# Collaborative transmission in wireless sensor networks

Distributed Adaptive Beamforming

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1/65

#### **Overview and Structure**

- Introduction to context aware computing
- Wireless sensor networks
- Wireless communications
- Basics of probability theory
- Randomised search approaches
- Cooperative transmission schemes
- Distributed adaptive beamforming
  - Feedback based approaches
  - Asymptotic bounds on the synchronisation time
  - Alternative algorithmic approaches
  - Alternative Optimisation environments

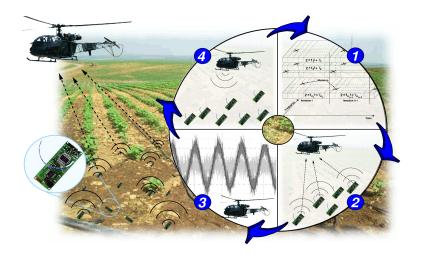
#### **Overview and Structure**

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#### **Outline**

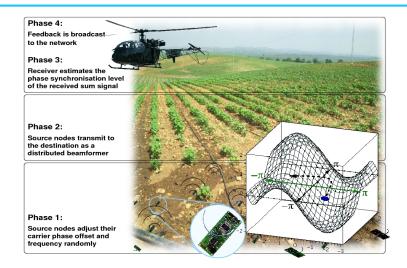
#### Feedback based distr. adaptive beamforming

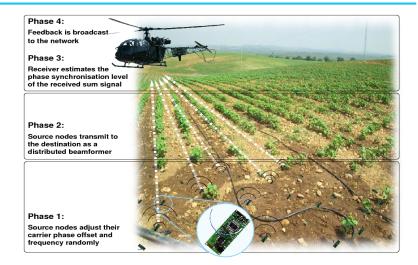
- Analysis of the problem scenario
  - Individual representation
  - Fitness function
  - Search space
  - Variation operators
- Analysis of the convergence time
  - An upper bound on the synchronisation performance
  - A lower bound on the synchronisation performance

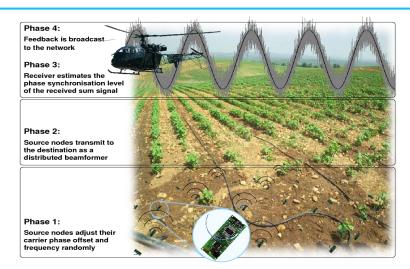


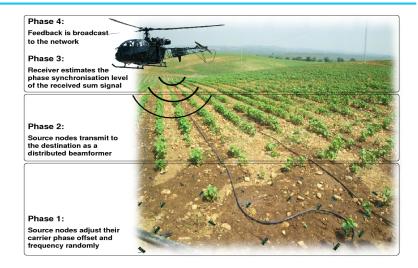


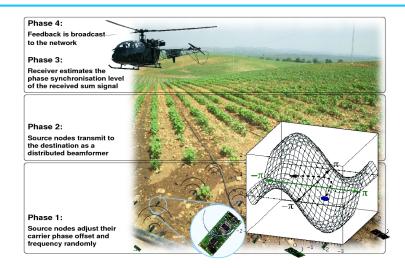
- 1-bit feedback based closed loop carrier synchronisation
  - Slow synchronisation
  - But: Computationally modest demands
  - Only: Adaptation of carrier phase based on binary feedback value
- Therefore: Well suited to be applied for WSNs











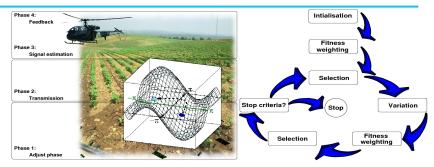


- Analysis of the underlying algorithmic problem
  - Precise mathematical understanding of the problem required
  - Modelling of
    - Search space
    - Optimisation aim
    - Representation of search points
    - Parameters that impact the synchronisation performance

#### **Outline**

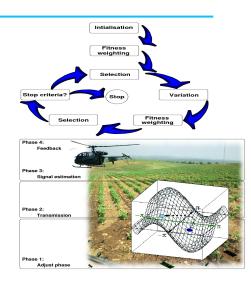
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- Observations
  - Iterative approach similar to evolutionary random search
    - New search points are requested by altering the carrier phases
    - Fitness function implemented by receiver feedback
    - Selection of individuals based on feedback values
    - Population size and offspring population size:  $\mu=
      u=1$

- Individual representation
  - Ordered set
  - Vector
  - Binary representation
- Fitness function
  - SNR
  - Simple distance
- Search space
  - Identical frequency
  - Distinct frequencies
- Variation operators
  - Mutation
  - Crossover



Analysis of the problem scenario

- Individual representation
  - Ordered set of phase and frequency pairs  $\gamma_i, f_i$

Advantage: Very near to the actual physical scenario

Disadvantage: Similarity measures between individuals not straightforward

• Vector  $V = v_1, \dots, v_{2n}$  of phases and/or frequencies

Advantage: Configurations as points in vector spaces, simple distance measure

Disadvantage: Representation very problem specific/untypical

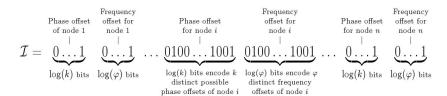
Binary representation of phase/frequency offsets

Advantage: Various results on binary search spaces in the

literature

Disadvantage: Hamming distance may not represent

neighbourhood similarities



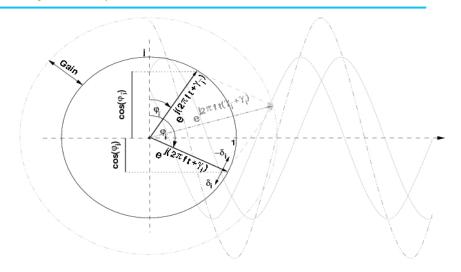
- Individual representation
  - Here: Binary representation of phase/frequency offsets
    - log(k) bits to represent k phase offsets
    - $log(\varphi)$  bits to represent  $\varphi$  frequency offsets
    - Configurations for all nodes concatenated
  - Phase and frequency offsets enumerated in ascending order
  - Neighbourhood: Gray encoded bit sequence to respect neighbourhood similarities

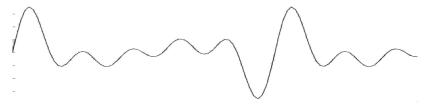


- Fitness function
  - Receiver estimates synchronisation quality of

$$\zeta_{\mathsf{sum}} = \Re\left(m(t)e^{j2\pi f_{\mathsf{c}}t}\sum_{i=1}^{n}\mathsf{RSS}_{i}e^{j(\gamma_{i}+\phi_{i}+\psi_{i})}\right)$$

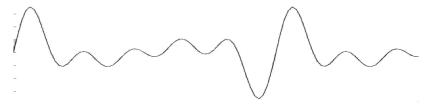
- SNR
- Numeric distance
- One bit feedback?





- Binary feedback
  - Minimum transmission load
  - Can be invested into higher redundancy schemes
  - Reduced information at source nodes
    - No adaptive operation
    - Less advanced optimisation schemes
    - Estimation of optimisation progress

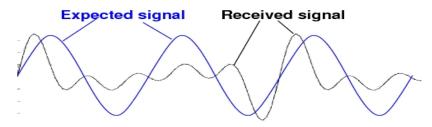
Analysis of the problem scenario



#### Fitness estimated by SNR:

- Calculate SNR of received sum signal
- Received signal strength above noise power
- Higher SNR interpreted as improved synchronisation quality
- Optimisation aim to reach minimum required SNR

Analysis of the problem scenario



#### Fitness estimated by simple distance:

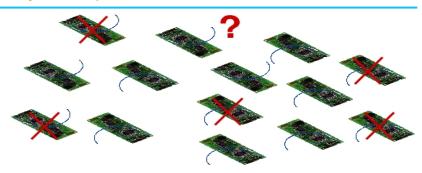
- Calculate surface between  $\zeta_{opt}$  and  $\zeta_{sum}$
- Smaller surface → better synchronisation quality
- Optimum signal:

$$\zeta_{\mathsf{opt}} = \Re\left(m(t)\mathsf{RSS}_{\mathsf{opt}}e^{j(2\pi f_c t + \gamma_{\mathsf{opt}} + \phi_{\mathsf{opt}} + \psi_{\mathsf{opt}})}\right)$$

$$\zeta_{\text{opt}} = \Re\left(m(t)\mathsf{RSS}_{\text{opt}}e^{j(2\pi f_c t + \gamma_{\text{opt}} + \phi_{\text{opt}} + \psi_{\text{opt}})}\right)$$

- Transmit sequence m(t) (preconditioned)
- Transmit frequency  $f_c$  (preconditioned)
- Average transmit power  $P_{\text{avg}}$  (preconditioned)
- Gain G<sub>i</sub>, G<sub>receiver</sub> (preconditioned)
- Distance d to network (Estimated by RTT)
- Number of transmitting nodes  $n \rightarrow ???$
- $RSS_{opt} = n \cdot \left( P_{avg} \cdot \left( \frac{\lambda}{2\pi \cdot d} \right)^2 \cdot G_i \cdot G_{receiver} \right)$

Analysis of the problem scenario



#### Estimate the count of transmitting nodes:

- Possible to estimate count of transmitting nodes
- From superimposed signal of simultaneously transmitting nodes<sup>1</sup>

<sup>&</sup>lt;sup>1</sup>A.Krohn, Superimposed Radio Signals for Wireless Sensor Networks, PhD thesis, 2007 Stephan Sigg Collaborative transmission in wireless sensor networks

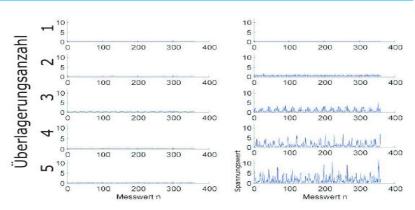
Analysis of the problem scenario



Estimate the count of transmitting nodes <sup>2</sup>

<sup>&</sup>lt;sup>2</sup> A.Krohn, Superimposed Radio Signals for Wireless Sensor Networks, PhD thesis, 2007 Stephan Sigg Collaborative transmission in wireless sensor networks

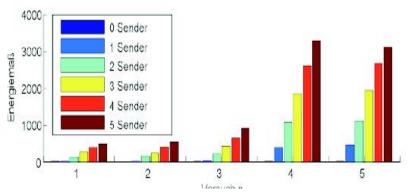
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Analysis of the problem scenario



Estimate the count of transmitting nodes <sup>4</sup>

<sup>&</sup>lt;sup>4</sup> A.Krohn, Superimposed Radio Signals for Wireless Sensor Networks, PhD thesis, 2007 Stephan Sigg Collaborative transmission in wireless sensor networks

Analysis of the problem scenario

catsächliche Anzahl

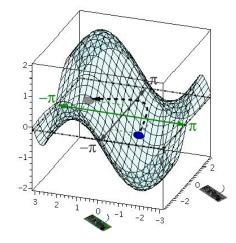
#### Geschätzte Anzahl

	0	1	2	3	4	5
0	287	0	0	0	0	0
1	0	327	3	0	0	0
2	0	0	330	0	0	0
3	0	0	32	321	14	19
4	0	0	11	69	211	124
5	0	0	0	17	39	220

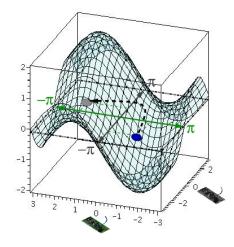
Estimate the count of transmitting nodes <sup>5</sup>

<sup>&</sup>lt;sup>5</sup> A. Krohn, Superimposed Radio Signals for Wireless Sensor Networks, PhD thesis, 2007 Stephan Sigg Collaborative transmission in wireless sensor networks

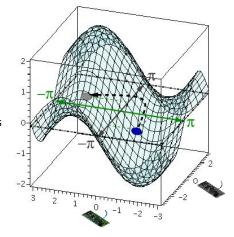
- Search space
  - Optimisation performance dependent on search space
  - Global or local optima?



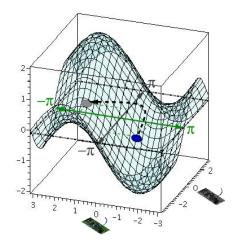
- Search space
  - Feedback function not unimodal
  - In two global optima, carrier signals are shifted by fixed amount
  - Fitness function weak multimodal
    - Many global optima
    - No local optima



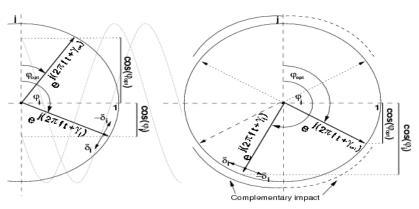
- Search space
  - Identical transmit frequencies
  - Distinct transmit frequencies



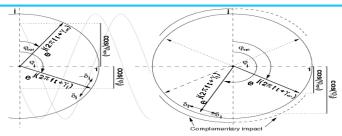
- Identical transmit frequencies:  $e^{j(2\pi ft + \gamma_i)}$ ;  $\forall i \in \{1, ..., n\}$ 
  - Local optimum:  $\exists$  search point  $s_{\overline{c}} \neq s_{\text{opt}}$  with
  - All small phase modulations decrease fitness value
  - Smallest possible modification: Single carrier signal altered



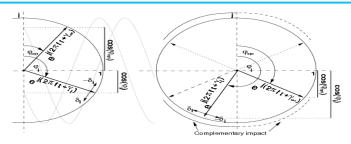
Analysis of the problem scenario



• Fitness dependent on distance  $|\cos(\varphi_{\text{opt}}) - \cos(\varphi_i)|$ 

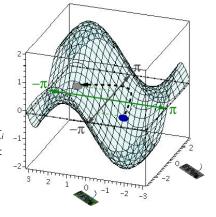


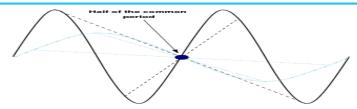
- Phase shift of  $\delta_i \neq 0$  alters the fitness value
- For some t the fitness increases while for others it decreases.
- Assume  $(\varphi_i + \delta_i) \varphi_{\sf opt} < 180^\circ$  and  $\varphi_i > \varphi_{\sf opt}$
- For  $[\varphi_i>180^\circ \land \varphi_{\sf opt}<180^\circ]$  or  $[\varphi_i>360^\circ \land \varphi_{\sf opt}<360^\circ]$ 
  - Contribution to  $\mathcal{F}$  zero
- Else:  $\delta_i$  has either always positive or always negative impact



- Compared to sopt
  - No configuration short of the optimum configuration s<sub>i</sub> = s<sub>opt</sub> exists
  - For which distance is increased for phase offset  $\delta_i$
  - regardless of the sign of  $\delta_i$
- No local optima

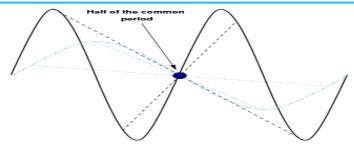
- Distinct transmit frequencies:  $e^{j(2\pi f_i t + \gamma_i)}$ ;  $\forall i \in \{1, ..., n\}$ 
  - Consider phase offset between two signals:
    - Modified signal component  $\zeta_i$
    - Nearest global optimum  $\zeta_{opt}$





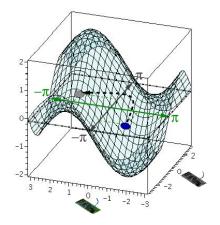
- Distinct transmit frequencies:  $e^{j(2\pi f_i t + \gamma_i)}$ ;  $\forall i \in \{1, ..., n\}$ 
  - Feedback function not affected by phase modifications only
  - ullet Periodic function: Reflection in half of common period  $\Phi$
  - For every positive contribution also negative contribution

$$\begin{array}{ll} & e^{j(2\pi(f_1)t \mod \varPhi + \gamma_1)} - e^{j(2\pi ft \mod \varPhi)} \\ = & - \left( e^{j(2\pi(f_1)t' \mod \varPhi + \gamma_1)} - e^{j(2\pi ft' \mod \varPhi)} \right) \end{array}$$



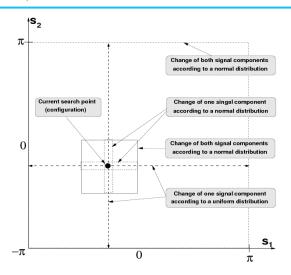
- Distinct transmit frequencies:  $e^{j(2\pi f_i t + \gamma_i)}$ ;  $\forall i \in \{1, \dots, n\}$ 
  - signal quality is not affected by phase adaptations when frequencies are unsynchonised
  - without frequency synchronisation, phase synchronisation alone is useless in order to improve the signal quality
- In both cases no local optima but several global optima

- Variation operators
  - Mutation
  - Crossover



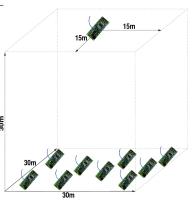
- Variation operators Mutation
  - Small modifications on individuals
  - Target individuals with small distance more probable
  - Phase modification of one or more carrier signals  $\zeta_i$
  - Design parameters:
    - Count of altered carrier signal components
    - Method for alteration of a single carrier

- Variation operators Mutation
  - Count of altered carrier signal components
    - Fixed number (how to implement in sensor network?)
    - Random number (Probability for each node)
  - Method for alteration of a single carrier
    - Neighbourhood bounds vs. Probability distribution
    - Uniform vs. Normal
    - Standard deviation  $\sigma$  (search neighbourhood)
    - Mean  $\mu$  (search direction)



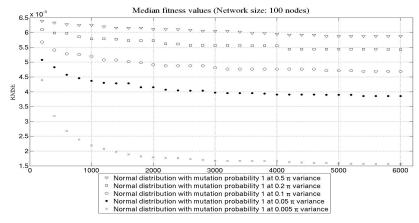
Analysis of the problem scenario

Value
30 <i>m</i> × 30 <i>m</i>
(15m, 15m, 30m)
stationary nodes
$f_{base} = 2.4 \text{ GHz}$
$P_{t imes}=1\;\mathrm{mW}$
$G_{tx}=0 \text{ dB}$
$G_{rx}=0 \text{ dB}$
6000
10
$-103~\mathrm{dBm}$
$P_{tx} \left(\frac{\lambda}{2\pi d}\right)^2 G_{tx} G_{rx}$



Variation operators - Mutation - example

Analysis of the problem scenario



Variation operators - Mutation - example

- Variation operators Crossover
  - Not yet considered in the literature
  - ullet (1 + 1)-EA straightforward as it consides one individual at a time
  - Multiple individuals possible by
    - Simultaneous transmission on distinct transmit signals
    - 2 Time-shifted transmission of several individuals

- Summary
  - 1-bit feedback based phase synchronisation always converges<sup>6</sup>
  - We can now come to the same result:
    - No local optima in the search space
    - Algorithm does never accept worse points
  - But: What is the expected time to reach an optimum?

<sup>&</sup>lt;sup>6</sup>R. Mudumbai, J. Hespanha, U. Madhow, G. Barriac: Distributed transmit beamforming using feedback control. IEEE Transactions on Information Theory (In review)

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#### Feedback based distr. adaptive beamforming

- Analysis of the problem scenario
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  - Fitness function
  - Search space
  - Variation operators
- Analysis of the convergence time
  - An upper bound on the synchronisation performance
  - A lower bound on the synchronisation performance

Analysis of the convergence time

#### Assumptions:

- Network of n nodes
- Each node changes the phase of its carrier signal with probability  $\frac{1}{n}$
- Carrier phase altered uniformly at random from  $[0,2\pi]$
- Feedback function  $\mathcal{F}: \zeta^*_{\mathsf{sum}} \to \mathbb{R}$  maps

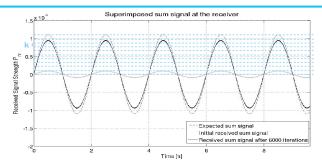
$$\zeta_{\mathsf{sum}} = \Re\left(m(t)e^{j2\pi f_c t}\sum_{i=1}^n \mathsf{RSS}_i e^{j(\gamma_i + \phi_i + \psi_i)}\right)$$

to a real-valued fitness score.

Possible feedback:

$$\mathcal{F}\left(\zeta_{\mathsf{sum}}\right) = \int_{t=0}^{2\pi} \left|\zeta_{\mathsf{sum}} - \zeta_{\mathsf{opt}}\right|$$

Analysis of the convergence time



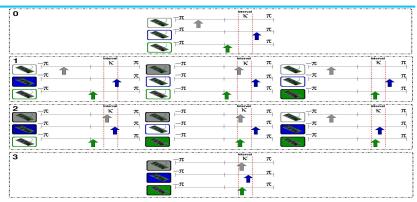
#### Optimisation aim:

- Achieve maximum relative phase offset of  $\frac{2\pi}{k}$
- Between any two carrier signals
- For arbitrary k
- Divide phase space into k intervals of width  $\frac{2\pi}{k}$

Analysis of the convergence time

- An upper bound on the synchronisation performance
  - Upper bound by method of fitness based partitions
  - Value of fitness function increases with number of carrier signals  $\zeta_i$  that share same interval for phase offset  $\gamma_i$
  - Assume, that  $\kappa \in [1, k]$  is interval with most carrier phases
  - Worse fitness values are not accepted
  - Count iterations required for all carrier signals to change to interval  $\kappa$ 
    - Note: We disregard positive possibilities to reach any other optimum
    - Possible since only upper bound is calculated

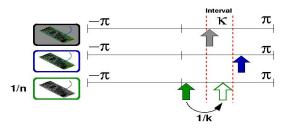
Analysis of the convergence time



### Divide values of the fitness function into k partitions:

•  $L_1, \ldots, L_n$ , depending on the count of carrier signals with phase offset in  $\kappa$ 

Analysis of the convergence time



#### Divide values of the fitness function into k partitions:

- Probability to adapt phase to specific interval:  $\frac{1}{\iota}$
- Probability to reach at least to next partition

$$\frac{1}{k} \cdot (n - L_i) \cdot \frac{1}{n}$$

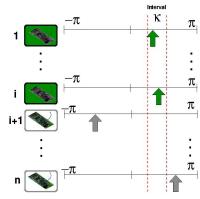
Analysis of the convergence time

• In partition i, one of

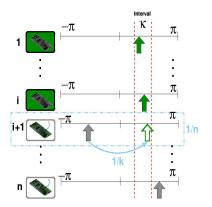
$$\left(\begin{array}{c} n-i\\ 1 \end{array}\right)=n-i$$

carrier signals suffice to improve the fitness value

- this happens with probability  $\frac{1}{n} \cdot \frac{1}{k}$
- At least one shall be correctly altered while all other n-1 signals remain unchanged



Analysis of the convergence time

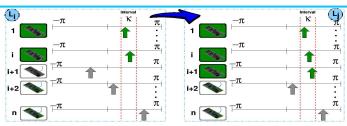


- Alter 1 carrier and keep n-1 signals
- This happens with probability

$$\begin{pmatrix} n-i \\ 1 \end{pmatrix} \cdot \frac{1}{n} \cdot \frac{1}{k} \cdot \left(1 - \frac{1}{n}\right)^{n-1}$$

$$= \left(\frac{n-i}{n \cdot k}\right) \cdot \left(1 - \frac{1}{n}\right)^{n-1}$$

Analysis of the convergence time



Since

$$\left(1-\frac{1}{n}\right)^n<\frac{1}{e}<\left(1-\frac{1}{n}\right)^{n-1}$$

• Probability that  $L_i$  is left for partition j, j > i:

$$P[L_i] \ge \frac{n-i}{n \cdot e \cdot k}$$

Analysis of the convergence time

• Expected number of iterations to change layer bounded from above by  $P[L_i]^{-1}$ :

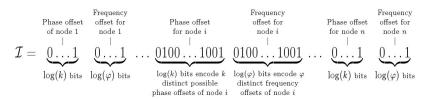
$$E[T_{\mathcal{P}}] \leq \sum_{i=0}^{n-1} \frac{e \cdot n \cdot k}{n-i}$$

$$= e \cdot n \cdot k \cdot \sum_{i=1}^{n} \frac{1}{i}$$

$$< e \cdot n \cdot k \cdot (\ln(n) + 1)$$

$$= O(n \cdot k \cdot \log n)$$

Analysis of the convergence time



- A lower bound on the synchronisation performance
  - We utilise the method of the expected progress
  - After initialisation, phases of carrier signals are identically and independently distributed.
  - Each bit in the binary representation of search point  $s_{\zeta}$  has equal probability to be 1 or 0.

Analysis of the convergence time

																	$\sum$
$\mathcal{I}_i =$	1	0	1	1	0	1	1	1	1	0	1	1	0	1	1	1	
$\mathcal{I}_i = \ \mathcal{I}_{ ext{opt}} = $	1	0	1	0	1	1	0	1	1	0	1	0	1	1	0	1	
$h(\mathcal{I}_i, \mathcal{I}_{ ext{opt}}) =$	0	0	0	1	1	0	1	0	0	0	0	1	1	0	1	0	6

• Probability to start with hamming distance  $h(s_{\text{opt}}, s_{\zeta}) \leq I$ ;  $I \ll n \cdot \log(k)$  to global optima  $s_{\text{opt}}$  at most

$$P[h(s_{\text{opt}}, s_{\zeta}) \leq I] = \sum_{i=0}^{I} \binom{n \cdot \log(k)}{n \cdot \log(k) - i} \cdot \frac{k}{2^{n \cdot \log(k) - i}}$$

$$\leq \frac{(n \cdot \log(k))^{I+2}}{2^{n \cdot \log(k) - I}}$$

Analysis of the convergence time

																	$\sum$
$\mathcal{I}_i =$	1	0	1	1	0	1	1	1	1	O	1	1	0	1	1	1	
${\mathcal I}_i = \ {\mathcal I}_{ m opt} =$	1	O	1	O	1	1	O	1	1	O	1	O	1	1	O	1	
$h(\mathcal{I}_i, \mathcal{I}_{\mathrm{opt}}) =$																	6

$$P[h(s_{\text{opt}}, s_{\zeta}) \le I] \le \frac{(n \cdot \log(k))^{l+2}}{2^{n \cdot \log(k) - I}}$$

• Count of configurations with *i* bit errors to s<sub>opt</sub>:

$$\left(\begin{array}{c} n \cdot \log(k) \\ n \cdot \log(k) - i \end{array}\right)$$

- Probability for all these bits to be correct:  $\frac{1}{2^{n \cdot \log(k) i}}$
- Count of global optima: k

Analysis of the convergence time

$$\mathcal{I} = \underbrace{0 \dots 1}_{\text{log}(k) \text{ bits}} \underbrace{\log(\varphi) \text{ bits}}_{\text{offset for of seeks offset of of seeks of seeks}}^{\text{Frequency}} \underbrace{0 \text{ fiset for phase offset of for node } i}_{\text{phase offset of node } i} \underbrace{0 \dots 1}_{\text{phase offset of of node } i} \underbrace{0 \dots 1}_{\text{phase offset of node } i} \underbrace{0 \dots 1}_{\text{phase o$$

$$P[h(s_{\text{opt}}, s_{\zeta}) \leq I] = \sum_{i=0}^{I} \binom{n \cdot \log(k)}{n \cdot \log(k) - i} \cdot \frac{k}{2^{n \cdot \log(k) - i}}$$

$$\leq \frac{(n \cdot \log(k))^{I+2}}{2^{n \cdot \log(k) - I}}$$

• This means that with high probability (w.h.p.) the hamming distance to the nearest global optimum is at least /.

Analysis of the convergence time

- Use method of expected progress to calculate lower bound:
- $(s_{\zeta}, t)$  denotes that  $s_{\zeta}$  is achieved after t iterations
- Assume Progress measure  $\Lambda: \mathbb{B}^{n \cdot \log(k)} o \mathbb{R}_0^+$
- $\Lambda(s_{\zeta},t) < \Delta$ : Global optimum not found in first t iterations
- ullet For every  $t\in {
  m I\! N}$  we have

$$E[T_{\mathcal{P}}] \geq t \cdot P[T_{\mathcal{P}} > t]$$

$$= t \cdot P[\Lambda(s_{\zeta}, t) < \Delta]$$

$$= t \cdot (1 - P[\Lambda(s_{\zeta}, t) \geq \Delta])$$

Analysis of the convergence time

$$E[T_{\mathcal{P}}] \geq t \cdot (1 - P[\Lambda(s_{\zeta}, t) \geq \Delta])$$

• With the help of the Markov-inequality we obtain

$$P[\Lambda(s_{\zeta},t) \geq \Delta] \leq \frac{E[\Lambda(s_{\zeta},t)]}{\Lambda}$$

and therefore

$$E[T_{\mathcal{P}}] \geq t \cdot \left(1 - \frac{E[\Lambda(s_{\zeta}, t)]}{\Delta}\right)$$

Obtain lower bound by providing expected progress after t iterations

#### Analysis of the convergence time

Probability for I bits to correctly flip at most

$$\left(1 - \frac{1}{n \cdot \log(k)}\right)^{n \cdot \log(k) - l} \cdot \left(\frac{1}{n \cdot \log(k)}\right)^{l} \le \frac{1}{(n \cdot \log(k))^{l}}$$

Probability that no correct but remaining / bits flip:

$$\left(1 - \frac{1}{n \cdot \log(k)}\right)^{n \cdot \log(k) - l}$$

- I bits mutate with probability  $\left(\frac{1}{n \cdot \log(k)}\right)^I$
- Expected progress in one iteration:

$$E[\Lambda(s_{\zeta},t),\Lambda(s_{\zeta'},t+1)] \leq \sum_{i=1}^{l} \frac{i}{(n \cdot \log(k))^{i}} < \frac{2}{n \cdot \log(k)}$$

• Expected progress in t iterations:  $\leq \frac{2t}{n \cdot \log(k)}$ 

Analysis of the convergence time

- Choose  $t = \frac{n \cdot \log(k) \cdot \Delta}{4} 1$
- Double of expected progress still smaller than  $\Delta$ .
- With Markov inequality: Progress not achieved with prob.  $\frac{1}{2}$ .
- Expected optimisation time bounded from below by

$$E[T_{\mathcal{P}}] \geq t \cdot \left(1 - \frac{E[\Lambda(s_{\zeta}, t)]}{\Delta}\right)$$

$$\geq \frac{n \cdot \log(k) \cdot \Delta}{4} \cdot \left(1 - \frac{\frac{2 \cdot n \cdot \log(k)}{4 \cdot n \cdot \log(k)} \cdot \Delta}{\Delta}\right)$$

$$= \Omega(n \cdot \log(k) \cdot \Delta)$$

• With  $\Delta = k \cdot \frac{\log(n)}{\log(k)}$ : Same order as upper bound:

$$E[T_{\mathcal{P}}] = \Theta(n \cdot k \cdot \log(n))$$