

Kapitel 5.7:
Sortieren in Linearzeit
Algorithmen und Datenstrukturen
WS 2022/23

Prof. Dr. Sándor Fekete



Satz 5.6. *Für n Objekte x_1, \dots, x_n*



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wenn man die Objekte nur paarweise vergleichen kann.





Lower bound of $\Omega(n \log n)$ on sorting



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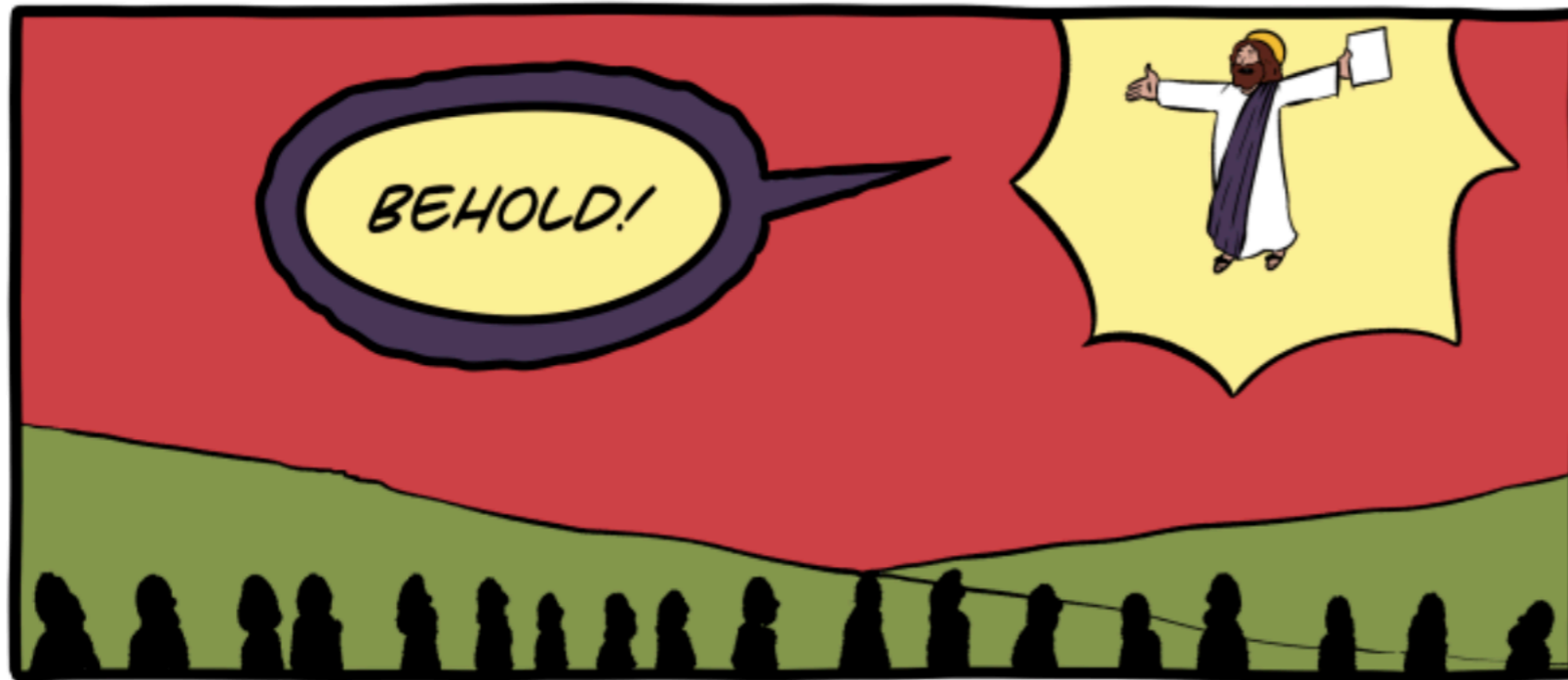
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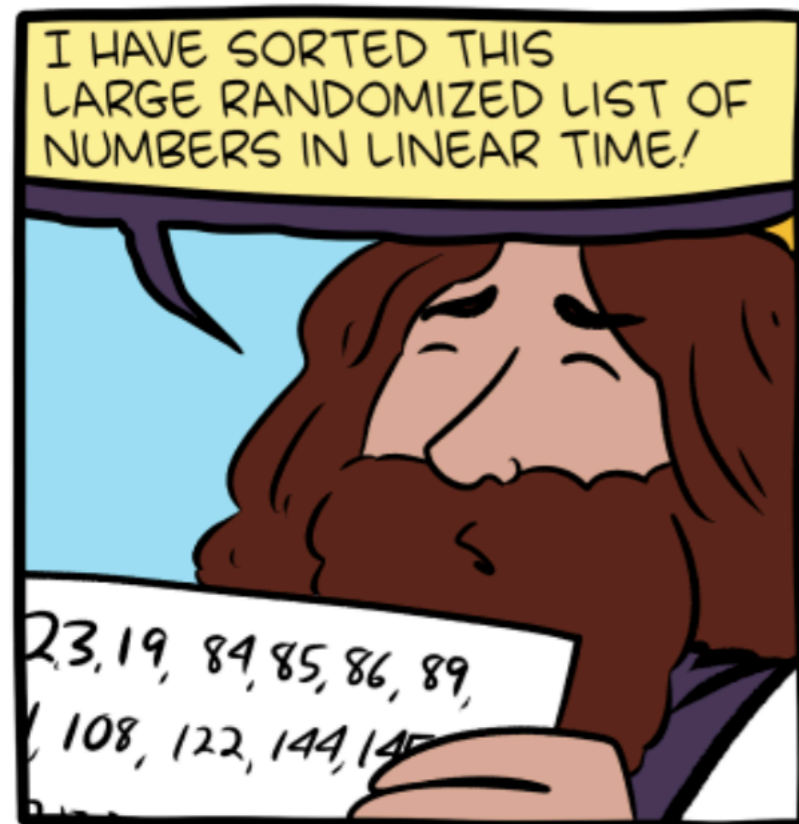
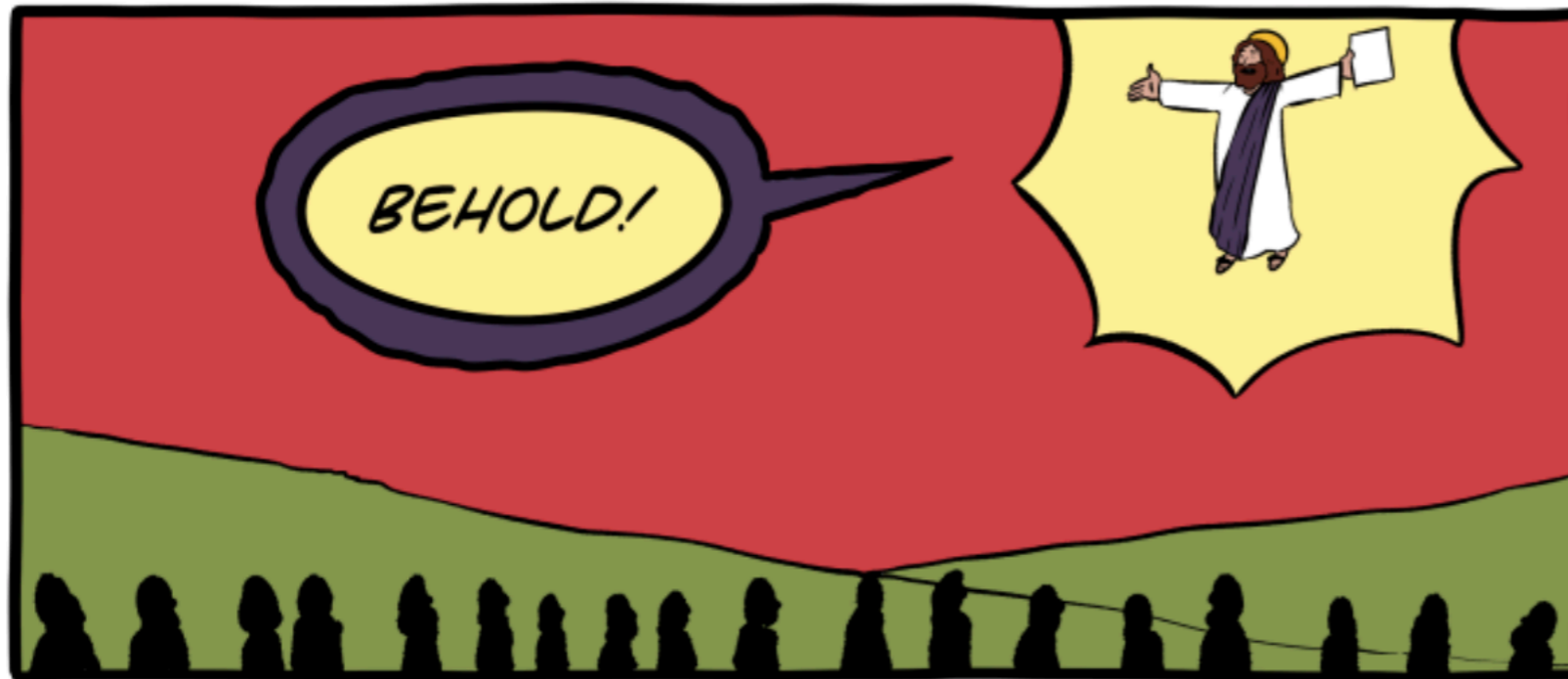
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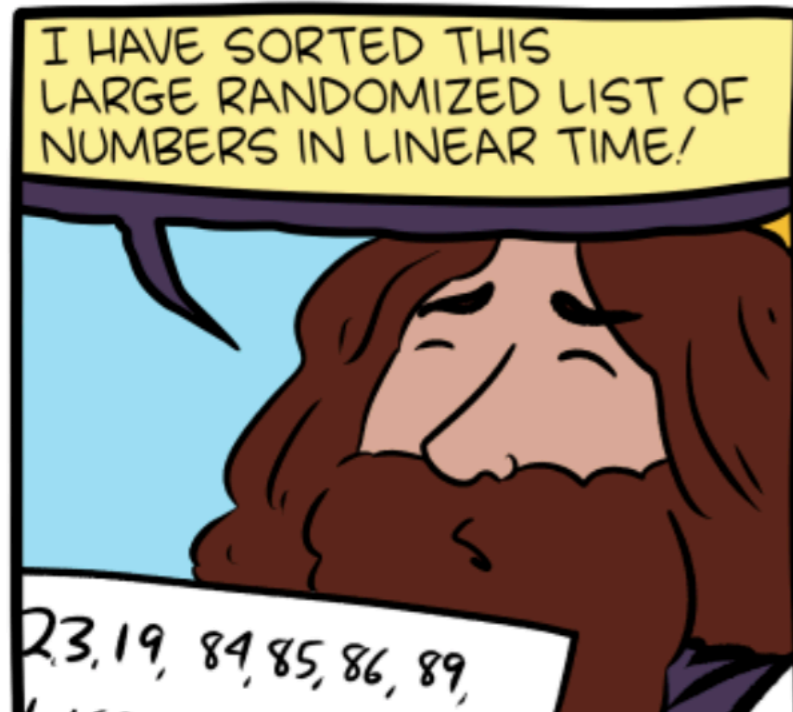
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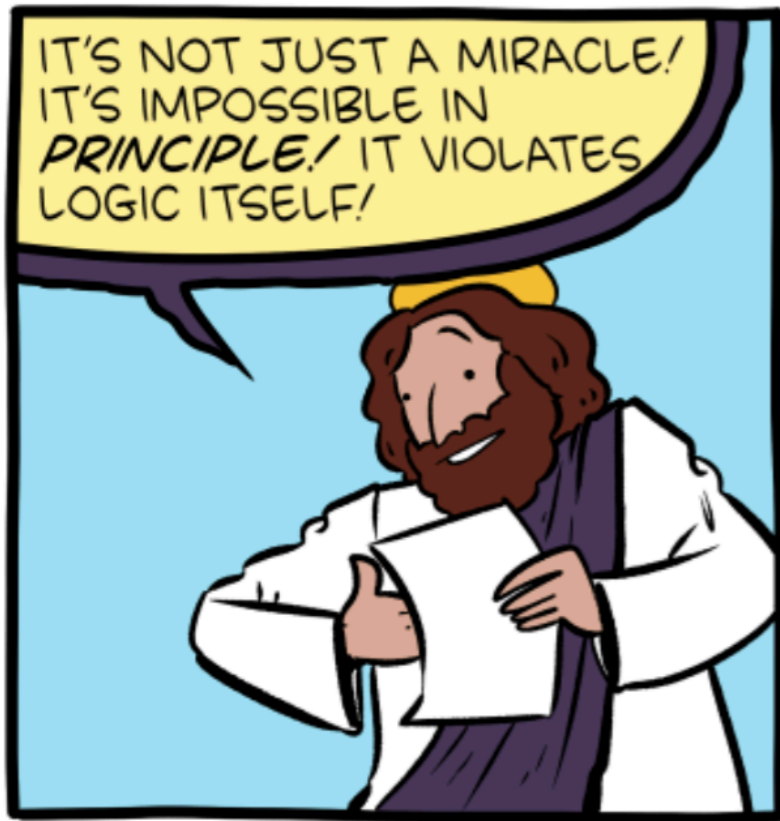
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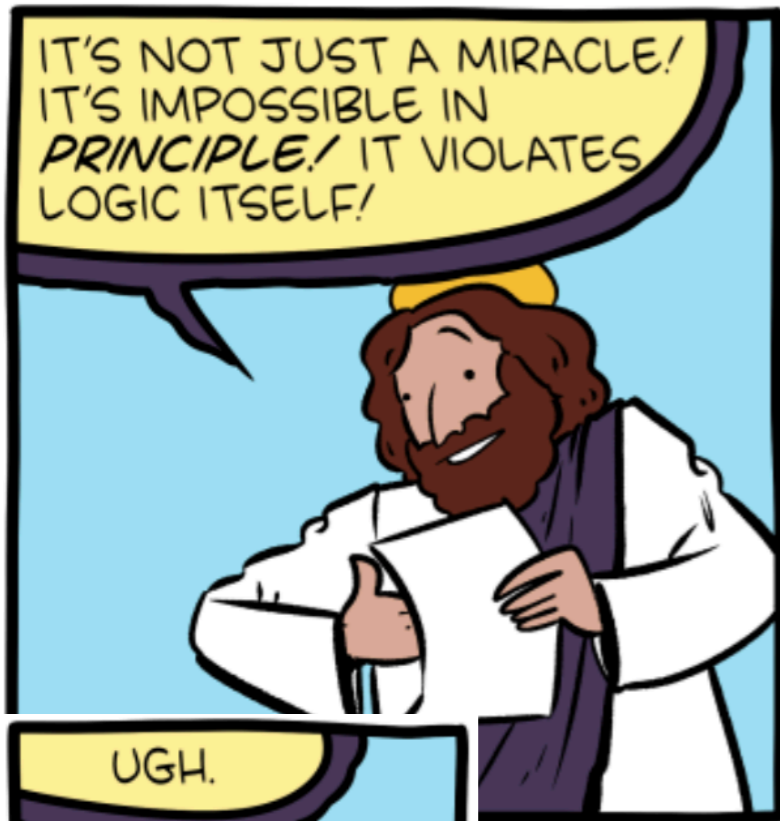
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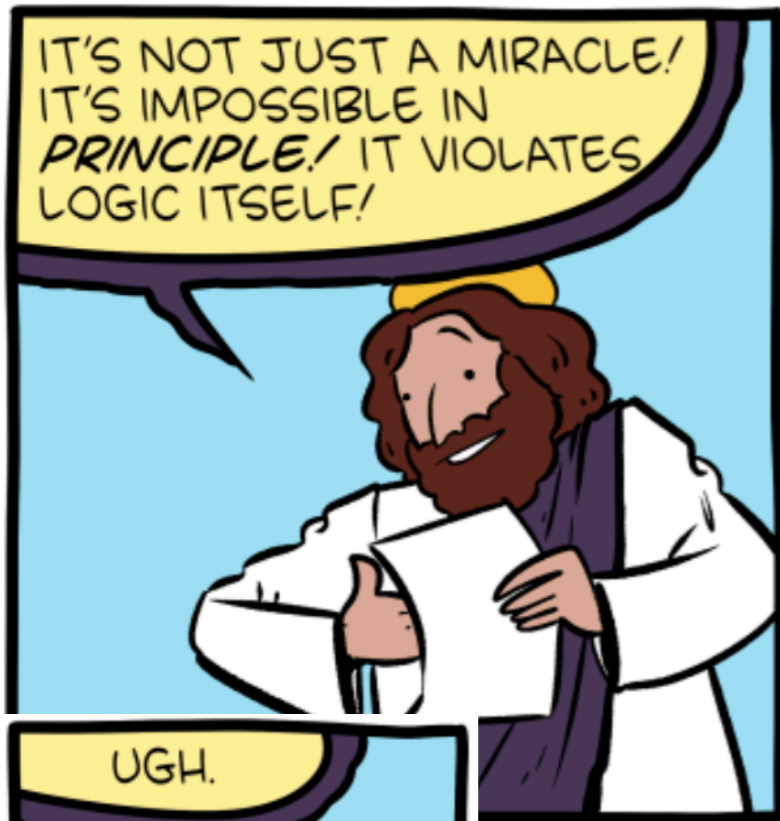
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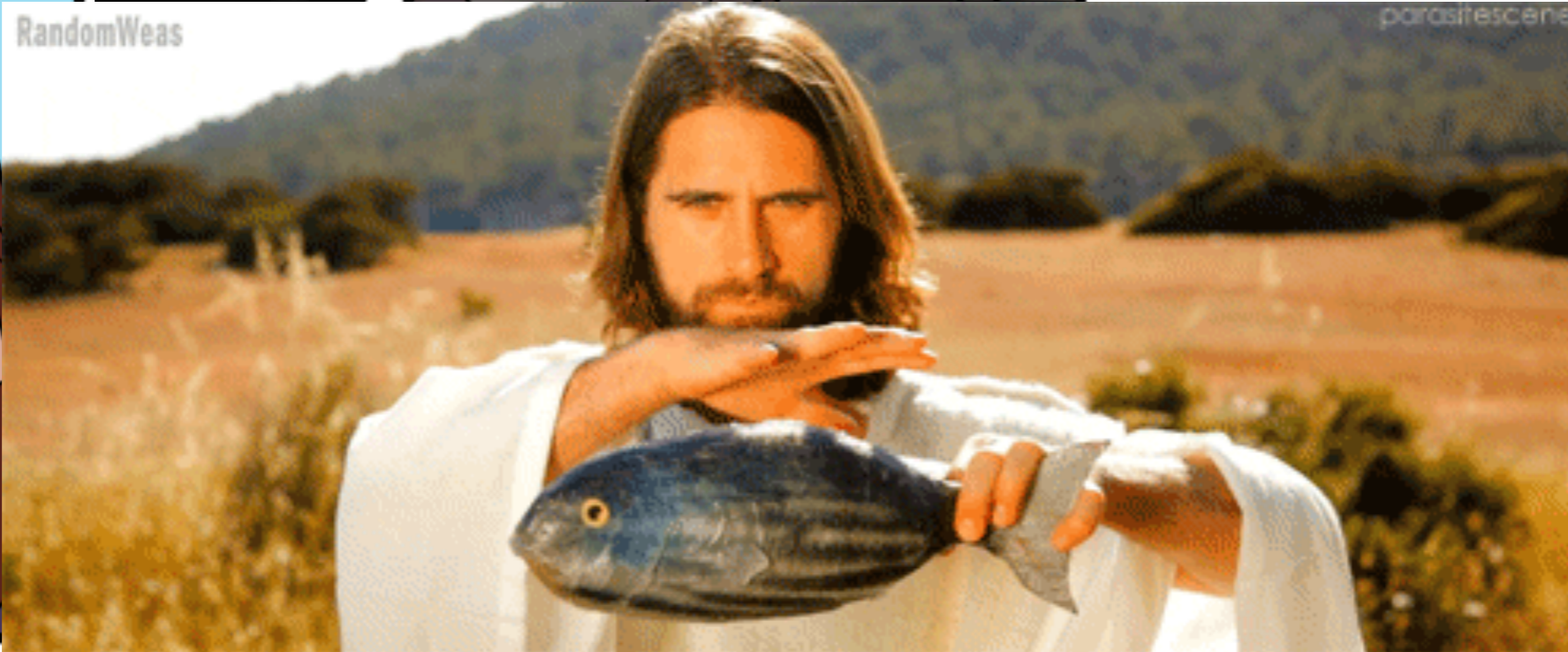
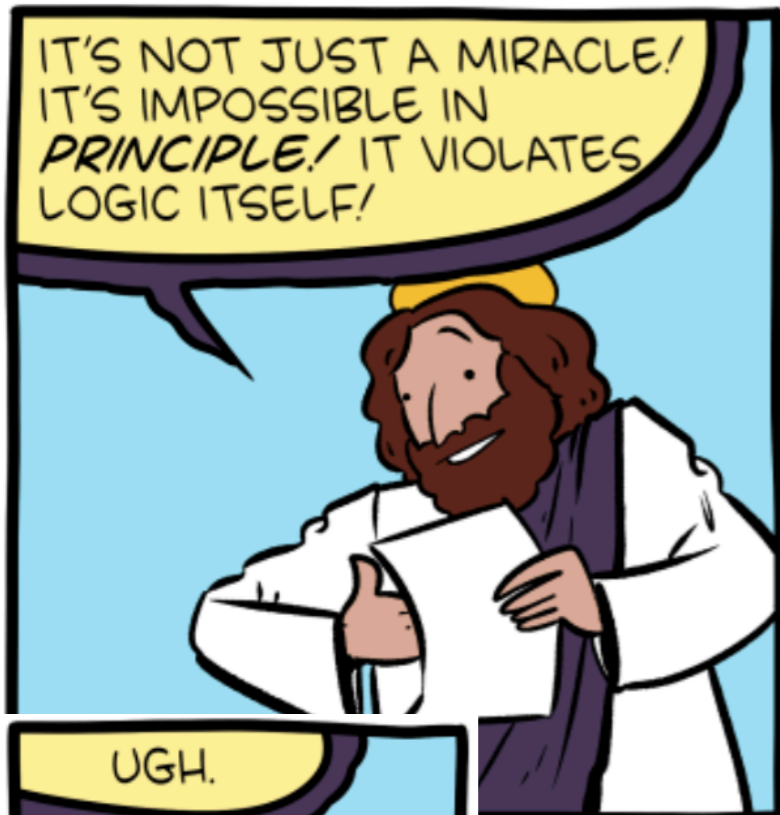
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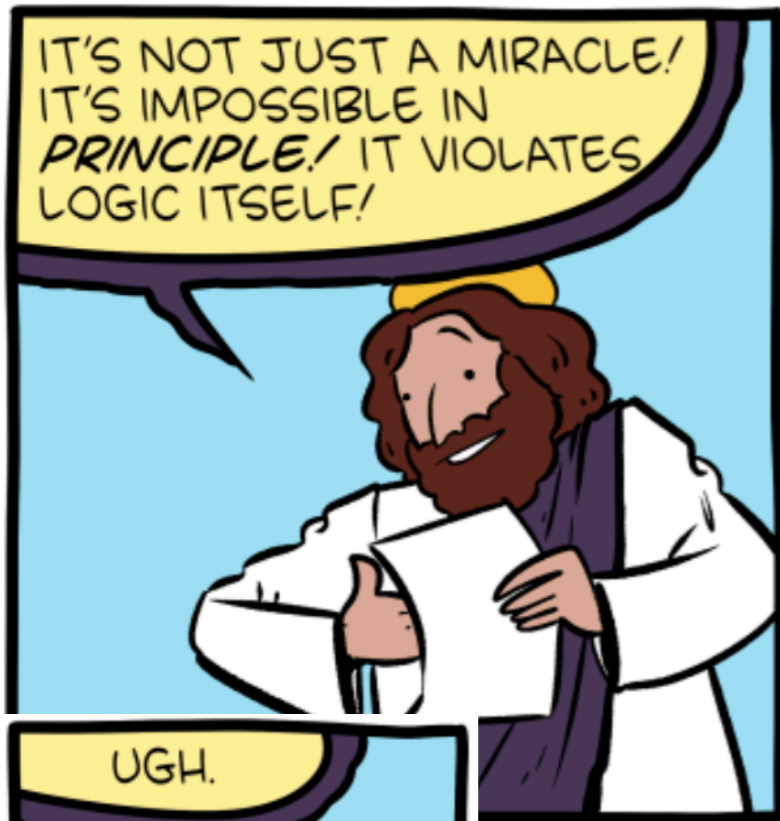
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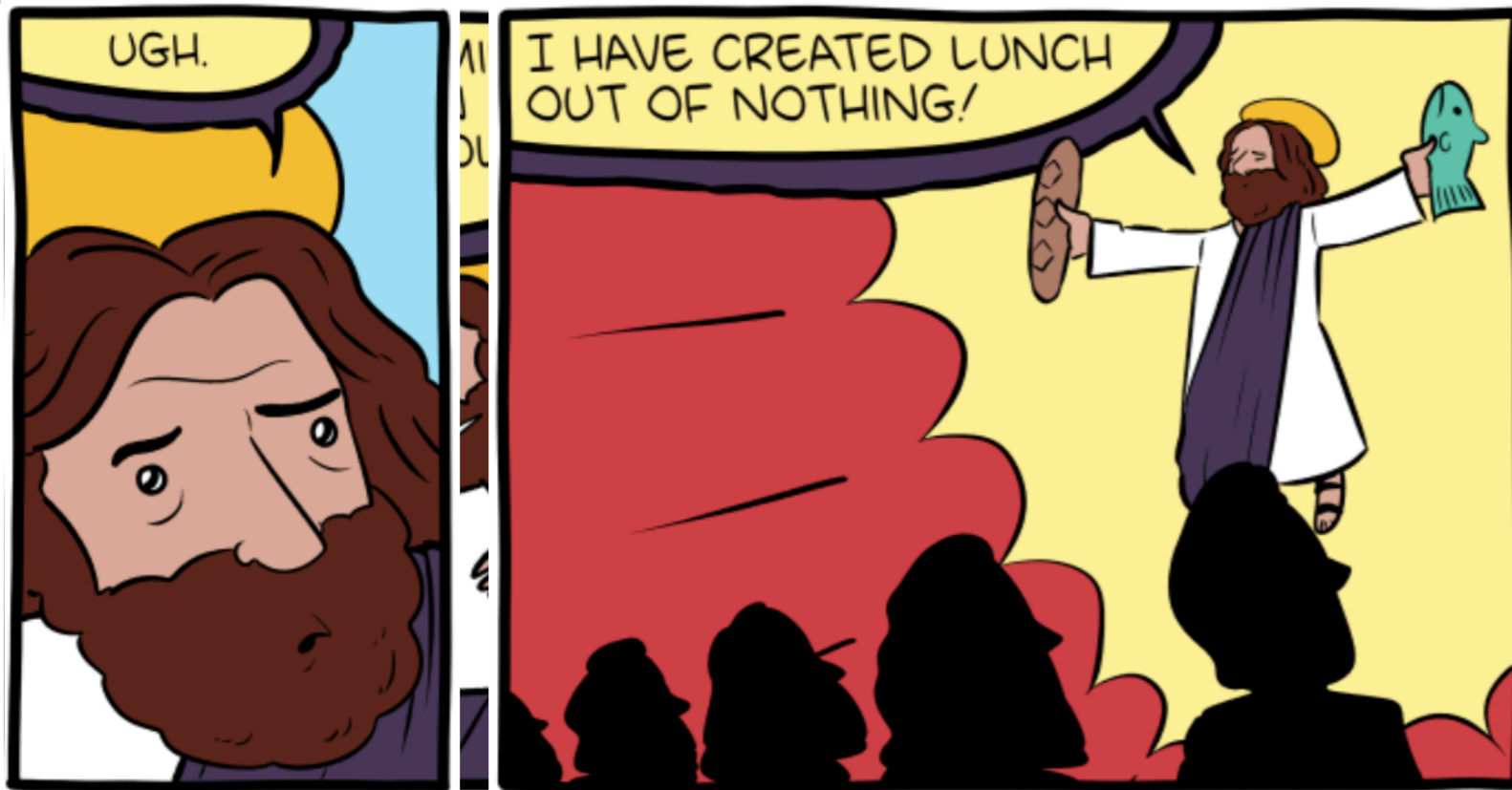
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Ab an die Tafel!

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