

Mergesort(A, p, r)

Input : Subarray von $A[1, \dots, n]$, der bei p beginnt und bei r endet ($A[p, \dots, r]$)
Output : Den sortierten Subarray

```
1 if (p < r) then
2   | q ← ⌊ $\frac{p+r}{2}$ ⌋
3   | Mergesort(A, p, q)
4   | Mergesort(A, q + 1, r)
5   | Merge(A, p, q, r)
6 end
```

Merge(A, p, q, r):

```
1  $n_1 \leftarrow q - p + 1$ 
2  $n_2 \leftarrow r - q$ 
3 create arrays  $L[1, \dots, n_1 + 1]$  and  $R[1, \dots, n_2 + 1]$ 
4 for (i ← 1 to  $n_1$ ) do
5   |  $L[i] \leftarrow A[p + i - 1]$ 
6 end
7 for (j ← 1 to  $n_2$ ) do
8   |  $R[j] \leftarrow A[q + j]$ 
9 end
10  $L[n_1 + 1] \leftarrow \infty$ 
11  $R[n_2 + 1] \leftarrow \infty$ 
12 i ← 1
13 j ← 1
14 for k ← p to r do
15   | if ( $L[i] \leq R[j]$ ) then
16     |  $A[k] \leftarrow L[i]$  i = i + 1
17   | end
18   | else
19     |  $A[k] \leftarrow R[j]$ 
20     | j ← j + 1
21   | end
22 end
```



(ii) Mergesort

A:	2	8	6	5		
	2	8		6	5	
	2		8	6	5	
	2	8		6		5
	2	8		5	6	
	2	5	6	8		

A[1]=2, A[2]=8, A[3]=6, A[4]=5

Mergesort(A, 1, 4)

q = floor(5/2) = 2

Mergesort(A, 1, 2)

q = floor(3/2) = 1

Mergesort(A, 1, 1) ← nichts passiert

Mergesort(A, 2, 2) ← " " "

Merge(A, 1, 1, 2) → 2 8 (*) ← nicht, sondern (xx)

Mergesort(A, 3, 4)

q = floor(5/2) = 3

Mergesort(A, 3, 3) ← nichts passiert

Mergesort(A, 4, 4) ← nichts passiert

Merge(A, 3, 3, 4) → 5 6 (xx)

Merge(A, 1, 2, 4) → 2 5 6 7

(*):

Merge (A, 1, 1, 2)
 $\uparrow \uparrow \uparrow$
 p q r

$$n_1 = q - p + 1 = 1 - 1 + 1 = 1$$

$$n_2 = r - q = 2 - 1 = 1$$

L[1, 2] R[1, 2]

$$i=1: L[1] = A[1+1-1] = A[1] = 2$$

$$j=1: R[1] = A[1+1] = A[2] = 8$$

$$L[2] = \infty$$

$$R[2] = \infty$$

$$i=1$$

$$j=1$$

$$k=1: L[1] \leq R[1]? (2 \leq 8? \text{ ja})$$

$$\hookrightarrow A[1] = L[1] = 2$$

$$i=2$$

$$k=2: L[2] \neq R[1] (\infty \neq 8)$$

$$\hookrightarrow A[2] = R[1] = 8$$

end

Somit: A[1]=2, A[2]=8

(**)

Merge (A, 3, 3, 4)
 $\uparrow \uparrow \uparrow$
 p q r

$$n_1 = q - p + 1 = 3 - 3 + 1 = 1$$

$$n_2 = r - q = 4 - 3 = 1$$

L[1, 2], R[1, 2]

$$i=1: L[1] = A[3+1-1] = A[3] = 6$$

$$j=1: R[1] = A[3+1] = A[4] = 5$$

$$L[2] = \infty$$

$$R[2] = 6$$

$$i = 1$$

$$j = 1$$

$$k = 3: L[1] \neq R[1]$$

$$\hookrightarrow A[3] = R[1] = 5$$

$$j = 2$$

$$k = 4: L[1] \leq R[2] \quad \checkmark \quad (6 \leq \infty)$$

$$\hookrightarrow A[4] = L[1] = 6$$

$$i = 2$$

Somit $A[3] = 5, A[4] = 6$

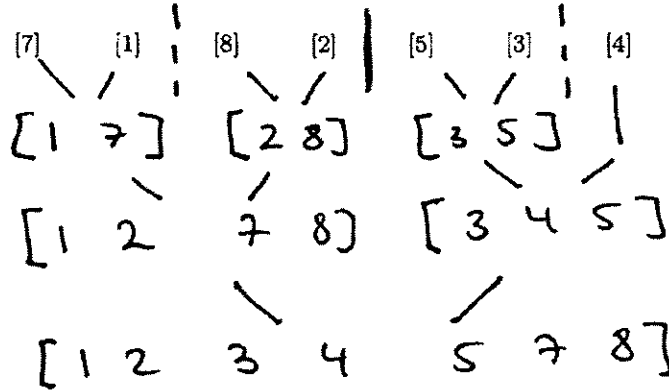
6. Aufgabe: Sortieren

9 Punkte

Sortiere die folgenden Zahlen mit dem in der Vorlesung vorgestellten Mergesort. Kennzeichne in jedem Schritt, welche Teilfolgen gemischt werden.

$$p = 1, q = 7$$

$$q = \left\lfloor \frac{p+r}{2} \right\rfloor = 4$$



$$\begin{aligned}
 \text{a) } T(n) &= 256 T\left(\frac{n}{4}\right) + n^3 \\
 &= \sum_{i=1}^{256} T\left(\frac{1}{4} \cdot n\right) + \Theta(n^3)
 \end{aligned}$$

$$\text{Also: } \alpha_i = \frac{1}{4}, i=1, \dots, 256$$

$$m = 256$$

$$k = 3$$

$$\sum_{i=1}^m \alpha_i^k = \sum_{i=1}^{256} \left(\frac{1}{4}\right)^3 = 256 \cdot \frac{1}{64} > 1$$

↳ 3. Fall, wir suchen also c mit $\sum_{i=1}^m \alpha_i^c = 1$

$$\sum_{i=1}^{256} \left(\frac{1}{4}\right)^c = 1$$

$$\Leftrightarrow 256 \left(\frac{1}{4}\right)^c = 1$$

$$\Leftrightarrow \left(\frac{1}{4}\right)^c = \frac{1}{256}$$

$$\Leftrightarrow 4^c = 256$$

$$\Leftrightarrow c = \log_4 256$$

$$\Leftrightarrow c = 4$$

$$\Rightarrow T(n) \in \Theta(n^c) = \Theta(n^4)$$

$$b) T(n) = 27 \cdot T\left(\frac{n}{3}\right) + n^3$$

$$\alpha_i = \frac{1}{3}, i=1, \dots, 27$$

$$m = 27$$

$$k = 3$$

$$\sum_{i=1}^m \alpha_i^k = \sum_{i=1}^{27} \left(\frac{1}{3}\right)^3 = 27 \cdot \frac{1}{27} = 1$$

→ Fall 2

$$\Rightarrow T(n) \in \Theta(n^k \log n) = \Theta(n^3 \log n)$$

$$c) T(n) = 3 T\left(\frac{n}{4}\right) + n^2$$

$$\alpha_i = \frac{1}{4}, i=1, \dots, 3$$

$$m = 3$$

$$k = 2$$

$$\sum_{i=1}^m \alpha_i^k = \sum_{i=1}^3 \left(\frac{1}{4}\right)^2 = 3 \cdot \frac{1}{16} < 1$$

→ Fall 1

$$\Rightarrow T(n) \in \Theta(n^k) = \Theta(n^2)$$