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## Computational Geometry Homework Set 3, 02.12.2014

Solutions are due Tuesday, December 9, 2014, until 9:45 in the mailbox for homework sheets or at the beginning of the lecture. Please put your name on all pages!


Exercise 1 (Convex hulls):
Let $W$ be a simple closed path that encloses the convex hull of a point set $S$. Show that the boundary of $\operatorname{ch}(S)$ is at most as long as $W$.

Exercise 2 (Triangulations and convex hull):
Show that each triangulation of a point set $S$ contains the edges of the points' convex hull.
(5 points)
Exercise 3 (Convex vertices):
Show that any polygon must have at least three convex vertices.

Exercise 4 (Monotonicity):
a) Give an algorithm that tests in $O(n)$ steps whether a simple polygon $P$ is monotone with respect to a line $g$.
You may assume $P$ to be given either as a Doubly-Connected Edge List (DCEL) or simply as a list of vertices and edges.
Hint: You may assume that no edge of the polygon is perpendicular to $g$.
b) Given a simple polygon $P$, give an algorithm that decides in $O(n)$ steps whether there exists a line $g$, such that $P$ is monotone with respect to $g$. Hint: Consider the interior angles at potential saddle points.

Of course it is possible/allowed to present an algorithm that solves both problems.

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(10+10 \text { points })
$$

## Exercise 5 (Triangulation):

a) Triangulate the polygon shown in Figure 1 using the algorithms from the lecture.
b) Give an algorithm that triangulates a polygon with holes in $O(n \log n)$.


Figure 1: A Polygon.

