# Collaborative transmission in wireless sensor networks

#### Alternative algorithmic approaches

Stephan Sigg

Institute of Distributed and Ubiquitous Systems Technische Universität Braunschweig

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Stephan Sigg

### **Overview and Structure**

- Introduction to context aware computing
- Wireless sensor networks
- Wireless communications
- Basics of probability theory
- Randomised search approaches
- Cooperative transmission schemes
- Distributed adaptive beamforming
  - Feedback based approaches
  - Asymptotic bounds on the synchronisation time
  - Alternative algorithmic approaches
  - Alternative Optimisation environments

### **Overview and Structure**

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  - Feedback based approaches
  - Asymptotic bounds on the synchronisation time
  - Alternative algorithmic approaches
  - Alternative Optimisation environments

# Outline

Alternative beamforming approaches

- Hierarchical clustering
- 2 Local random search
- 3 An asymptotically optimal algorithm
- Environmental changes
  - Velocity of nodes
  - Multiple receiver nodes
  - Increased population size
  - Receive beamforming

Hierarchical clustering

- For feedback based distributed adaptive transmit beamforming:
  - RSS<sub>sum</sub> changes linear with the network size *n*.
  - Bound on the synchronisation time is more than linear in n

Hierarchical clustering



 $E[T_{\mathcal{P}}] = \Theta(n \cdot k \cdot \log(n))$ 

Hierarchical clustering

- Hierarchical clustering
  - Oetermine clusters
  - Synchronise clusters successively (with possibly increased transmit power for nodes)
  - Build and synchronise overlay-cluster of representative nodes from all clusters.
  - Nodes alter carrier phase by phase offset experienced by representative node:

• 
$$\zeta_i = \Re \left( m(t) \mathsf{RSS}_i e^{j2\pi f_c t(\gamma_i + \phi_i + \psi_i)} \right)$$
 (before)

• 
$$\zeta'_i = \Re\left(m(t) \text{RSS}_i e^{j2\pi f_c t(\gamma'_i + \phi_i + \psi_i)}\right)$$
 (after)

Node h from same cluster alters carrier signal

• 
$$\zeta_h = \Re \left( m(t) \text{RSS}_h e^{j2\pi f_c t(\gamma_h + \phi_h + \psi_h)} \right)$$
 to

• 
$$\zeta'_h = \Re \left( m(t) \mathsf{RSS}_h e^{j2\pi f_c t(\gamma_h + \phi_h + \psi_h + \gamma_i - \gamma'_i)} \right)$$

Ideal conditions: All nodes should now in phase

Sinal synchronisation among all nodes

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#### Hierarchical clustering



Hierarchical clustering

Potential problem : Phase noise

- Only one cluster synchronised at a time
- Due to practical properties of oscillators, phases of nodes in the inactive clusters experience phase noise and start drifting out of phase
- Sufficient synchronisation possible in the order of milliseconds

Positive :

• No inter-node communication required

Open Issue :

- More than one hierarchy stage might be optimal
- for optimisation time
- for energy consumption
- Optimum hierarchy depth and cluster size derived by integer programming in time O(n<sup>2</sup>)

Hierarchical clustering

#### Determine optimum cluster size and hierarchy depth :

Expected optimisation time: E[T<sub>Pn</sub>] = c · k · n · log(n)
Expected energy consumption: E[E<sub>Pn</sub>] = c · k · n · log(n) · E<sub>Pn</sub>

Hierarchy and cluster structure that minimises these formulae optimal

Hierarchical clustering

Opt. cluster size and hierarchy depth (integer programming) :

• For a cluster size of *m*:

$$E[T_{\mathcal{P}n}] = E[T_{\mathcal{P}\frac{n}{m}}] \cdot \frac{n}{m} \cdot E[T_{\mathcal{P}m}]$$
$$E[\mathcal{E}_{\mathcal{P}n}] = E[\mathcal{E}_{\mathcal{P}\frac{n}{m}}] \cdot \frac{n}{m} \cdot E[\mathcal{E}_{\mathcal{P}m}].$$

• Define recursion by

$$E_{\text{opt}}[T_{\mathcal{P}n}] = \min_{m} \left[ E_{\text{opt}}[T_{\mathcal{P}\frac{n}{m}}] \cdot \frac{n}{m} \cdot E_{\text{opt}}[T_{\mathcal{P}m}] \right]$$
$$E_{\text{opt}}[\mathcal{E}_{\mathcal{P}n}] = \min_{m} \left[ E_{\text{opt}}[\mathcal{E}_{\mathcal{P}\frac{n}{m}}] \cdot \frac{n}{m} \cdot E_{\text{opt}}[\mathcal{E}_{\mathcal{P}m}] \right]$$

• Start of recursion ( $\eta$  min feasible cluster size):

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Hierarchical clustering

Opt. cluster size and hierarchy depth (integer programming) :

- Time required for calculation is quadratic.
  - With a network of *n* nodes, at most *n*<sup>2</sup> distinct terms
    - E<sub>opt</sub>[T<sub>Pi</sub>]
    - $E_{opt}[\mathcal{E}_{\mathcal{P}_i}]$
- Start calculation at
  - $E_{\text{opt}}[\mathcal{E}_{\mathcal{P}\eta}]$ •  $E_{\text{opt}}[\mathcal{T}_{\mathcal{P}\eta}]$
- All other values by table loop-up in time  $\mathcal{O}(n^2)$  according to
  - $E_{\text{opt}}[T_{\mathcal{P}n}]$
  - $E_{\text{opt}}[\mathcal{E}_{\mathcal{P}n}]$  in time  $\mathcal{O}(n^2)$

#### Hierarchical clustering



- Reduction of synchronisation time and transmission power
- Calculation of optimum cluster size and depth in  $\mathcal{O}(n^2)$

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### Environmental changes

- Velocity of nodes
- Multiple receiver nodes
- Increased population size
- Receive beamforming

#### Local random search



- Global random search:
  - Synchronisation performance might deteriorate when the optimum is near
- With small local search space:
  - Majority of worse points excluded

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An upper bound on the synchronisation performance



#### An upper bound on the synchronisation performance



Analysis in two phases for the synchronisation process

Phase 1: Optimum outside search neighbourhood for at least one node

Phase 2: Optimum within search neighbourhood for all nodes

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An upper bound on the synchronisation performance

### Phase 1: Optimum is outside the neighbourhood • Reach search point with improved fitness: $\geq \frac{1}{2}$



An upper bound on the synchronisation performance

When *i* signals synchronised:

- Improve n i non-optimal signals
- *i* already optimal ones unchanged:

$$(n-i) \cdot \frac{1}{n} \cdot \frac{1}{2} \cdot \left(1-\frac{1}{n}\right)^{i}$$
$$= \frac{n-i}{2n} \cdot \left(1-\frac{1}{n}\right)^{i}$$

since 
$$\left(1-\frac{1}{n}\right)^n < e < \left(1-\frac{1}{n}\right)^{n-1}$$

$$s_i \geq \frac{n-i}{2en}$$

 Expected number of mutations to increase fitness bounded by s<sub>i</sub><sup>-1</sup>.



An upper bound on the synchronisation performance

- Time until optimum is within the neighbourhood?
  - Constant time to leave slice
  - k distinct slices



$$\begin{split} \mathsf{E}[\mathcal{T}_{\mathcal{P}}] &\leq \quad c \cdot \sum_{i=0}^{k} \frac{2en}{n-i} \quad = 2 \operatorname{cen} \cdot \sum_{i=1}^{k+1} i^{-1} \\ &< \quad 2 \operatorname{cen} \cdot \ln(k+1) \quad = \mathcal{O}\left(n \cdot \log(k)\right) \end{split}$$

#### An upper bound on the synchronisation performance



Phase 2: Optimum within search neighbourhood

- Worst case: Increase fitness with probability  $\frac{1}{M}$
- Similar to consideration above:

$$\mathcal{O}(N \cdot n \cdot \log(k))$$

#### Overall synchronisation time :

$$\mathcal{O}(N \cdot n \cdot \log(k)).$$

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#### Local random search



#### A lower bound on the synchronisation time :

- Method of the expected progress
- Similar to estimation for global random search
- Basically: Substitute network size *n* by neighbourhood size *N*

Local random search

### A lower bound on the synchronisation time $% \left( {{{\mathbf{x}}_{i}}} \right)$ :

- Method of the expected progress
- Similar to estimation for global random search
- Basically: Substitute network size *n* by neighbourhood size *N* 
  - Probability to alter individual bit

$$\frac{1}{N \cdot \log(k)}$$
Instead of
$$\frac{1}{n \cdot \log(k)}$$

#### Local random search



A lower bound on the synchronisation time :

• With similar arguments as for global random search, lower bound

$$\Omega(N \cdot \log(k) \cdot \Delta)$$

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### Mathematical simulation environment

#### Impact of the node choice



#### • Fitness measure:

$$RMSE = \sqrt{\sum_{t=0}^{\tau} \frac{\left(\sum_{i=1}^{n} s_i + s_{noise}(i) - s^*\right)^2}{n}}$$

Local random search



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**Received sum signal** 

- Reduce the amount of randomness in the optimisation
- Improve the synchronisation performance
- Improve the synchronisation quality

Search space

- Search space:
  - Spanned by all Configurations of carrier phase offsets  $\gamma_i$
- Search point / Configuration:
  - One possible configuration of carrier phase offsets



**Received sum signal** 

- Fitness function observed by single node
- Constant carrier phase offset for *n* 1 nodes
- Fitness function:

$$\mathcal{F}(\Phi_i) = A\sin(\Phi_i + \phi) + c$$



**Received sum signal** 

Approach:

- Measure feedback at 3 points
- Solve multivariable equations
- Apply optimum phase offset calculated



Received sum signal

#### • Problem:

• Calculation not accurate when two or more nodes alter the phase of their transmit signals



Solution

### An active node will :

- O Transmit with three distinct phase offsets γ<sub>1</sub> ≠ γ<sub>2</sub> ≠ γ<sub>3</sub> and measure feedback.
- From these three feedback values and phase offsets, estimate feedback function and optimum phase offset γ<sup>\*</sup><sub>i</sub>.
- Transmit a fourth time with  $\gamma_4 = \gamma_i^*$ .
- If the deviation is less than 1% save γ<sub>i</sub><sup>\*</sup> as optimal phase offset, otherwise discard it.

A passive node will :

• Transmit 4 times with identical phase offset  $\gamma_i$ .

Solution

- Node estimates the quality of the function estimation itself
- Transmit with optimum phase offset and measure channel again
- When Expected fitness deviates significantly from measured fitness, discard altered phase offset
- Deviation:

 $\begin{array}{l} 1 \mbox{ node: } \approx 0.6\% \\ 2 \mbox{ nodes: } \approx 1.5\% \\ 3 \mbox{ nodes: } > 3\% \end{array}$ 





Synchronisation process





- **()** Transmit with phase offsets  $\gamma_1 \neq \gamma_2 \neq \gamma_3$ ; measure feedback
- Stimate feedback function and calculate  $\gamma_i^*$
- Transmit with  $\gamma_4 = \gamma_i^*$
- Solution Smaller 1% finished, otherwise discard  $\gamma_i^*$

**Received sum signal** 

• Asymptotic synchronisation time:

 $\mathcal{O}(n)$ 

Classic approach:<sup>1</sup>

 $\Theta(n \cdot k \cdot \log(n))$ 

<sup>&</sup>lt;sup>1</sup>Sigg, El Masri and Beigl, A sharp asymptotic bound for feedback based closed-loop distributed adaptive beamforming in wireless sensor networks (submitted to Transactions on Mobile Computing)

Collaborative transmission in wireless sensor networks

#### **Performance estimation**



#### **Performance estimation**



#### **Performance estimation**



• Phase offset of distinct nodes is within  $+/-0.05\pi$  for up to 99% of all nodes.

#### **Performance estimation**



- Asymptotically optimal synchronisation time
- Simulations:  $\approx 12n$
- Further improvement:
  - 3 iterations per turn
  - Utilise last transmission from previous iteration

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Introduction

- Velocity of nodes
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Velocity of nodes

#### Moving receiver :

- Straight line
- Random walk

#### Moving transmitter :

Straight line Random walk

Velocity of nodes

#### Moving receiver :

- Straight line
- Random walk

#### Aspects :

- Only one moving node
- Simple case
- Also applicable when all transmitters move identically

Velocity of nodes

### Moving transmit nodes :

- Straight line
- Random walk

#### Aspects :

Multiple nodes moving Hard case

#### Velocity of nodes



Random walk - receiver :

• Maximum velocity for classic algorithm: 5m/sec

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#### Velocity of nodes



#### Random walk - receiver :

• Max. velocity for Multivariable equations: 5m/sec easily supported

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#### Velocity of nodes



### Random walk - transmitter :

• Maximum velocity for classic algorithm: 2m/sec

#### Velocity of nodes



#### Random walk - transmitter :

• Max. velocity for Multivariable equations: 5m/sec supported

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#### Velocity of nodes



straight line - maximum relative speed :

- Maximum velocity for classic algorithm: 30m/sec
- Regardless if transmitter or receiver move

#### Velocity of nodes



straight line - maximum relative speed :

- Maximum velocity for Multivariable equations algorithm: 60m/sec
- Regardless if transmitter or receiver move

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#### Multiple receiver nodes



#### Multiple receiver nodes

3m					
Knoten	<b>n</b> 0	<i>n</i> 1	n2	m0	m1
Gain zur Anfangsamplitude (Median) [dB]	0,96	2,39	1,40	1,46	1,10
Gain zu einem Knoten (Median) [dB]	2,33	2,32	2,37	3,50	4,05
Anzahl letztes Feedback	5/11	3/11	3/11	8/11	7/11
Amplitude nach Synchronisation [%]	92,4	51,4	65,3	91,0	90,7

#### 12m

Knoten	<b>n</b> 0	<i>n</i> 1	n2	<b>m</b> 0	<i>m</i> 1
Gain zur Anfangsamplitude (Median) [dB]	1,24	0,63	1,39	2,06	1,47
Gain zu einem Knoten (Median) [dB]	2,53	1,09	2,00	2,74	4,18
Anzahl letztes Feedback	2/10	4/10	4/10	5/10	5/10
Amplitude nach Synchronisation [%]	57,1	92,0	86,5	86,4	86,6

### 24m

Knoten	n0	n1	n2	m()	m1
Gain zur Anfangsamplitude (Median) [dB]	1,12	2,33	2,76	3,61	1,67
Gain zu einem Knoten (Median) [dB]	1,2	2,54	2,03	$^{5,15}$	3,76
Anzahl letztes Feedback	4/5	0/5	1/5	4/5	3/5
Amplitude nach Synchronisation [%]	94,2	80,0	61,4	95,8	97,9

Multiple receiver nodes



#### Multiple receiver nodes



Multiple receiver nodes

#### Multiple receiver nodes - issues :

- Only binary feedback value
  - Therefore only classic optimisation approach
- Distance between transmit and receive nodes relative to spatial diversity of nodes in one network
  - Better synchronisation when nodes in one network in spatial proximity
  - When nodes in one network communicate: No issue

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Increased population size

Increased population size - Discussion :

How to achieve population size greater than one?

- Separate transmit times
- WCDMA
- Distinct frequencies simultaneously

Only separate transmit times feasible for WSN

More time for each iteration

- Initial solution: Random search
- Not clear if performance improvement possible by crossover

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Receive beamforming

### Receive beamforming – Discussion :

- Transmit node transmits only once
- Receiver nodes combine received signal fragments in the network
- Tradeoff:
  - Transmission power for in-network communication
  - Transmission over several iterations with receiver node
- More complex computation of transmit nodes

### **Questions?**

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