LINEAR PROGRAMMING

[V. CH3]: DEGENERACY

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November 21, 2023

DEFINITION/EXAMPLES OF DEGENERACY

PERTURBATION METHOD

FUNDAMENTAL THEOREM

GEOMETRY

OTHER PIVOT RULES

We say that a dictionary is *degenerate* if $\bar{b}_i = 0$ for some $i \in \mathcal{B}$.

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It is clear that the problem is unbounded and therefore no more pivots were required.

→ A degenerate dictionary could cause difficulties for the simplex method. Problems arise, when a degenerate dictionary produces *degenerate pivots*.

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$$\frac{\zeta = 3 - 0.5x_1 + 2x_2 - 1.5w_1}{x_3 = 1 - 0.5x_1 - 0.5w_1}$$

$$w_2 = 0 + x_1 - 1.0x_2 + w_1$$

For this dictionary, the entering variable is x_2 and the ratio computed to determine the leaving variable is

$$\{ i \in \mathcal{B} : \frac{\bar{b}_i}{\bar{a}_{i2}} \text{ with } \bar{a}_{i2} > 0 \} = \{ \frac{0}{1.0} \}$$

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 \leadsto The fact that the ratio is zero means that as soon as x_2 is increased from zero to a positive value, w_2 will go negative.

$\zeta =$	3 +	$1.5 x_1 -$	$2w_2 +$	$0.5 \ w_1$
$x_3 =$	1 -	$0.5x_1$	_	$0.5w_1$
$x_2 =$		x_1 -	$w_2 +$	w_1

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But we are just representing the solution in a new way.

The entering variable for the next iteration is x_1 and the leaving variable is x_3 , producing a *nondegenerate* pivot that leads to

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$$\frac{\zeta = 6 - 3x_3 - 2w_2 - w_1}{x_1 = 2 - 2x_3 - w_2}$$

$$x_2 = 2 - 2x_3 - w_2$$

While it is typical for some pivot to "break away" from the degeneracy, the real danger is that

the simplex method will make a sequence of degenerate pivots and eventually return to a dictionary that has appeared before, in which case the simplex method enters an infinite loop and never finds an optimal solution.

This behavior is called *cycling*.

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This behavior is called *cycling*.

→ Cycling is rare! A program that generates random 2×4 fully degenerate problems was run more than one billion times and did not find one example!

Unfortunately, under certain pivoting rules, cycling is possible.

In fact, cycling is possible even when using one of the most popular pivoting rules:

- Entering variable: Choose the entering variable as the one with the largest coefficient in the ζ-row of the dictionary.
- Leaving variable: When two or more variables compete for leaving the basis, reading left to right, pick the first leaving-variable candidate from the list:

$$x_1, x_2, \cdots, x_n, w_1, w_2, \cdots, w_m$$

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$$x_1, x_2, \cdots, x_n, w_1, w_2, \cdots, w_m$$

Here is an example that cycles:

$\zeta =$	+	x_1 -	$2x_2$	_	$2x_4$
$\overline{w_1} =$	_	$0.5x_1 +$	$3.5x_2 +$	$2x_3 -$	$4x_4$
$w_2 =$	_	$0.5x_1 +$	$x_2 + $	$0.5x_3 -$	$0.5x_4$
$w_3 =$	1 -	x_1			

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$\zeta =$		_	$2w_2 + $	$1 x_3 -$	$3x_4$
$x_1 =$	+	$0.8w_1 -$	$2.8w_2 -$	$0.2x_3 +$	$1.8x_4$
$x_2 =$	+	$0.4w_1 -$	$0.4w_2 -$	$0.6x_3 +$	$1.4x_4$
$w_3 =$	1 -	$0.8w_1 +$	$2.8w_2 +$	$0.2x_3 -$	$1.8x_4$

For the third iteration, x_3 enters and x_1 leaves:

$\zeta =$	+	$4 w_1 -$	$16w_2 -$	$5x_1 + $	$6 x_4$
$x_3 =$	+	$4w_1 -$	$14w_2 -$	$5x_1 + $	$9x_4$
$x_2 =$	_	$2w_1 + $	$8w_2 + $	$3x_1$ -	$4x_4$
$w_3 =$	1		_	x_1	

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Lastly, for the sixth iteration, w_2 enters and x_4 leaves:

$\zeta =$		_	$2x_4 + $	$1x_1 - $	$2x_2$
$\overline{w_1} =$	+	$2x_3 -$	$4x_4 -$	$0.5x_1 +$	$3.5x_{2}$
$w_2 =$	+	$0.5x_3 -$	$0.5x_4 -$	$0.5x_1 +$	x_2
$w_3 =$	1		_	x_1	

Lastly, for the sixth iteration, w_2 enters and x_4 leaves:

$\zeta =$		_	$2x_4 + $	$1x_1 -$	$2x_2$
$\overline{w_1} =$	+	$2x_3 -$	$4x_4 -$	$0.5x_1 +$	$3.5x_2$
$w_2 =$	+	$0.5x_3 -$	$0.5x_4$ -	$0.5x_1 +$	x_2
$w_3 =$	1		_	x_1	

Note that we have come back to the original dictionary:

→ from here on the simplex method simply cycles through these six dictionaries and never makes any further progress toward an optimal solution.

As bad as cycling is, the following theorem tells us that nothing worse can happen:

THEOREM

If the simplex method fails to terminate, then it must cycle.

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PROOF.

A dictionary is completely determined by specifying which variables are basic and which are nonbasic. There are only

$$\left(\begin{array}{c} n+m \\ m \end{array}\right) = \frac{(n+m)!}{n!m!}$$

different possibilities. This number is big, but it is finite. If the simplex method fails to terminate, it must visit some of these dictionaries more than once. Hence, the algorithm cycles.

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Whenever a vanishing "rhs", (zero \bar{b}_i), appears *perturb* it.

- set :
$$ar{b}_i = ar{b}_i + \epsilon_i$$

If there are lots of them, say k, perturb them all. Make the perturbations at different *scales*:

$$0<\epsilon_k\ll\cdots\ll\epsilon_2\ll\epsilon_1\ll$$
 other nonzero data

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 other nonzero data

Try to keep in mind that:

- no linear combination of the ϵ_i 's using coefficients that might arise in the course of the simplex method can ever produce a number whose size is of the same order as the data in the problem.
- each of the "lower down" ϵ_i 's can never "escalate" to a higher level.

Hence, cancellations can only occur on a given scale.

First example:

EXAMPLE

$\zeta =$	0	+	$6x_1 +$	$4x_2$
$w_1 =$	0	+	$9x_1 + $	$4x_2$
$w_2 =$	0	_	$4x_1 -$	$2x_2$
$w_3 =$	1		_	x_2

The first step is to introduce symbolic parameters

$$0 < \epsilon_3 \ll \epsilon_2 \ll \epsilon_1$$

to get a perturbed problem:

This dictionary is *not degenerate*. The entering variable is x_1 and the leaving variable is unambiguously w_2 .

The next dictionary is

$\zeta =$	0	+	$1.5\epsilon_2$	_	$1.5w_2 +$	$1 x_2$
$\overline{w_1} =$	0+	ϵ_1 +	$2.25\epsilon_2$	_	$2.25w_2$ -	$0.5x_2$
$x_1 =$	0	+	$0.25\epsilon_2$	_	$0.25w_2 -$	$0.5x_2$
$w_3 =$	1		+	ϵ_3	_	x_2

The next dictionary is

For the next pivot, the entering variable is x_2 and, using the fact that $\epsilon_2 \ll \epsilon_1$, we see that the leaving variable is x_1 . The new dictionary is

$\zeta =$	0	+	$2\epsilon_2$	_	$2w_2$ –	$2x_1$
$\overline{w_1} =$	0+	ϵ_1 +	$2\epsilon_2$	_	$2w_2 +$	$\overline{x_1}$
$x_2 =$	0	+	$0.5\epsilon_2$	_	$0.5w_2 -$	$2x_1$
$w_3 =$	1	_	$0.5\epsilon_{2} +$	ϵ_3 +	$0.5w_2 +$	$2x_1$

This last dictionary is optimal. At this point, we simply drop the symbolic ϵ_i parameters and get an optimal dictionary for the unperturbed problem.

$\zeta =$	0 —	$2w_2$ -	$2x_1$
$w_1 =$	0 —	$2w_2 + $	x_1
$x_2 =$	0 -	$0.5w_2 -$	$2x_1$
$w_3 =$	1 +	$0.5w_2 +$	$2x_1$

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When treating the ϵ_i 's as symbols, the method is called the *lexicographic* method.

- → The lexicographic method does not affect the choice of entering variable but does amount to a precise prescription for the choice of leaving variable.
- → The lexicographic method produces a variant of the simplex method that never cycles:

THEOREM

The simplex method always terminates provided that the leaving variable is selected by the lexicographic rule.



It suffices to show that no degenerate dictionary is ever produced.

As we have discussed before, the ϵ_i 's operate on different scales and hence cannot cancel with each other. Therefore, can think of the ϵ_i 's as a collection of independent variables.

PROOF.

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Extracting the ϵ terms from the first dictionary, we see that we start with the following pattern:





PROOF CONT.

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This system of linear combinations is obtained from the original system by pivot operations and, since pivot operations are reversible, it follows that the rank of the two systems must be the same.

PROOF CONT.

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This system of linear combinations is obtained from the original system by pivot operations and, since pivot operations are reversible, it follows that the rank of the two systems must be the same.

Since the original system had rank m, we see that every subsequent system must have rank m.

- \rightarrow There must be at least one nonzero r_{ij} in every row i, which of course implies that
- \rightarrow None of the rows can be degenerate.

Hence, no dictionary can be degenerate.

Second example:

EXAMPLE

$\zeta =$	0	+	$2x_1 + $	$4x_2$
$w_1 =$	0	+	x_1 -	x_2
$w_2 =$	0	+	$3x_1 -$	x_2
$w_3 =$	0	_	$4x_1 + $	x_2

Perturb vanishing rhs

$\zeta =$	0.0			+	$2x_1 + $	$4x_2$
$w_1 =$	0.0 +	ϵ_1		+	x_1 -	x_2
$w_2 =$	0.0	+	ϵ_2	+	$3x_1 - $	x_2
$w_3 =$	0.0		+	ϵ_3 –	$4x_1 + $	x_2

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or equivalently

 x_2 enters and as $\epsilon_2 \ll \epsilon_1, w_2$ leaves.

Perturb vanishing rhs

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 x_2 enters and as $\epsilon_2 \ll \epsilon_1, w_2$ leaves.

 \rightsquigarrow Note that, as ϵ_i 's have different scale, the will not cancel each other.

We get

$\zeta =$	0.0 +	$0.0\epsilon_1$ +	$4.0\epsilon_2$ +	$0.0\epsilon_3$ +	$14 x_1 -$	$4w_2$
$w_1 =$	0.0 +	$1.0\epsilon_1$ –	$1.0\epsilon_{2} +$	$0.0\epsilon_3$ –	$2x_1 + $	w_2
$x_2 =$	0.0 +	$0.0\epsilon_1$ +	$1.0\epsilon_2$ +	$0.0\epsilon_{3} +$	$3x_1$ -	w_2
$w_3 =$	0.0 +	$0.0\epsilon_{1} +$	$1.0\epsilon_{2} +$	$1.0\epsilon_3$ –	x_1 -	w_2

We get

 x_1 enters and $\ \ {
m as}\ \epsilon_2+\epsilon_3\ll\epsilon_1-\epsilon_2, w_3 \ \ {
m leaves.} \ \ {
m We get}$

$\zeta =$	0.0 +	$0.0\epsilon_1$ +	$18\epsilon_2$ +	$14\epsilon_3$ –	$14w_3 -$	$18w_2$
$\overline{w_1} =$	0.0 +	$1.0\epsilon_1$ –	$3.0\epsilon_2$ –	$2.0\epsilon_{3} +$	$2w_3 +$	$3w_2$
$x_2 =$	0.0 +	$0.0\epsilon_1$ +	$4.0\epsilon_2$ +	$3.0\epsilon_3$ –	$3w_3 -$	$4w_2$
$x_1 =$	0.0 +	$0.0\epsilon_{1} + $	$1.0\epsilon_2$ +	$1.0\epsilon_3$ –	w_3 -	w_2

Done!

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Now that we have a Phase I algorithm and a variant of the simplex method that is guaranteed to terminate, we can summarize the main points of this chapter in the following theorem:

THEOREM

For an arbitrary linear program in standard form, the following statements are true:

- If there is no optimal solution, then the problem is either infeasible or unbounded.
- If a feasible solution exists, then a basic feasible solution exists.
- If an optimal solution exists, then a basic optimal solution exists.

DEFINITION/EXAMPLES OF DEGENERACY

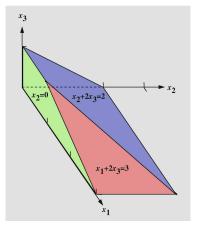
PERTURBATION METHOD

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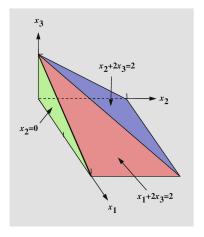
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The second pivoting rule we consider is called Bland's rule or smallest index rule.

It stipulates that:

Both the entering and the leaving variable be selected from their respective sets of choices by choosing the variable x_k with the smallest index k.

THEOREM

The simplex method always terminates provided that both the entering and the leaving variable are chosen according to Bland's rule.

Other rule called random selection rule.

Select at random from the set of possibilities.

Other rule called greatest increase rule.

Pick the entering/leaving pair so as to maximize the increase of the objective function over all other possibilities.

Too much computations.