

Evaluation of the 6TiSCH Network Formation

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Communication Technologies, Munich, Germany



Industrial Wireless Sensor Networks (IWSNs)



Targeted applications:
Process monitoring and control

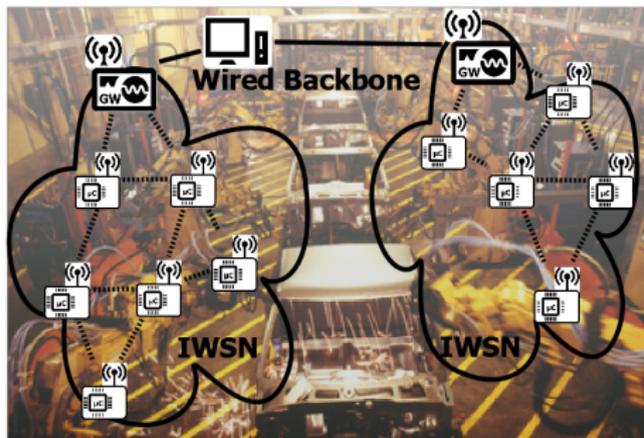
Strict requirements:

- Reliability up to 99,999%
- Lifetime > 5 years
- Latency: tens of milliseconds

Main characteristics of a IWSN:

- mesh, multi-hop, lossy network
- harsh environment
- a gateway and up to 100 resource constrained nodes
- specific designed communication protocols

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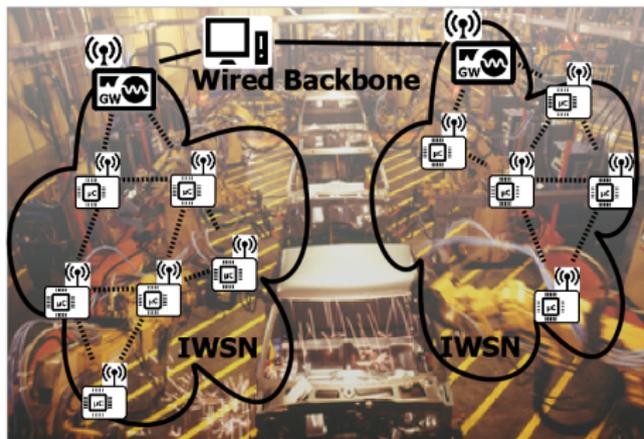
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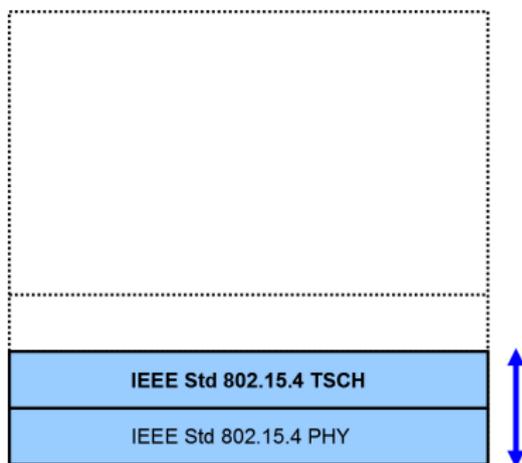
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6TiSCH-Stack: IETF Suite of Protocols for Industrial IoTs

Applications		
CoAP (OSCORE)	6LoWPAN ND	RPL
UDP	ICMPv6	
IPv6		
6LoWPAN HC / 6LoRH HC		
6top (6TiSCH)		
IEEE Std 802.15.4 TSCH		
IEEE Std 802.15.4 PHY		

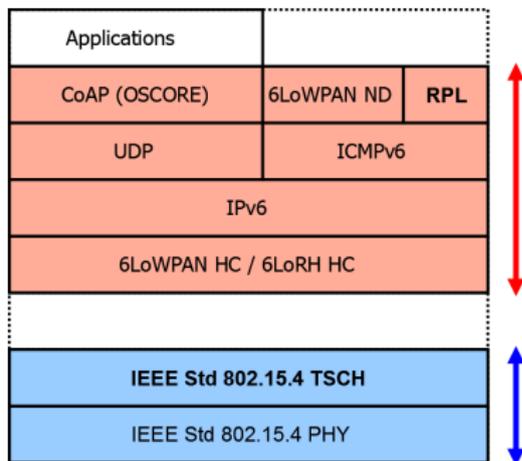
- Upper layers: IPv6-connectivity
 - 6LoWPAN, IPv6, CoAP etc.
 - *RPL as distributed routing protocol*
- Glueing together: 6top protocol proposed by *IETF 6TiSCH*
 - assignment of communication links
 - *definition of bootstrapping procedures*
- At the bottom: IEEE 802.15.4-2015
 - 2.4 GHz low-power radio
 - *MAC: Time Slotted Channel Hopping for industrial performance*

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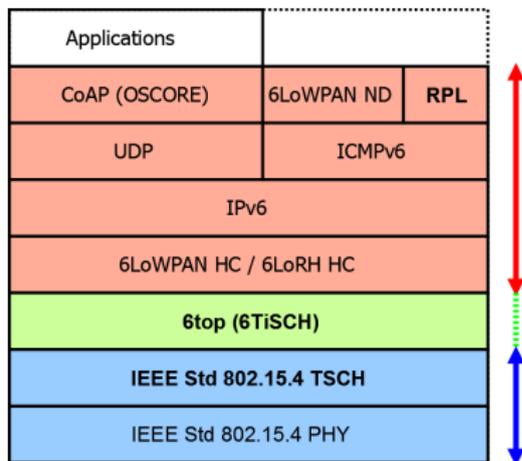
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Scenario:

- Mesh, multi-hop network for industrial wireless
- 6TiSCH-Stack as IETF proposal for industrial IoT

Problem statement:

- Interplay of MAC and Routing protocols affects network performance
- Initial network formation is challenging

Our simulative study highlights...

- why a blind adoption of IETF 6TiSCH proposal is risky
- how to tune the MAC and Routing protocol for a successful network formation (situation dependent)

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How a 6TiSCH network is formed (1)

At least two processes, before network is operational:

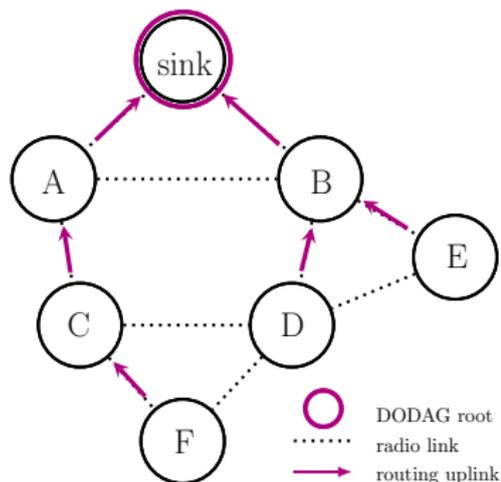
1 TSCH synchronisation

- Goal: build a globally synchronized mesh network

Exchange of Enhanced Beacon
 Formation of time information

2 RPL DODAG construction

Goal: build a mesh network with a
 DODAG (Destination Oriented
 Directed Acyclic Graph) structure
 Exchange of RPL messages
 Construction of the DODAG



How can we coordinate these two interplaying processes?

How a 6TiSCH network is formed (1)

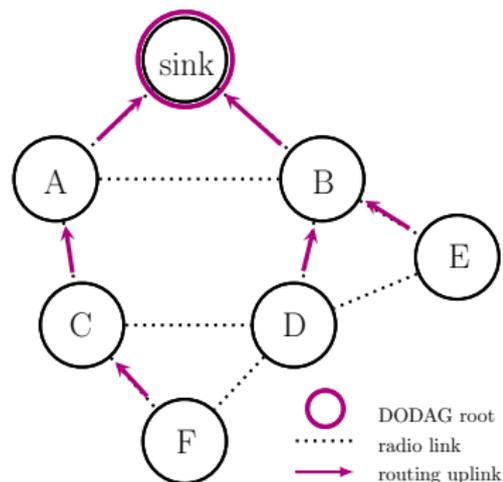
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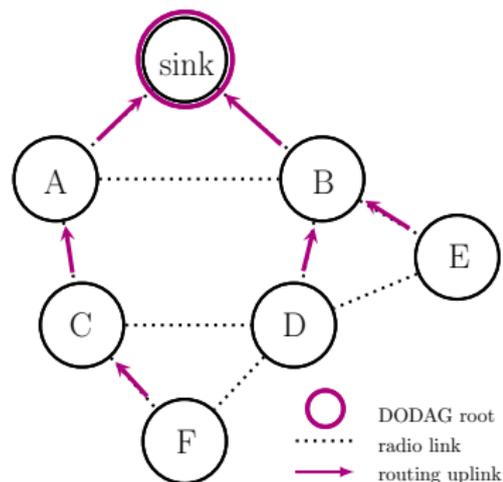
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- Goal: build a globally synchronized mesh network
- Exchanging Enhanced Beacon (EB) frames with time information
 - emitted every t_{EB}

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- Goal: organize nodes as a directed tree rooted at the sink
- RPL uses hop-by-hop link quality estimations



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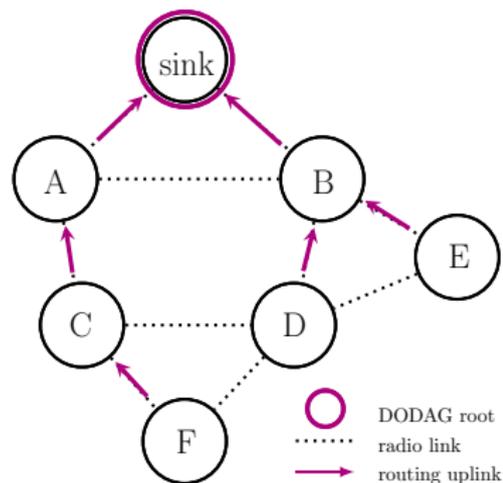
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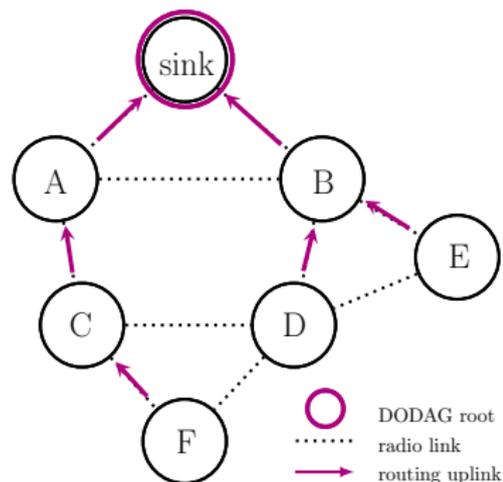
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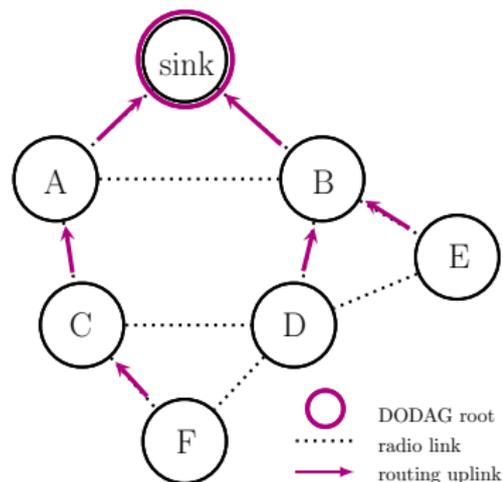
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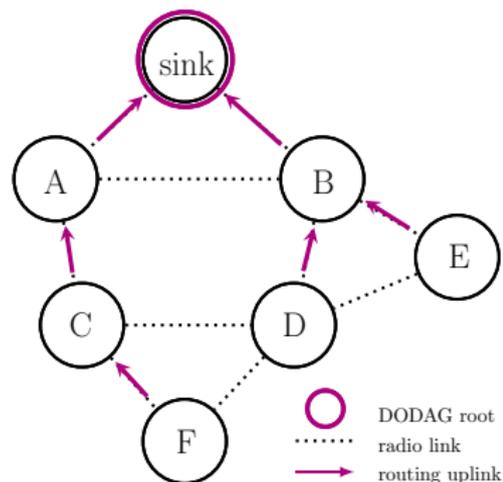


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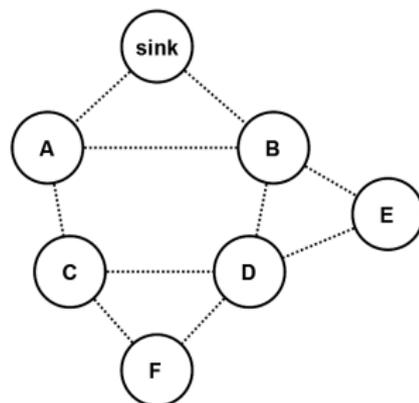


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- 1 Sink sets TSCH-schedule with one shared slot and sends EBs and DIOs
- 2 Joining Nodes keep their radio on and listen for EB
- 3 After hearing an EB: Node learns the *minimal schedule* and is synchronised
- 4 After hearing a DIO: Node selects a preferred parent and broadcasts EBs and DIOs messages on its turn.



Legend:
 radio link

At the end: every node knows the *minimal schedule* and is in the DODAG

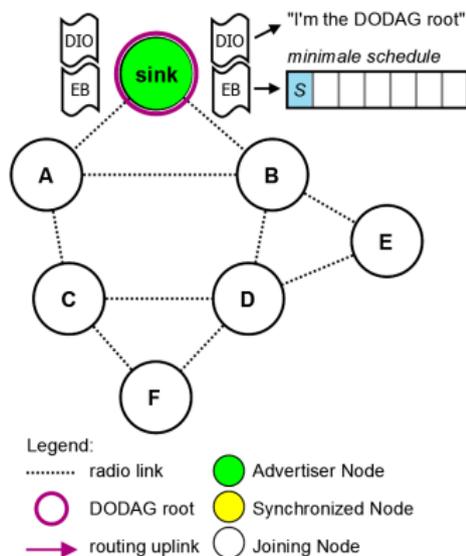
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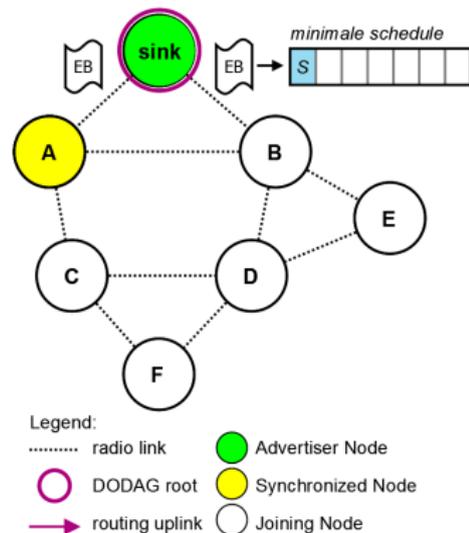


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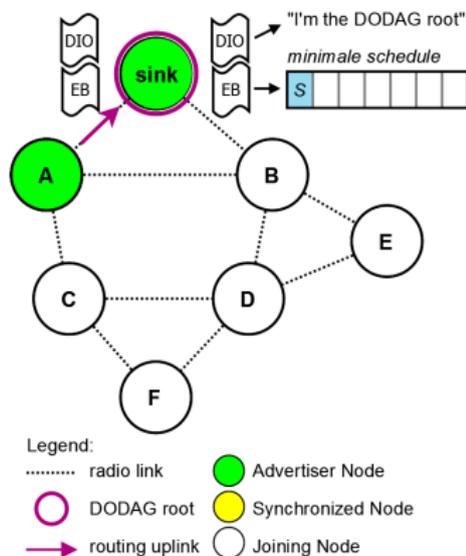
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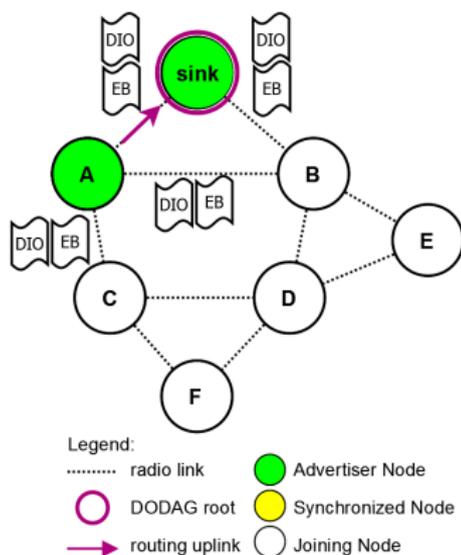
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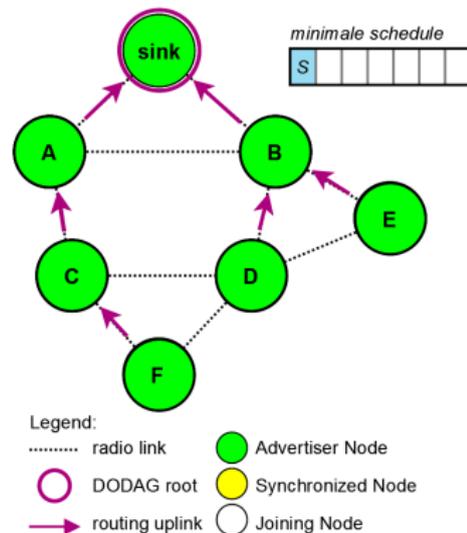
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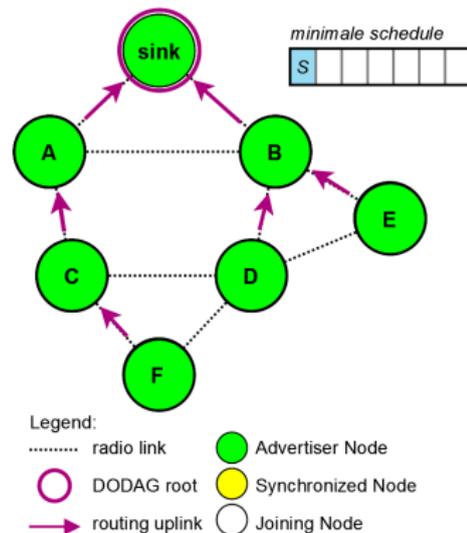
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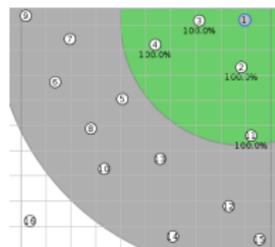
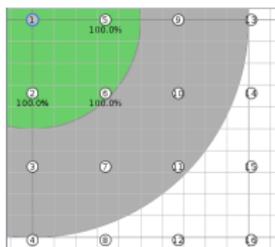
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Simulation Setup

- Contiki OS and Cooja simulator
 - (1) open-source, (2) popular and (3) compliance with 6TiSCH-stack
- Different topologies and three network sizes $N_{size} \in \{9, 16, 25\}$

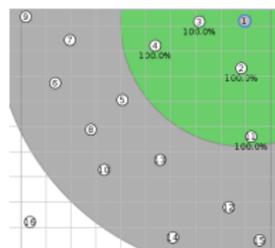
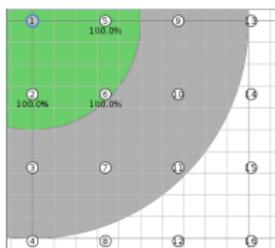
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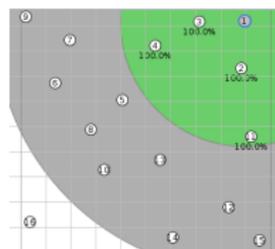
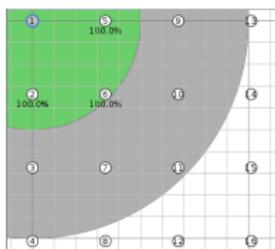


- Varying crucial parameter of TSCH and RPL Trickle:

Parameter	Symbol	Value
TSCH number of channels	N_c	$\{4, 16\}$
TSCH EB period	t_{eb}	$\{2048, 4096, 8192, 16384\}$ ms
RPL minimal interval	l_{min}	$\{128, 256, \dots, 4096\}$ ms

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- Performance metrics:
 - time, (2) charge consumed and (3) number of control frames exchanged until completed network formation

Results: Limits of 6TiSCH-MC

In dense network or with improper setting of TSCH and RPL parameters:

- 1 some nodes are not yet operational after 30 minutes
- 2 high battery consumption in several nodes

Table: Successful DODAG formations within 30 min

		Grid			Ellipse			Random		
		N_{size}			N_{size}			N_{size}		
t_{eb}	N_c	9	16	25	9	16	25	9	16	25
2048 ms	4	0%	0%	0%	70%	16%	0%	0%	0%	0%
4096 ms	4	84%	0%	0%	100%	100%	86%	46%	0%	0%
8192 ms	4	100%	20%	0%	100%	100%	100%	100%	2%	0%
16384 ms	4	100%	70%	0%	100%	100%	100%	100%	14%	0%

With slotframe duration $T_{sf} = 1.01 s$ (i.e. $N_s = 101$, $t_s = 10 ms$) and $l_{min} = 1024 ms$

Due to...

- collisions of control frame
- queuing delay of DIO packets

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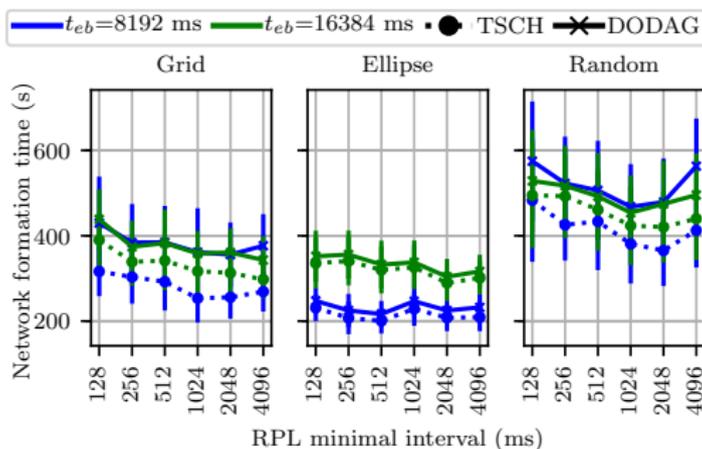
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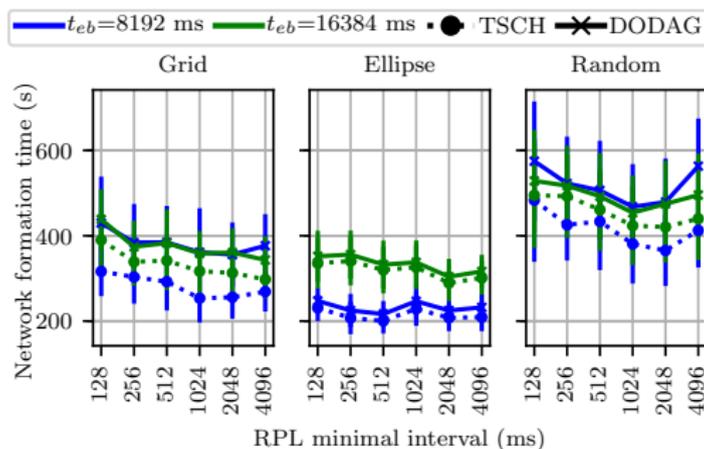
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Recommendations for implementers

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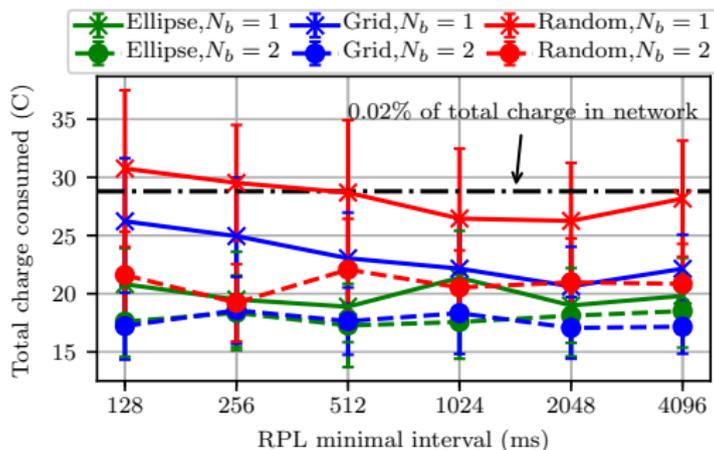
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Main Observations and Recommendations (2)

Extending 6TiSCH-MC with $N_b = 2$ shared slots:



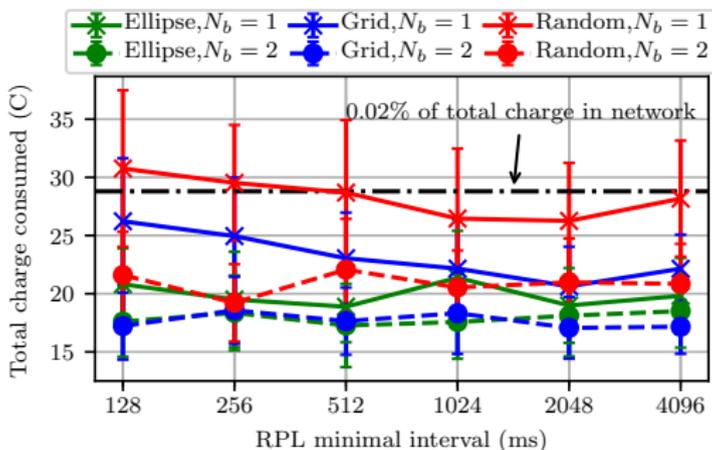
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Recommendation for implementers:

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Herein:

- Overview of the IETF 6TiSCH *minimal configuration* (6TiSCH-MC)
- Extensive simulations to characterize its behaviour

Conclusions:

- Potential downsides of 6TiSCH-MC with dense topologies
- Recommendations for setting TSCH and RPL parameters

Future work:

- Validate the results with testbeds/realistic channel
- Develop an algorithm for allocation of broadcast links in TSCH

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