + (\lambda_{1} + \lambda_{2})\pi_{n_{th},0}$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{t,0} = \mu\pi_{t+1,0} + \lambda_{1}\pi_{t-1,0}, \quad n_{th} + 1 < t \leq n - 1$$

$$-(\lambda_{1} + \lambda_{2} + \lambda_{3} + \mu)\pi_{0,s} = \lambda_{3}\pi_{0,s-1} + \mu\pi_{1,s} + \mu\pi_{0,s+1}, \quad s \leq n_{th}$$

$$-(\lambda_{1} + \lambda_{3} + \mu)\pi_{0,s} = \lambda_{3}\pi_{0,s-1} + \mu\pi_{1,s} + \mu\pi_{0,s+1}, \quad s > n_{th}$$

$$-(\lambda_{1} + \lambda_{2} + \lambda_{3})\pi_{0,0} = \mu\pi_{1,0} + \mu\pi_{0,1}$$

Now, we can write the following recursion relations for  $\pi_{t,s}$ , t>0:

The above recurrence suggests that all  $\pi_{t,s}$  can be written as functions of  $\pi_{0,s}$ , i.e.,

$$\pi_{t,s} = \sum_{l=0}^{n} C_{t,s}(l) \pi_{0,l}, \tag{9}$$

It is easier to calculate the recurrence for the coefficients,  $C_{t,s}(l)$ , rather than directly for  $\pi_{t,s}$ .

First, we calculate the coefficients of  $\pi_{1,s}$  by

$$C_{1,s}(s) = -q_{0,s,0,s}/q_{1,s,0,s} \quad s = 0, 1, 2, \dots, n-1$$

$$C_{1,s}(s-1) = -q_{0,s-1,0,s}/q_{1,s,0,s} \quad s = 1, 2, \dots, n-1$$

$$C_{1,s}(s+1) = -q_{0,s+1,0,s}/q_{1,s,0,s} \quad s = 0, 1, 2, \dots, n-1$$

$$C_{1,s}(l) = 0 \quad |l-s| > 1$$

Next, we calculate the coefficients of  $\pi_{t,s}$  for  $t=2,3,\ldots,n-1$ :

$$C_{t,s}(m) = (q_{t-1,s,t-1,s}C_{t-1,s}(m) - q_{t-1,s-1,t-1,s}C_{t-1,s-1}(m) - q_{t-2,s,t-1,s}C_{t-2,s}(m))/q_{t,s,t-1,s}$$
(11)

The recurrence calculation requires  $O(n^3)$  operations.

The following n + 1 linear equation system, is made from n (out of n + 1) unused equilibrium equations plus the probability conservation equation:

$$-q_{t,n-t,t,n-t}\pi_{t,n-t} = q_{t,n-(t+1),t,n-t}\pi_{t,n-(t+1)} + (12)$$

$$q_{t-1,n-t,t,n-t}\pi_{t-1,n-t}$$

$$1 \le t \le n-1$$

$$-q_{0,n,0,n}\pi_{0,n} = q_{0,n-1,0,n}\pi_{0,n-1}$$

$$\sum_{(t,s)} \pi_{t,s} = 1$$

Solution complexity is lower than  $O(n^3)$ .

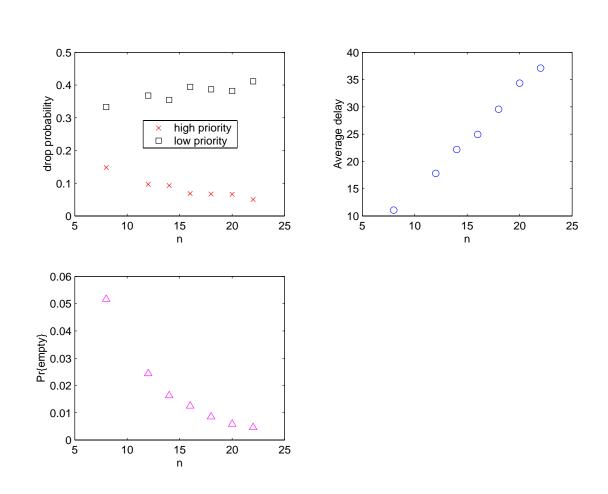
Using the recurrence on the coefficients we can write Eq. 12 as

$$-\sum_{m=0}^{n} q_{t,n-t,t,n-t} C_{t,n-t}(m) \pi_{0,m} = \sum_{m=0}^{n} \left( q_{t,n-(t+1),t,n-t} C_{t,n-(t+1)}(m) + q_{t-1,n-t,t,n-t} C_{t-1,n-t}(m) \right) \pi_{0,m}$$

$$1 \le t \le n-1$$

and rewrite the probability conservation equation as

$$\sum_{t=0}^{n-1} \sum_{s=0}^{n-t} \sum_{m=0}^{n} C_{t,s}(m) \pi_{0,m} = 1$$



The analysis as a function of the buffer size, n.

$$\lambda_1 = \lambda_2 = \lambda_3 = 0.4$$
 and  $\alpha_{th} = 0.8$ 

### **Simulation Results**

Line rate is 1Gbps

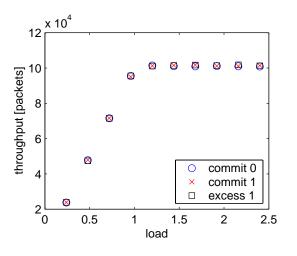
The buffer size is set to 100 packets

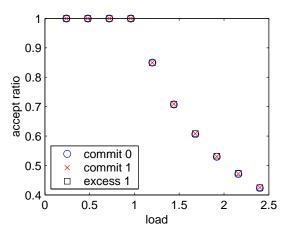
The packet arrival process is Poisson with combined rate 1.2

$$(\lambda_1 = \lambda_2 = \lambda_3 = 0.4).$$

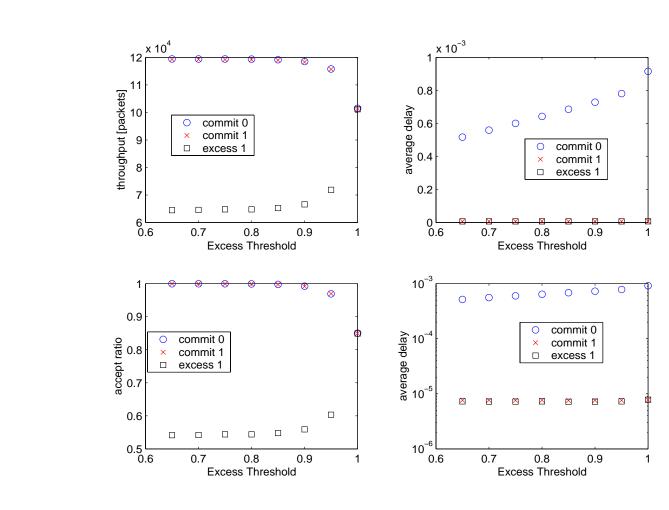
The packet length is Exponentially distributed, but bounded to be between 40 and 1500 bytes.

#### No thresholds

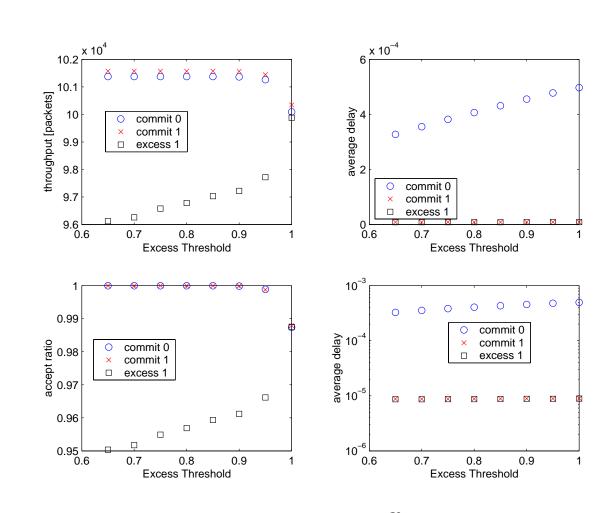




No admission control to the buffer



Thrsehhold admission control to the buffer, where  $\lambda_1 = \lambda_2 = \lambda_3 = 0.4$ .



Thrsehhold admission control to the buffer, where  $\lambda_1 = \lambda_2 = \lambda_3 = 0.34$ .

Low threshold for excess class 1 traffic  $\Rightarrow$  good protection and lower delay for committed class 0 traffic.

High threshold for excess class 1 traffic  $\Rightarrow$  good capability to pass bursts of excess traffic.

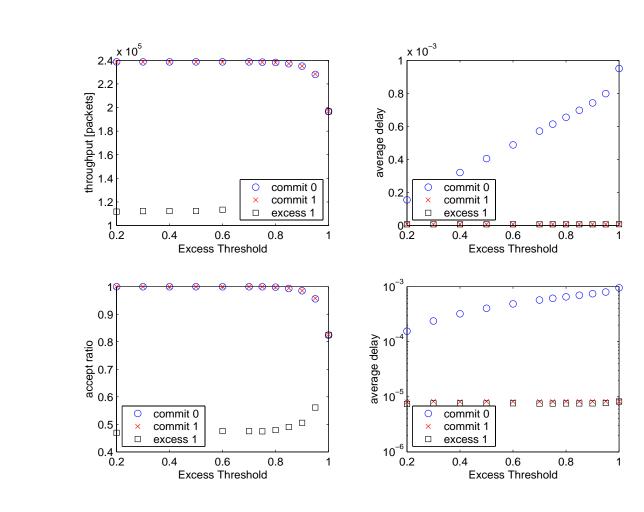
We want to have both!

#### Cross traffic threshold

We ALSO drop excess high priority traffic when the queue of class 0 committed traffic crosses a threshold.

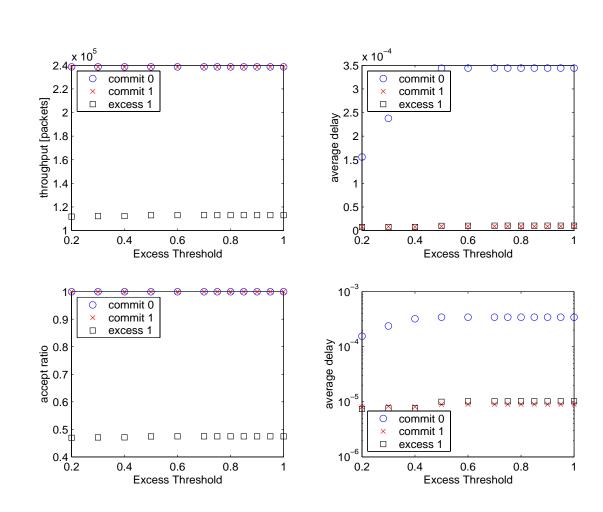
This enables us to set the previous threshold quite high.

We can use the similar analysis to drive numerical results.



Thrsehhold admission control to the buffer, where

$$\lambda_1 = \lambda_2 = \lambda_3 = 0.4.$$



Cross-threshold admission control to the buffer,  $\alpha_{1E0} = 0.4$ , where  $\lambda_1 = \lambda_2 = \lambda_3 = 0.4$ .

### **Conclusions**

- Initial results show potential
- We need to study the system numerically and by simulations to get better understanding of the capabilities.
- Need to simulate bursty arrivals.