



Please submit your handwritten answers in pairs, using the box in front of IZ338 before the exercise timeslot on the due date above. Make sure to include your full names, matriculation numbers, and the programmes that you are enrolled in. In accordance with the *guidelines* of the TU Braunschweig, using AI tools to solve any part of the exercises is **not permitted**.

Exercise 1 (Safehouse Problem).

(5+15 points)

You've successfully pulled off your biggest heist yet – you've stolen the golden Leibniz cookie! However, agents of *baked goods manufacturing co.* are now on your tail, and all roads out of the city have been locked down in search for the famous symbol of delicacy: You cannot leave. You decide that your best bet is to find a safehouse in town as far as possible from all cookie outlets, hide the cookie there, and lie low.

Given an axis-aligned rectangle R in the Euclidean plane \mathbb{R}^2 that defines the city limits and the locations of nearby (both inside and outside the city) cookie outlets c_1, c_2, \dots, c_n , you need a location inside R that maximizes the distance to the closest cookie outlet, see Fig. 1. (Hint: Start by thinking about a geometric suitable structure from the lecture to discretize the problem.)

- a) Identify a (*finite!*) set of candidate locations in R to choose from. Argue why an optimal solution is contained in this set.
- b) Design an $\mathcal{O}(n \log n)$ time algorithm based on your candidates and prove its correctness.

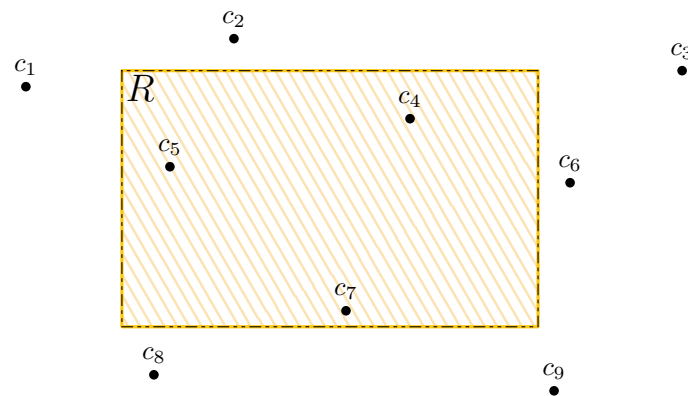


Figure 1: You require a location inside R that maximizes the distance to the nearest cookie outlet c_i .

Exercise 2 (Minimal enclosing disk).

(5+15 points)

Let \mathcal{P} be a set of n points in the Euclidean plane in *general position*, such that no four points in \mathcal{P} are concyclic: No four points lie on a common circle. A *minimal enclosing disk* $md(\mathcal{P})$ is a disk with minimal radius that contains \mathcal{P} . Let $c \in \mathbb{R}^2$ and $r \in \mathbb{R}$ be its center and radius.

Using concepts from the lecture and the tutorial, design an algorithm that determines $md(\mathcal{P})$ in less than $\mathcal{O}(n^2)$ time, or better. (Hint: Review Tutorial 4)