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Reactive Load Control based on Binary Packet Marking

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Overview

Motivation & Introduction

System & Prototype Design

Experimental Evaluation

Summary





Motivation & Introduction

Internet QoS

- highly debated, many proposals, little deployment
- traditional tools: admission control, traffic regulation, packet scheduling

Admission Control & Traffic Regulation

- service guarantees through overload protection

Differentiated Scheduling

- needed for delay differentiation

⇒ **PROACTIVE RESOURCE ALLOCATION MODEL**

Existing Internet Reactive Resource Allocation Model

- flow control
- adaptivity ⇒ elastic utility functions

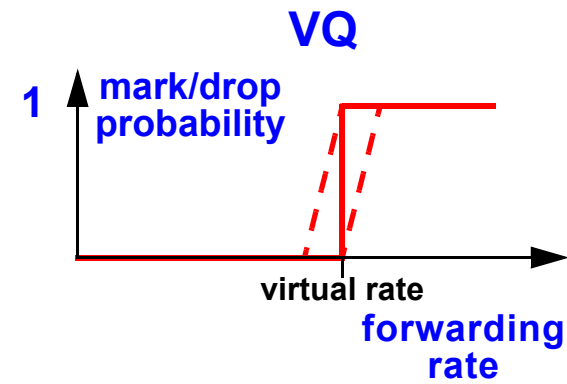
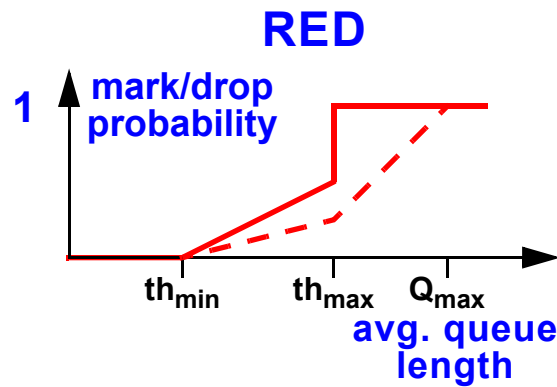
BUT: inelastic applications?

⇒ **GOAL: REACTIVE RESOURCE ALLOCATION FOR INELASTIC FLOWS**





Background – Packet Marking Algorithms



Random Early Detection (RED) & Variants

- **queue-based feedback**

Virtual Queue (VQ) & Variants

- **essentially rate- and queue-based feedback**
- **inelastic flows** → binary path marking rate
- **simplification: threshold-based marking (TBM)**

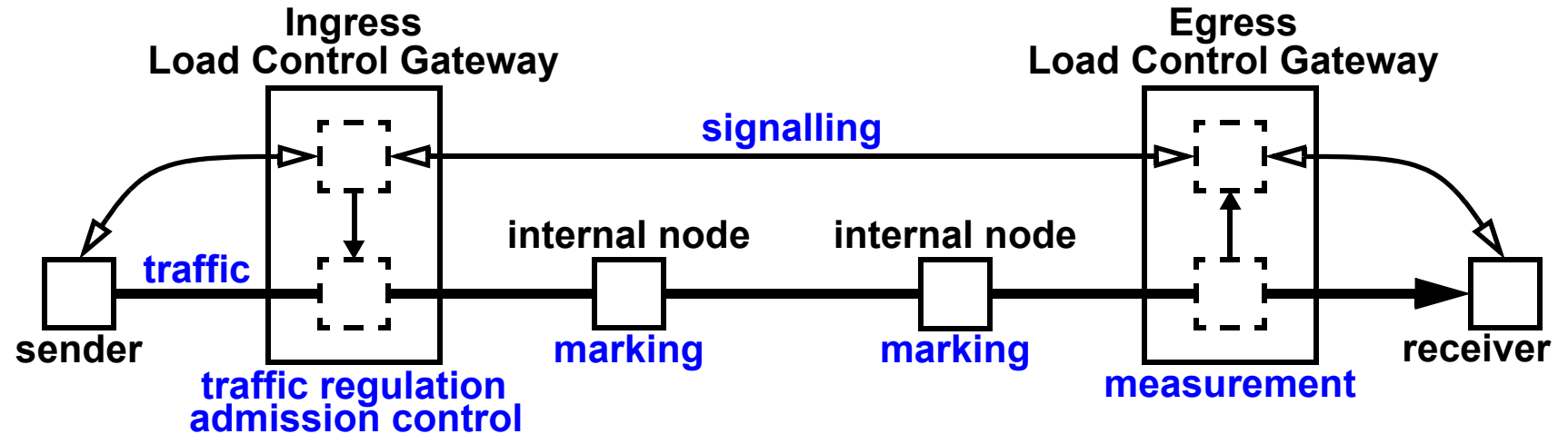




Reactive Load Control System Design

Reactive Resource Allocation: Admission Control for Inelastic Flows

- admission control based on network load feedback



- inherent feedback delay between egress and ingress

Implementation

- packet discrimination & marking → ECN (ECT & CE bit)
- signalling & admission control → RSVP (+ extensions)

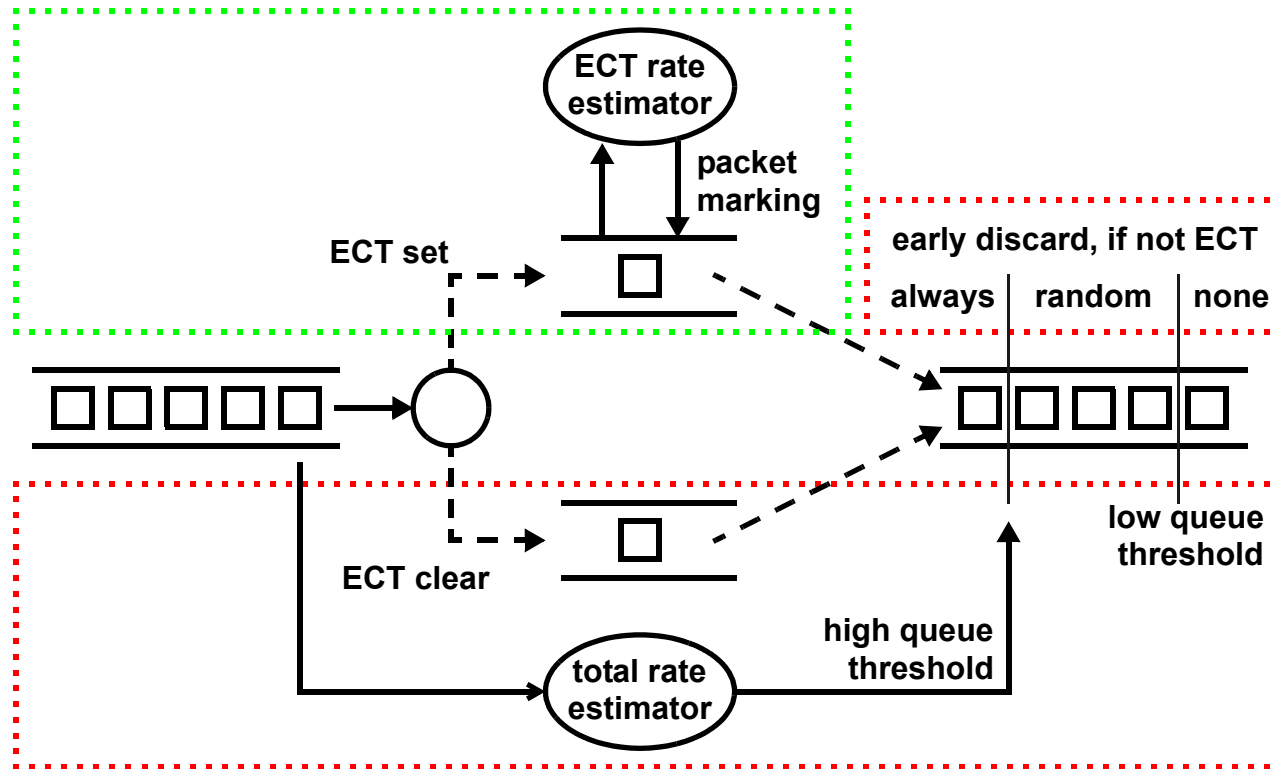




Internal Nodes

Dropping of Non-markable Packets?

⇒ Differentiated Queue Management (DQM) at Internal Nodes



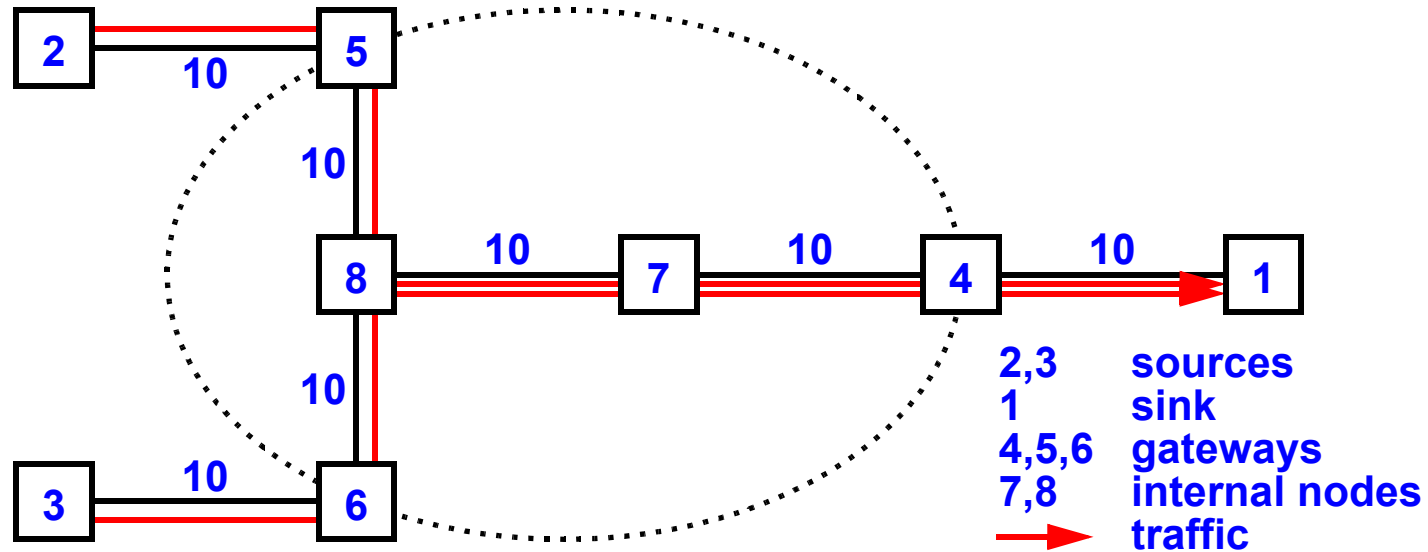
..... packet marking

..... adaptive early random drop (ERD)





Lab Experiments



Traffic Generation

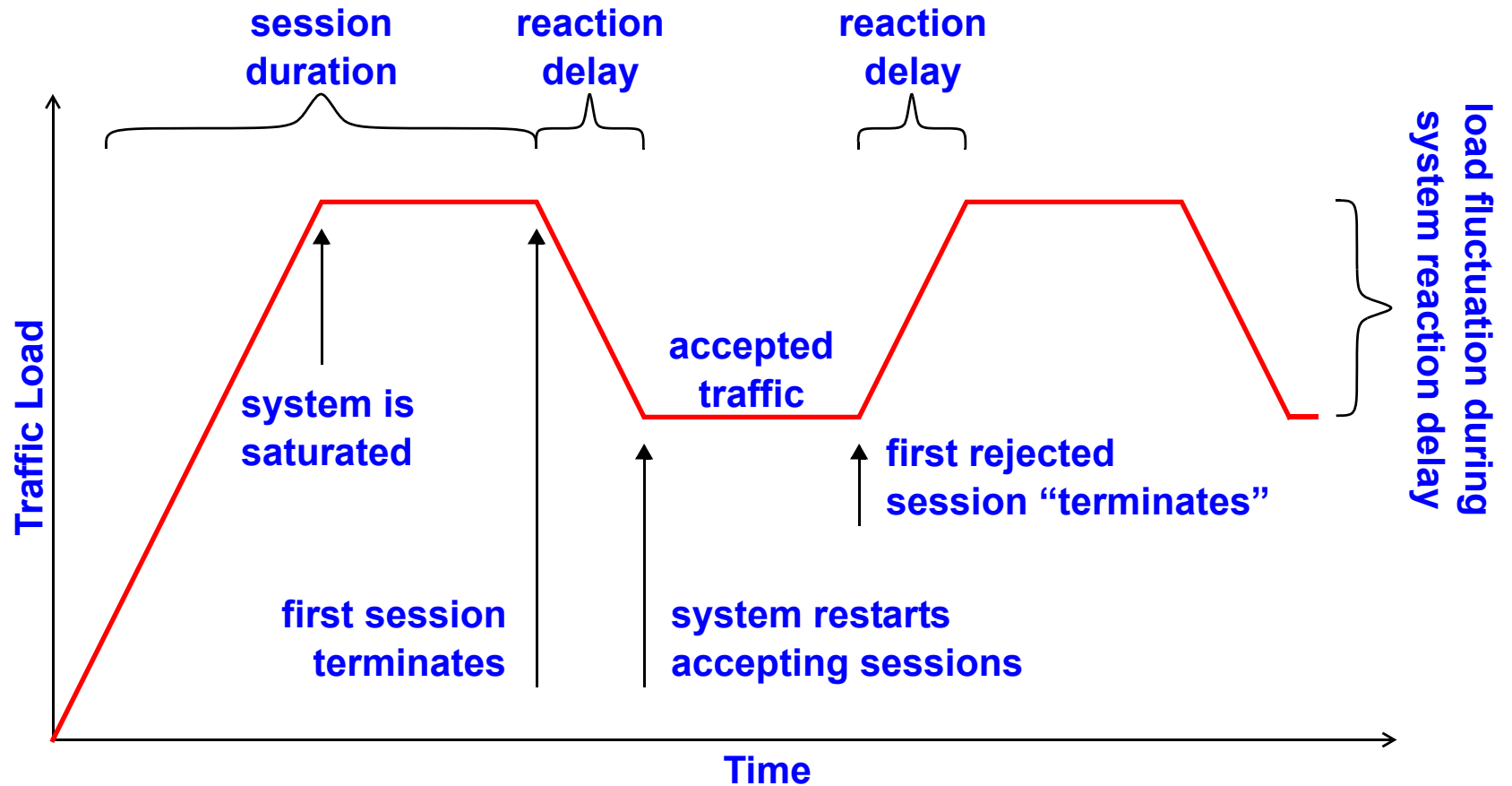
- **signalled traffic flows**
 - CBR traffic flows (VoIP-like, 64 Kbit/s IP payload)
 - strictly periodic or Poisson arrival
 - fixed or exponential duration
 - (no) persistent session setup (reject → retry after short random period)
- **(no) background traffic**
 - aggregate of Pareto-distributed on-off sources → self-similar traffic





System Behaviour – Model

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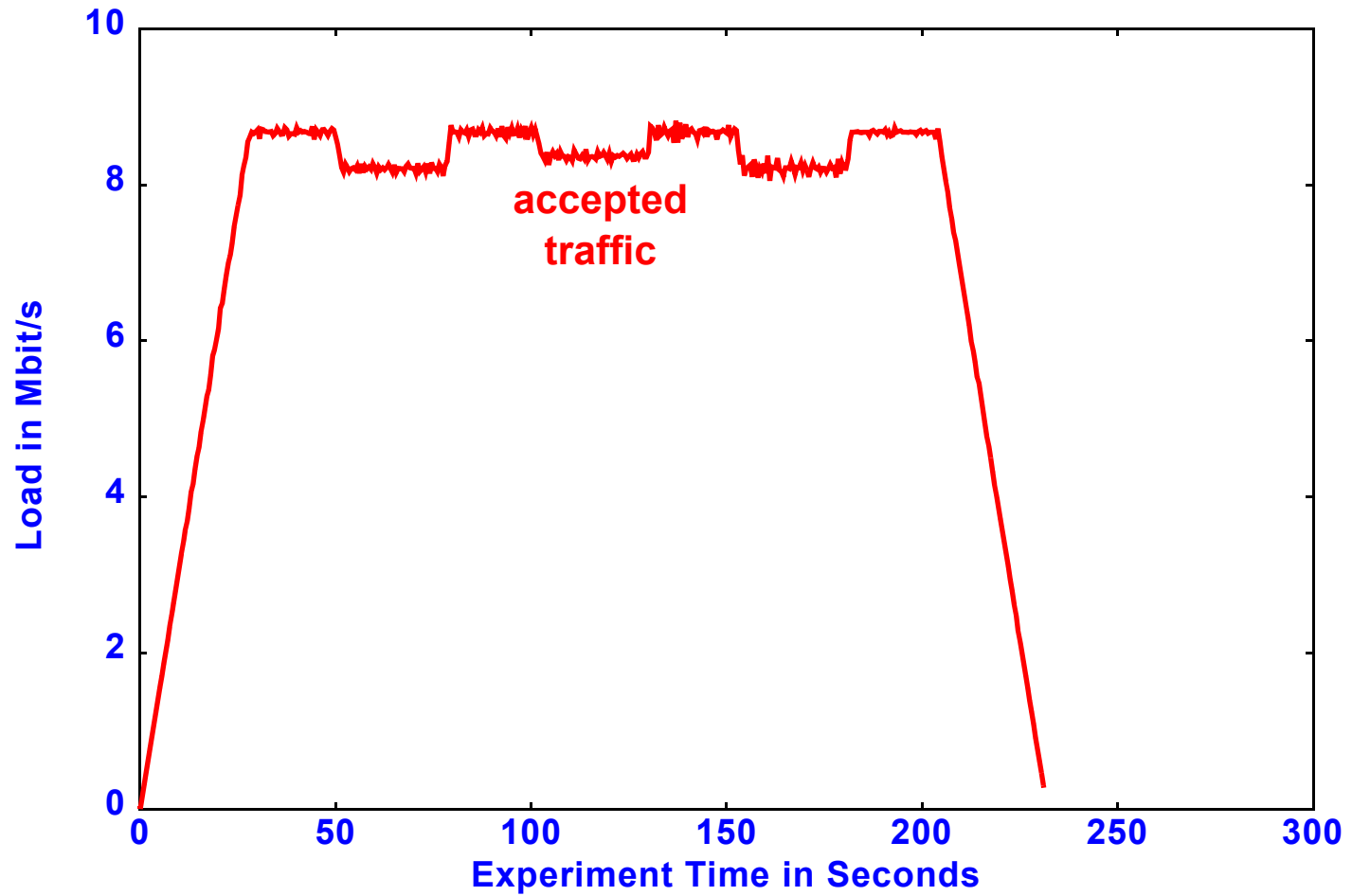
- deterministic flow arrival and duration
- flows "stay" in the system if rejected (but don't retry signalling)
- cumulative flows exceed available resources





System Behaviour – VQ

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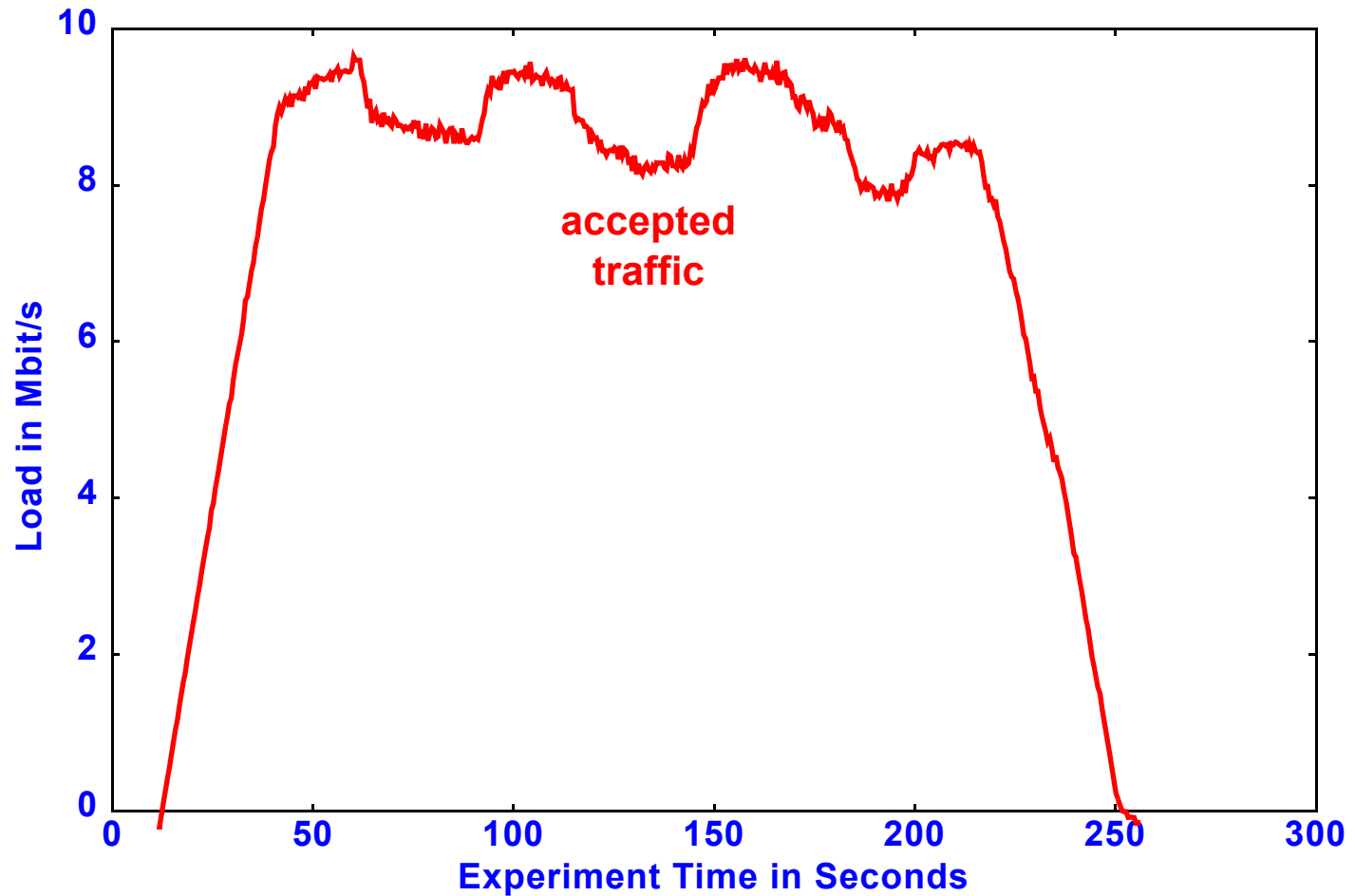
- system reaction delay: approx. 2 seconds
- TBM slightly faster





System Behaviour – RED

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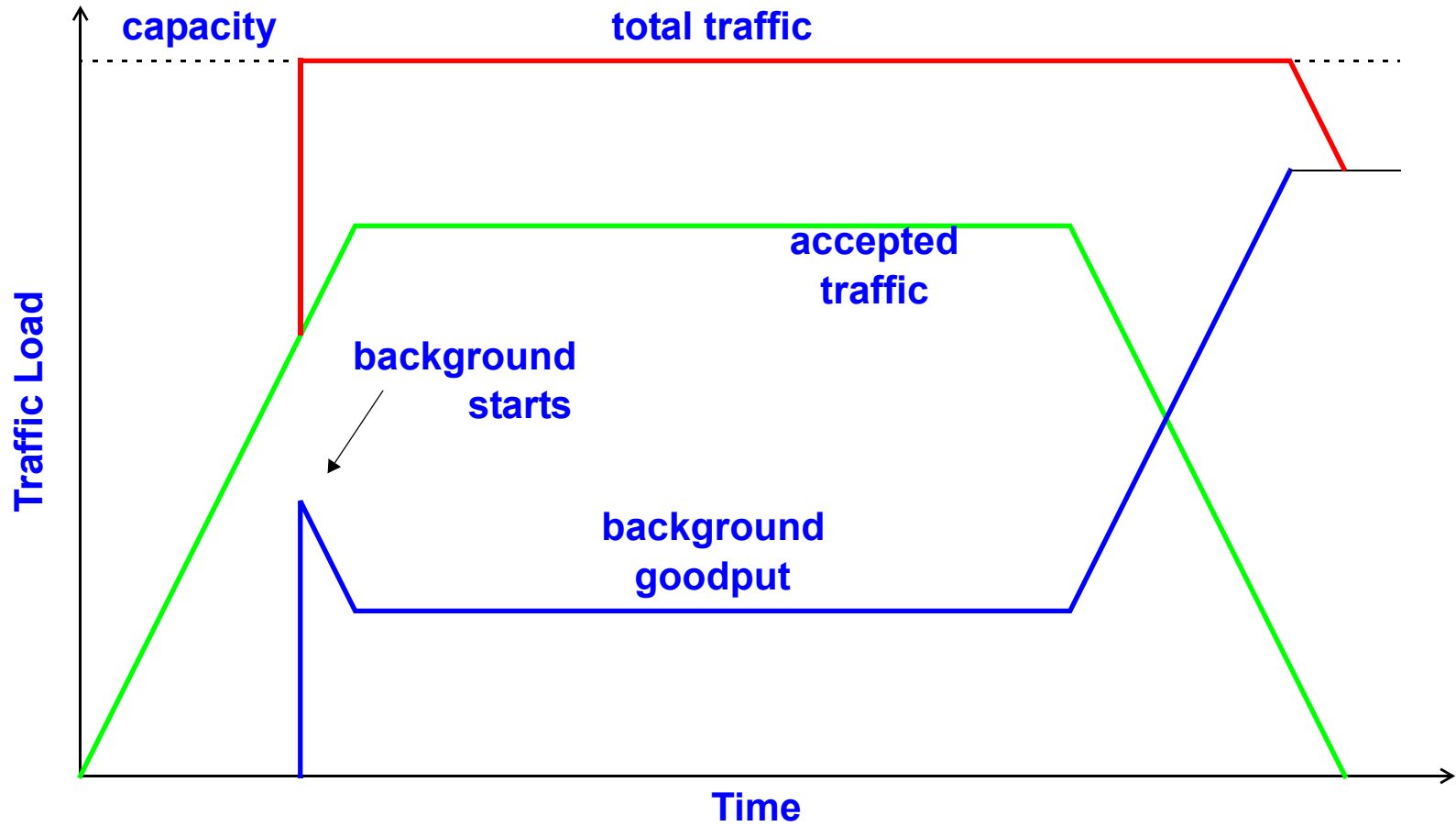


- system reaction delay: approx. 4 seconds
- aggressive RED parameters needed ...





Traffic Discrimination – Model



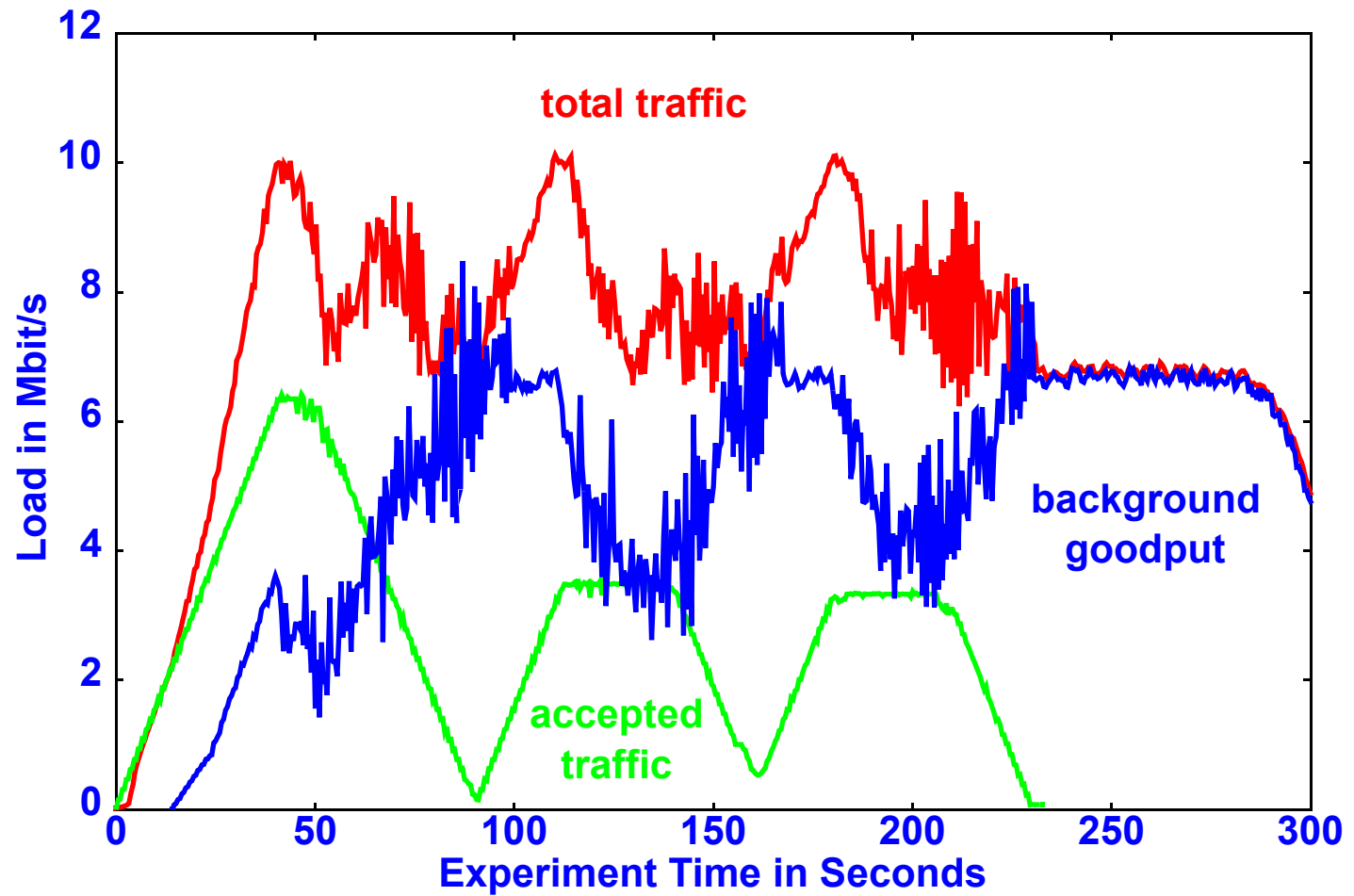
- deterministic flow arrival and duration not exceeding resources
- aggregate Pareto background traffic exceeding resources
- compare RED, VQ, and TBM





Traffic Discrimination – RED & TCP Background

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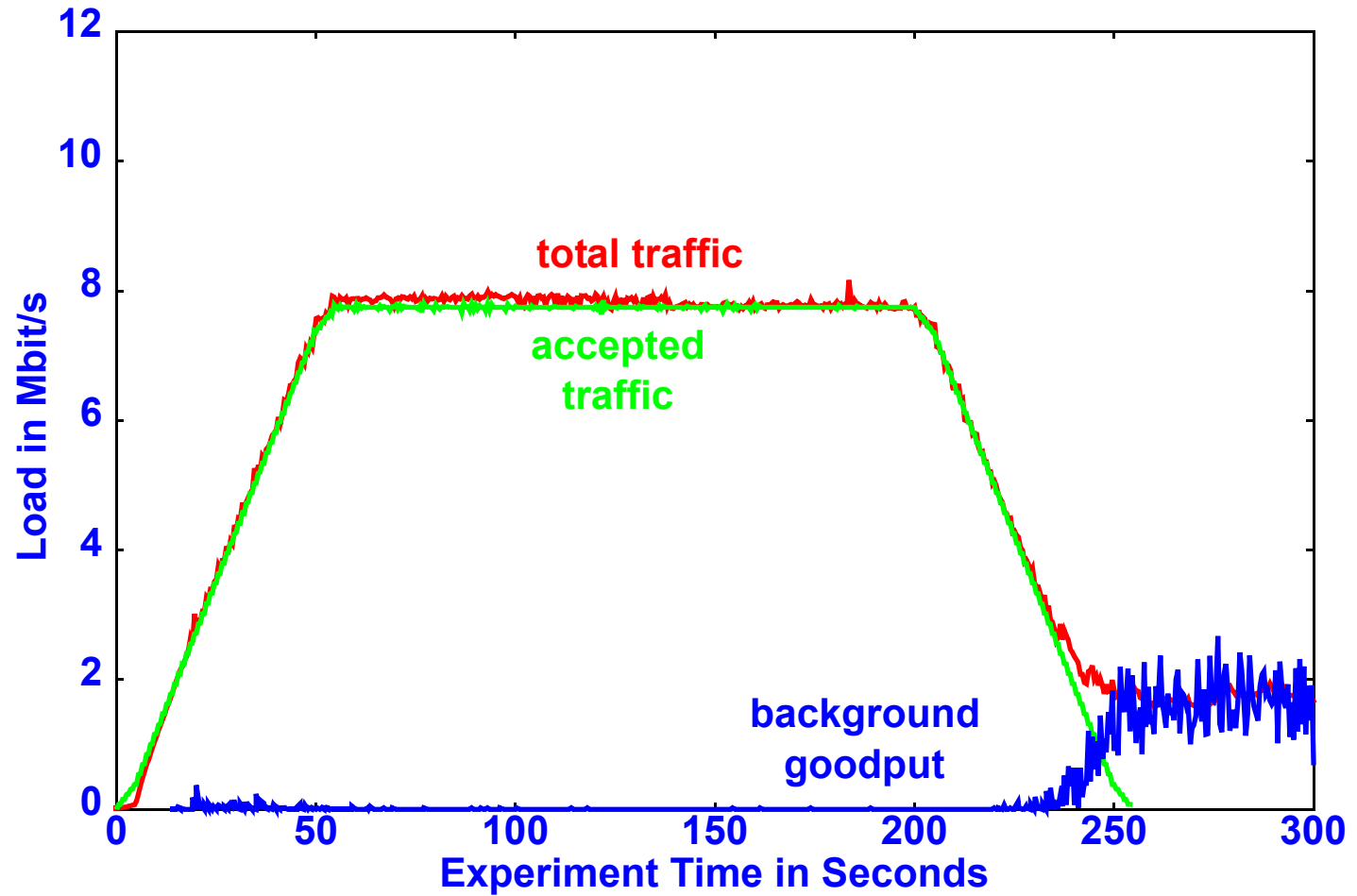
- total chaos ...





Traffic Discrimination – VQ & TCP Background

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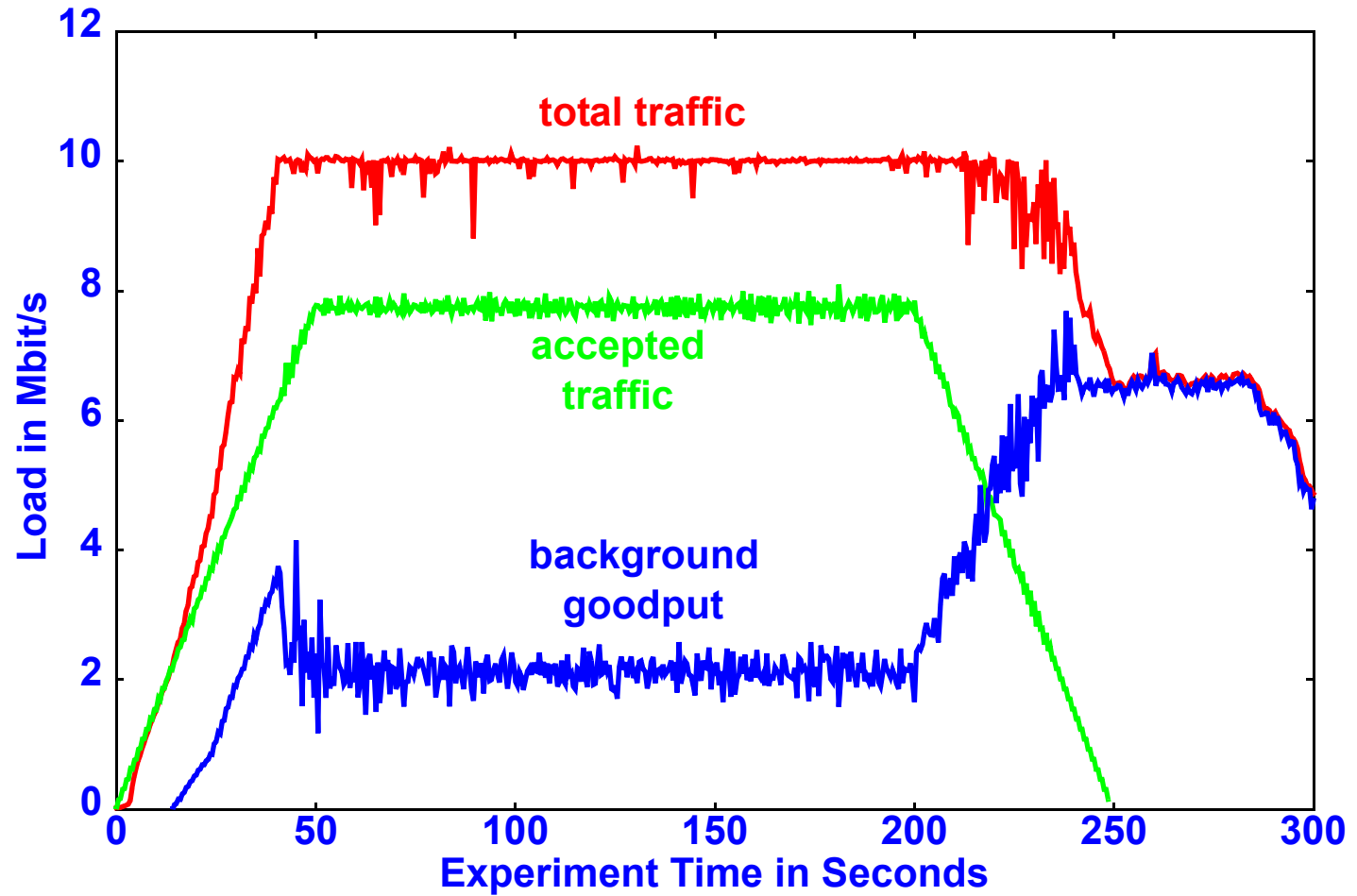
- aggressive bursty packet drop
- similar for TBM & TCP





Traffic Discrimination – VQ/DQM & TCP Background

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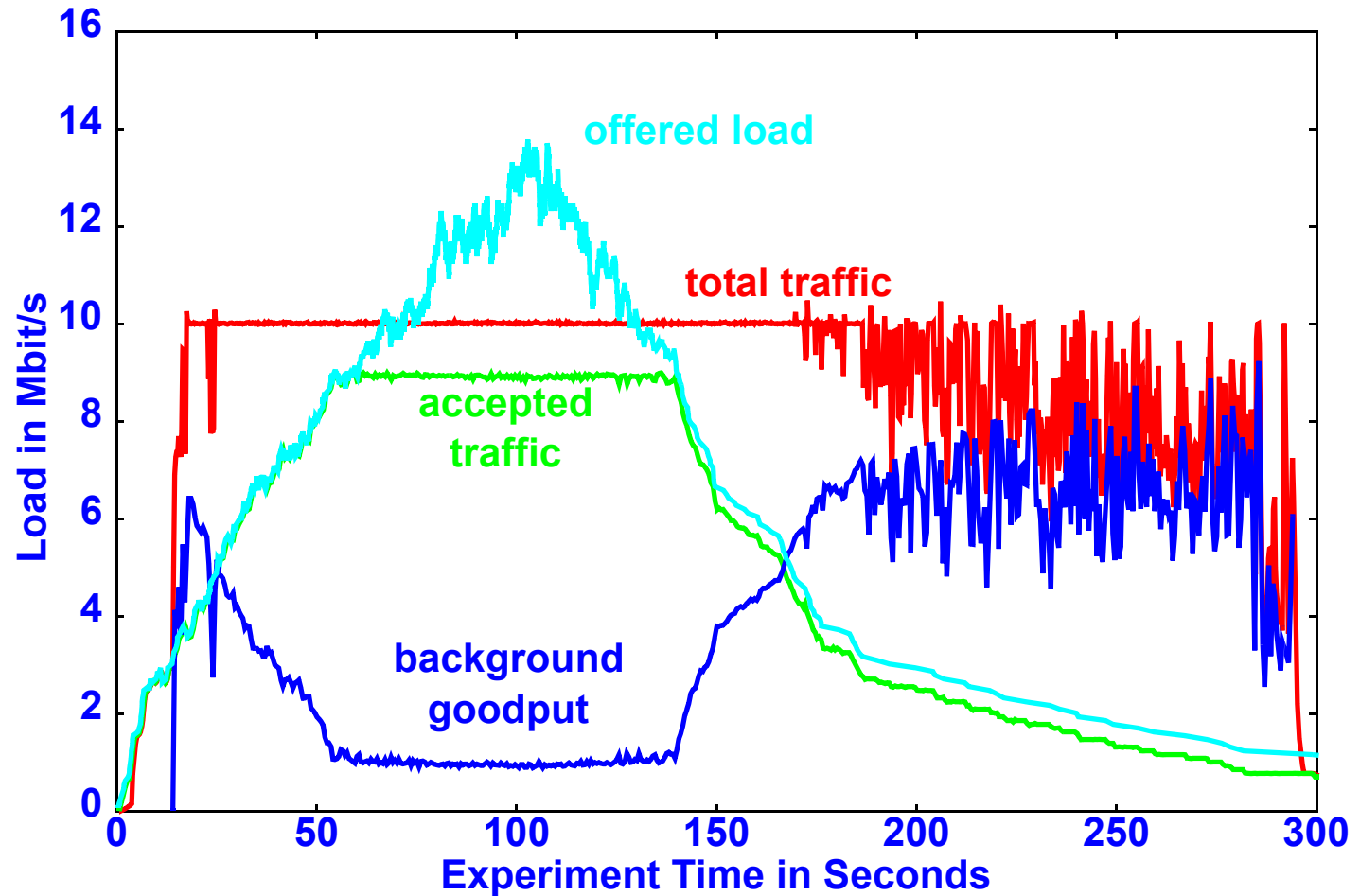


- proper traffic discrimination through DQM
- background TCP regulated by ERD





System Performance – Benchmark: IntServ



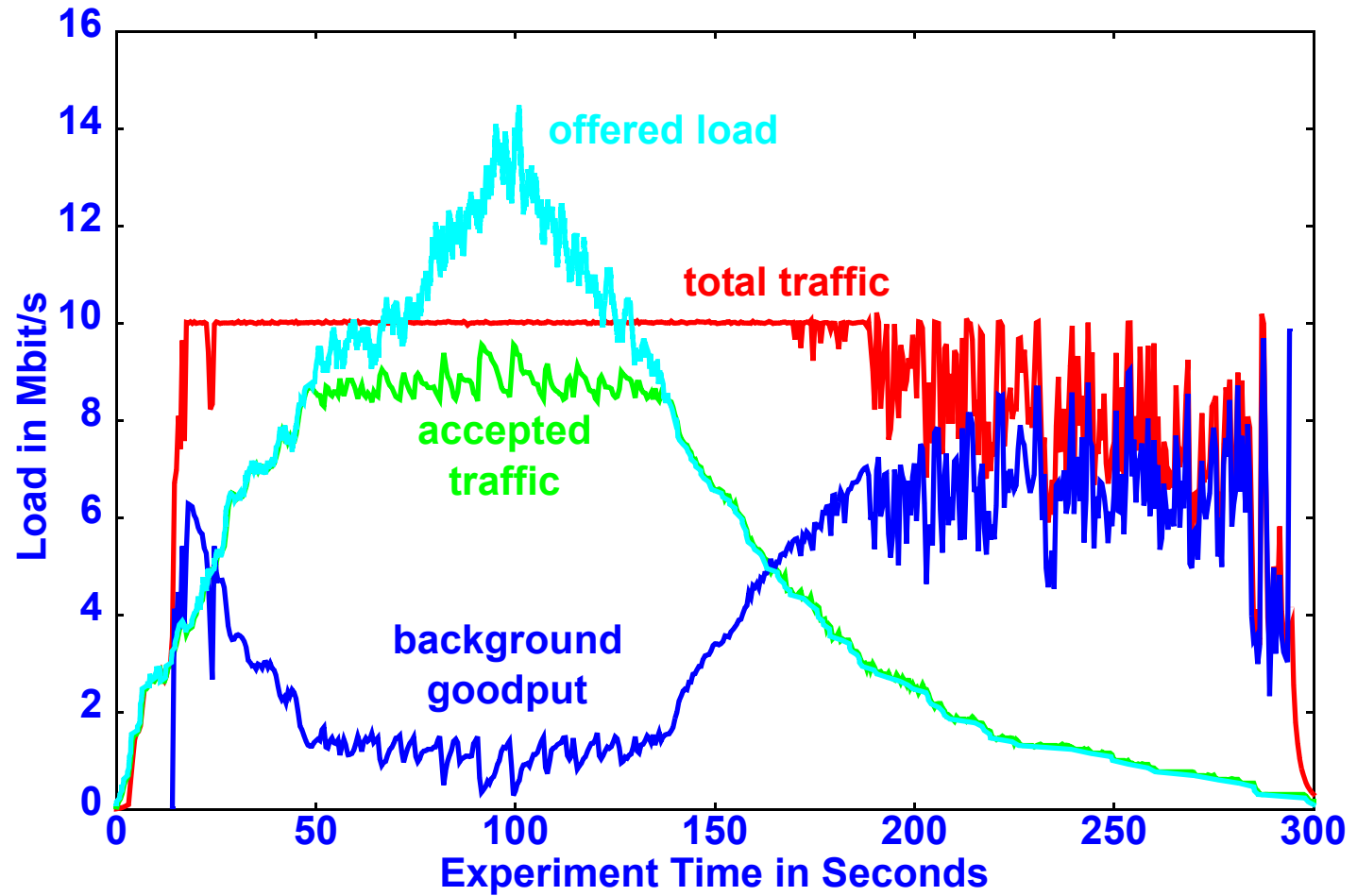
- overload from foreground traffic (exponential inter-arrival & duration)
- persistent session setup (retry of short random timeout)





System Performance – TBM/DQM

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- system reaction delay
- fast increase because of persistent sessions
- slightly better than VQ/DQM





Summary

Design & Implementation

- reactive load control design
- feedback signalling & admission control → RSVP extensions
- various packet marking algorithms
- prototype implementation (open source): lab and simulation

Experimental Evaluation

- verification of system behaviour
- system reaction delay
- suitability of different marking algorithms in different scenarios
- benefits of DQM
- overall performance compared to IntServ configuration

