## **Griffin: Towards an Agile, Predictive Infrastructure**

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# Outline

- Griffin
  - Motivation
  - Goals
  - Architecture
- Tapas
  - Motivation
  - Multi-layer protocol trace collection
  - Data preconditioning-based network modeling
  - Network feedback and prediction

# Near-Continuous, Highly-Variable Internet Connectivity

- Connectivity everywhere: campus, in-building, satellite...
  - Projects: Sahara (01-), Iceberg (98-01), Rover (95-97)
- Most applications support limited variability (1% to 2x)
  - Design environment for legacy apps is static desktop LAN
  - Strong abstraction boundaries (APIs) hide the # of RPCs
- But, today's apps see a wider range of variability
  - 3 $\rightarrow$ 5 orders of magnitude of bandwidth from 10's Kb/s  $\rightarrow$ 1 Gb/s
  - 4 $\rightarrow$ 6 orders of magnitude of latency from 1  $\mu$ sec  $\rightarrow$ 1,000's ms
  - 5→9 orders of magnitude of loss rates from  $10^{-3} \rightarrow 10^{-12}$  BER
  - Neither best-effort or unbounded retransmission may be ideal
  - Also, overloaded servers / limited resources on mobile devices
- Result: Poor/variable performance from legacy apps

## **Griffin Goals**

- Users always see excellent (= local, lightly loaded) application behavior and performance
  - Independent of the current infrastructure conditions
  - Move away from "reactive to change" model
  - Agility: key metric is time to react and adapt
- Help legacy applications handle changing conditions
  - Analyze, classify, and predict behavior
  - Pre-stage dynamic/static code/data (activate on demand)
- Architecture for developing new applications
  - Input/control mechanisms for new applications
  - Application developer tools
- Leverage Sahara policies and control mechanisms

# **Griffin: An Adaptive, Predictive Approach**

- Continuous, cross-layer, multi-timescale introspection
  - Collect & cluster link, network, and application protocol events
  - Broader-scale: Correlate AND communicate short-/long-term events and effects at multiple levels (breaks abstractions)
  - Challenge: Building accurate models of correlated events
- Convey app reqs/network info to/from lower-levels
  - Break abstraction boundaries in a controlled way
  - Challenge: Extensible interfaces to avoid existing least common denominator problems
- Overlay more powerful network model on top of IP
  - Avoid standardization delays/inertia
  - Enables dynamic service placement
  - Challenge: Efficient interoperation with IP routing policies

# Some Enabling Infrastructure Components

- Tapas network characteristics toolkit
  - Measuring/modeling/emulating/predicting delay, loss, ...
  - Provides Sahara with micro-scale network weather information
  - Mechanism for monitoring/predicting available QoS
- REAP protocol modifying / application building toolkit
  - Introspective mobile code/data support for legacy / new apps
  - Provides Sahara with dynamic placement of data and service sub-components
- Brocade, Mobile Tapestry, and Fault-Tolerant Tapestry
  - Overlay routing layer providing Sahara with efficient application-level object location and routing
  - Mobility support, fault-tolerance, varying delivery semantics

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## **Tapas Motivation**

- Accurate modeling and emulation for protocol design
  - Very difficult to gain access to new or experimental networks
  - Delay, error, congestion in IP, GSM, GPRS, 1xRTT, 802.11a/b
  - Study interactions between protocols at different levels
- Creating models/artificial traces that are statistically indistinguishable from traces from real networks
  - Such models have both predictive and descriptive power
  - Better understanding of network characteristics
  - Can be used to optimize new and existing protocols
- Answer several application design questions
  - Q1: Impact of network layering in application design
  - Q2: Effects of network model choices on results (survey?)
  - Q3: Using feedback for adaptive, predictive applications

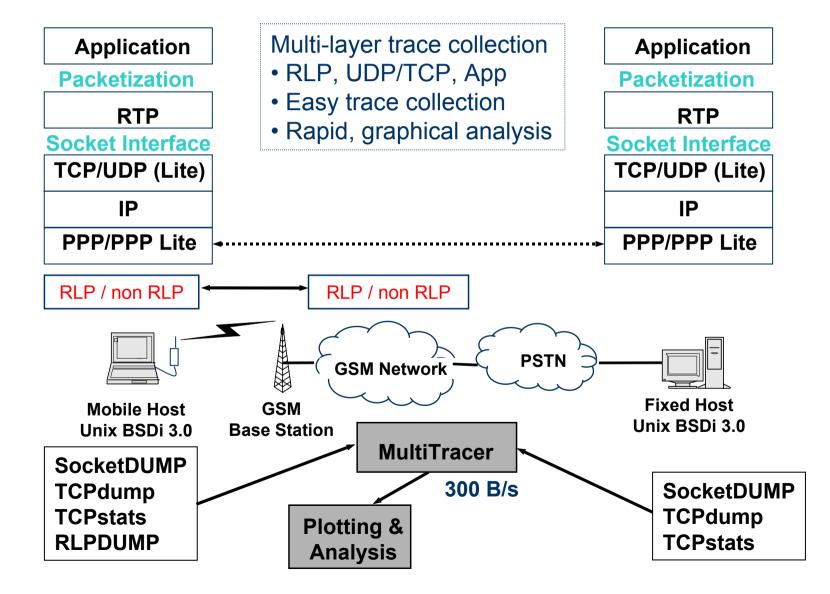
### Tapas

- Novel data preconditioning-based analysis approach
  - More accurately models/emulates long-/short-term dependence effects than classic approaches (Gilbert, Markov, HMM, Bernoulli)
- Analysis, simulation, modeling, prediction tools:
  - MultiTracer: Multi-layer trace collection and analysis (*download*)
  - Trace analysis and synthetic trace generator tools
    - Markov-based Trace Analysis, Modified hidden Markov Model
  - WSim: Wireless link simulator (currently trace-driven)
  - Simple feedback algorithm and API
  - Domain analysis tool: chooses most accurate model for a metric
- Error-tolerant radio / link layer protocols: RLPLite, PPPLite
- Collected >5,000 minutes of TCP, UDP, RLP traces in good/bad, stationary/mobile environments (*download*)

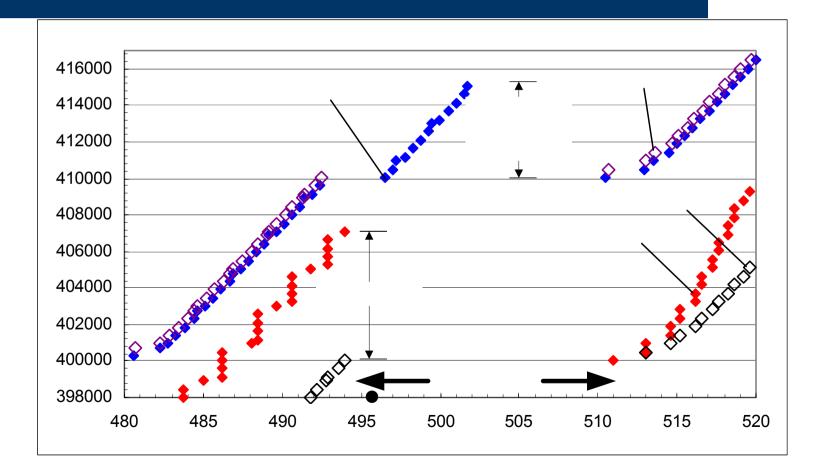
# **Optimizing GSM CS Data Protocol**

- Circuit-switched data Radio Link Protocol (RLP)
  - Semi-reliable ARQ Protocol at 9.6 Kbit/s
  - Link resets after N = 6 number of retransmissions
- RLP and TCP interaction measurement / analysis
  - Both are reliable protocols (link and transport layers)
  - Researchers claim *competing retransmissions* problem (Q1)
- Fixed logical frame size of 30 bytes
  - 6 bytes header/checksum
  - What are effects of alternative frame sizes on throughput?
  - What effect does the choice of model have on results? (Q2)
- MultiTracer trace collection/analysis tool quickly identifies *interaction* and *performance* effects

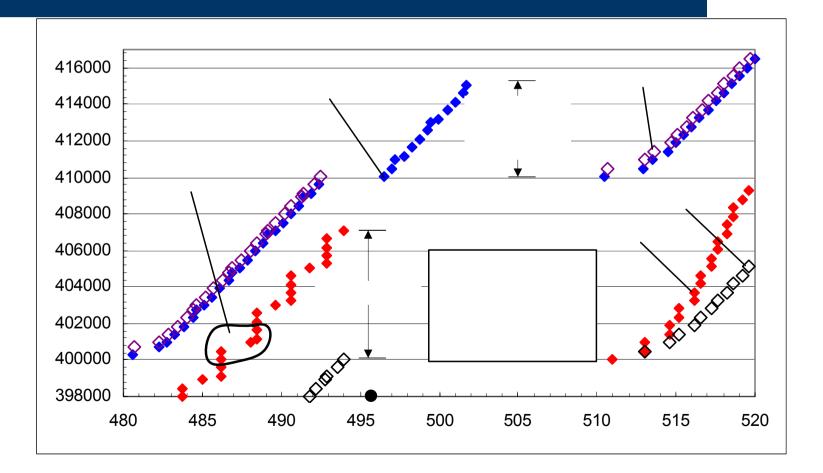
## **MultiTracer Measurement Testbed**



### **Time-Sequence TCP Plot**



### MultiTracer TCP and RLP Plot



# **Choosing the Right Network Model**

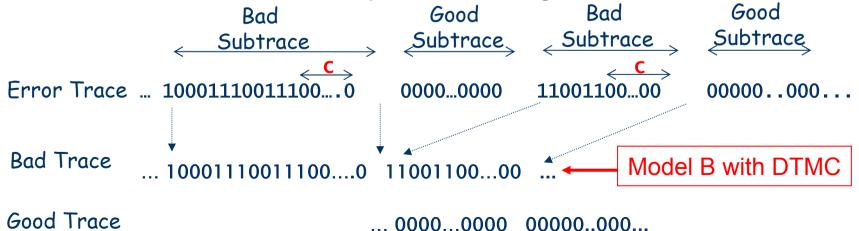
- Collect empirical packet trace: T = {1,0}\*
  - 1: corrupted/delayed packet, 0: correct/non-delayed packet
- Create mathematical models based on T



- Challenge: domain analysis which model to use?
- T may be *non-stationary* (statistics vary over time)
  - Classic models don't always work well (can't capture variations)
- MTA, M<sup>3</sup> Trace data preconditioning algorithms
  - Decompose T into stationary sub-traces & model transitions
  - Stationary sub-traces can be modeled with high-order DTMC
  - Markov-based Trace Analysis (MTA) and Modified hidden Markov
  - Model (M<sup>3</sup>) tools accurately model time varying links

# **Creating Stationarity in Traces**

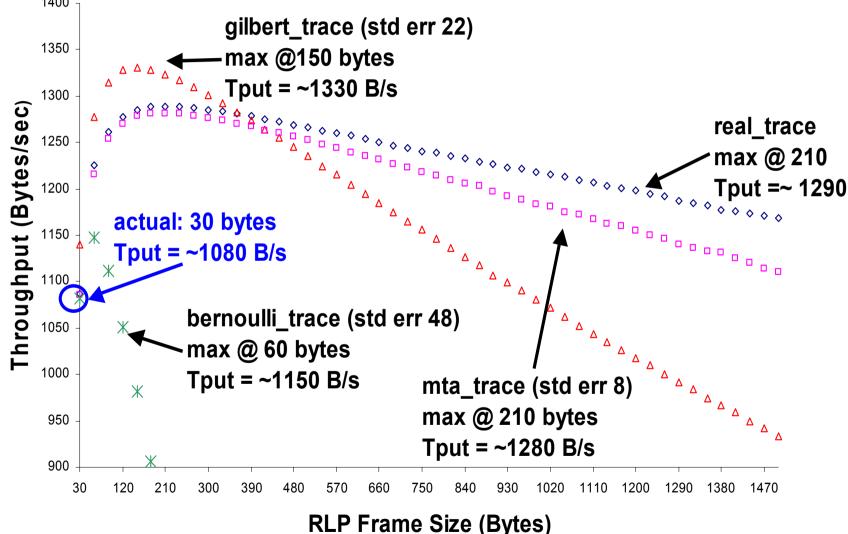
- Our idea for MTA and M<sup>3</sup>: decompose T into stationary sub-traces
  - Bad sub-traces  $B_{1..n} = 1\{1,0\}^*0^c$ , Good sub-traces  $G_{1..n} = 0^*$
  - C is a change-of-state constant: mean + std dev of length of 1\*
- MTA: Model B with a DTMC, model state lengths with exponential distribution, and compute transitions between states
- M<sup>3</sup>: Similar, but models multiple states using HMM to transition



# **Role of Accurate Network Modeling**

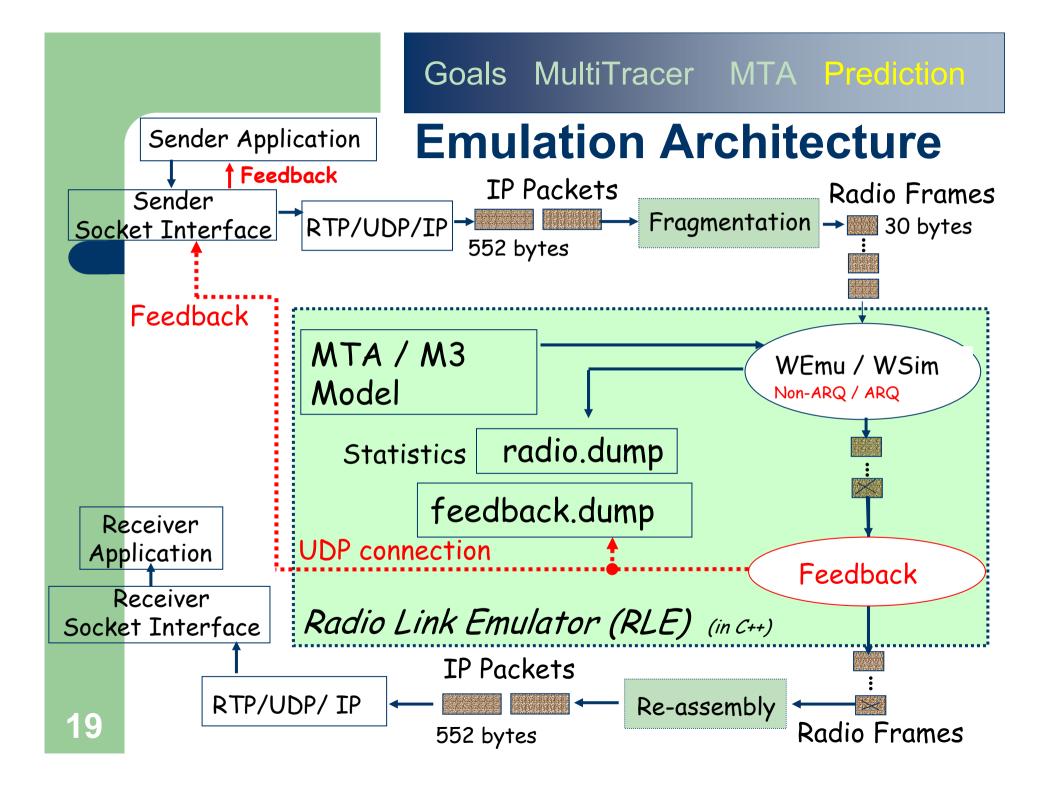
- Choosing optimal RLP frame size for bulk data transfer
  - Use real Block Error Traces from MultiTracer
  - Compare against synthetic traces with same BER generated by different models (MTA, Bernoulli, Gilbert)
  - Measure throughput for different frame sizes (30,60,...,1500 B)
- Optimal frame size is the one with highest throughput
  - Reduced overhead vs. increased retransmission delay
- MultiTracer analysis: errors occur in long bursts
  - Channel is either OK or very bad
  - A few long or many short packets are affected
- Burst effect not captured by classical models

## **Right Model** ⇒ **More Accurate Results**



# **Predictive, Error-Resilient Video**

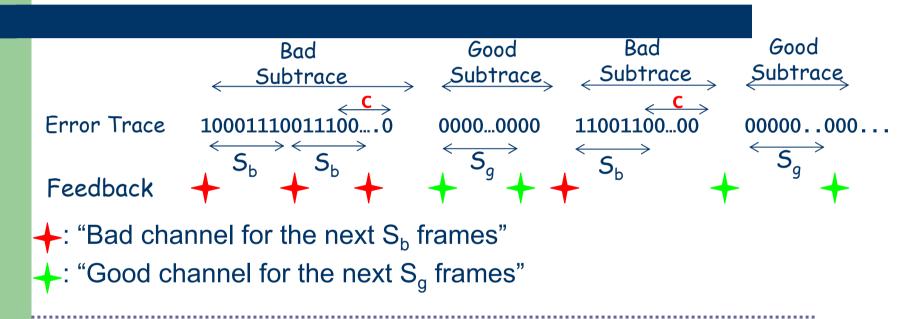
- Explore benefits of feedback and prediction
  - Dynamic adaptation of data rate to network conditions for adaptive, error resilient video codecs
    - H.263+ error-resilient QCIF codec (< 64 kb/s, 10-15fps)
- Link layer reliability helps with wireless errors
  - But, link layer reliability alone introduces delay  $\Rightarrow$  higher jitter
- Add a simple feedback algorithm
  - High and low resolution versions of video stored on server
  - Channel prediction bad state  $\Rightarrow$  send low resolution frame
  - Goal: switch data rates to minimize jitter effects by keeping interframe arrival times relatively flat



# **Streaming Video with Predictive Feedback**

- Simulation using wireless traces (4480 min):
  - ~4 min per video, bad channel (-105 to -99dB), ~ 1.5 % BLER
- Calculate S<sub>b</sub> and S<sub>g</sub> (avg bad/good subtrace lengths), and C
- Receiver receives radio frame:
  - In "bad subtrace"?  $\Rightarrow$  send bad state feedback every S<sub>b</sub> frames
  - In "good subtrace "?  $\Rightarrow$  send good state feedback every  $S_g$  frames

### **Predictive Feedback Algorithm**



• Preliminary results

- No feedback: jitter std dev 150 ms, many 200+ ms instances
  Feedback: jitter std dev 100 ms, only two 200+ ms instances
- Next steps: More sophisticated algorithms, full platform, RTCP feedback-based mechanism

### **Tapas Summary**

- A better understanding of effects of multi-layer effects
- Accurate models ⇒ Accurate simulation ⇒ Better protocol design
- 1<sup>st</sup> cut simple socket interface model for communicating with lower protocol stack layers
- Preliminary result: Prediction enables better response time to discontinuous changes in error rate
- On-going work:
  - Trace collection: CDMA 1xRTT, GPRS, & IEEE 802.11a

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