# Algorithms for context prediction in Ubiquitous Systems

The state prediction approach

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#### **Overview and Structure**

- Introduction to context aware computing
- Basics of probability theory
- Algorithms
  - Simple prediction approaches: ONISI and IPAM
  - Markov prediction approaches
  - The State predictor
  - Alignment prediction
  - Prediction with self organising maps
  - Stochastic prediction approaches: ARMA and Kalman filter
  - Alternative prediction approahces
    - Dempster shafer
    - Evolutionary algorithms
    - Neural networks
    - Simulated annealing

#### **Overview and Structure**

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### Outline

The state predictor method

- Introduction and basic techniques
- State predictors
- 3 Evaluation
- Estimation of Reliability
- 6 Hybrid predictors
- 6 Properties of the state predictor method

Historical remarks

- State predictor: Developed 2003-2005
- University of Augsburg: Jan Petzold<sup>1</sup>
- Origin: Branch prediction in microprocessors

<sup>&</sup>lt;sup>1</sup>Jan Petzold, *Zustandsprädiktoren zur Kontextvorhersage in ubiquitären Systemen*, PhD-thesis, 2005. Stephan Sigg Algorithms for context prediction in Ubiquitous Systems

Scenario

#### Smart Doorplates

- Display current work situation
  - Telephone interview
  - Meeting
  - Working at desk
- Display relevant invormation when person is absent
  - Current location
  - Receive/forward messages
  - Prediction of return time
  - Prediction of next location

Scenario

#### • Smart Doorplates



Scenario

- Smart Doorplates
  - Generation of measurements





Historical remarks

- Branch prediction
  - Loops in program code constitute branches
  - In a pipelined architecture, program execution cycles are interleaved
  - Branch target is only known after ALU.
  - At this time, other (possibly wrong) instructions are already loaded into the pipeline
  - Pipeline flush expensive

#### Historical remarks

#### Example: The DLX-Architecture



Historical remarks

#### Example: The DLX-Architecture



Historical remarks

#### • Example: The DLX-Architecture

IF	ID	EX	MEM	WB					
	IF	stall	stall	IF	ID	EX	MEM	WB	
					IF	ID	EX	MEM	WB
						IF	ID	EX	MEM
							IF	ID	EX
	Server.		i de					IF	ID
		10	100						IF
	IF	IF ID IF	IF ID EX IF stall	IF ID EX MEM IF stall stall	IF ID EX MEM WB IF stall stall IF	IF ID EX MEM WB IF stall stall IF ID IF	IF ID EX MEM WB IF stall stall IF ID EX IF ID IF ID IF	IF     ID     EX     MEM     WB       IF     stall     IF     ID     EX     MEM       IF     stall     IF     ID     EX     MEM       IF     ID     EX     IF     ID     EX       IF     IF     ID     IF     ID     IF	IF         ID         EX         MEM         WB           IF         stall         IF         ID         EX         MEM         WB           IF         stall         IF         ID         EX         MEM         WB           IF         IF         ID         EX         MEM         WB           IF         IF         ID         EX         MEM           IF         IF         ID         EX         MEM           IF         IF         ID         EX         MEM           IF         IF         ID         EX         IF           IF         IF         IF         ID         IF

Speedup due to pipelining

• Calculate the speedup achieved due to pipelining. Assume: Clock Cycle of 10 ns

	Taktzyklen (No pipelining)	Taktzyklen (Pipelining)	Occurence probability
Clock cycle	10ns	10ns	
ALU operations	4	4	40%
Branches	4	4	20%
Memory operations	5	5	40%
Pipelining overhead	0ns	1ns	

Speedup due to pipelining

Average instruction execution time = Clock cycle · Average CPI

- $= \text{Clock cycle} \cdot \text{Average CPI} \quad (1)$
- $= 10ns \cdot ((40\% + 20\%) \cdot 4 + 40\% \cdot 5)$

$$=$$
 10*ns* · 4.4

Pipelined implementation: Clock must be largest time for any stage in the pipeline (10ns) plus overhead: 10 + 1 = 11ns. This is the average instruction execution time. Speedup from pipelining:

Speedup from pipelining =  $\frac{\text{Average instr. time unpipelined}}{\text{Average instr. time pipelined}}$  (2) =  $\frac{44ns}{11ns}$ = 4 times

Historical remarks

- Branch prediction important to reduce program execution time
- Pipeline flush/stall expensive
- Important therefore: Accuracte prediction

Historical remarks

#### • Simple but effective branch prediction schemes:

- Always branch
- Sustain last decision

Branch prediction - dynamic branch prediction

- Dynamic branch prediction schemes
  - Branch decision based on recent behaviour
  - Behaviour stored in branching-tables
  - Tables updated in case of erroneous predictions
- Various implementations
  - Two-bit predictor
  - Two stage adaptive predictor
  - Hybrid predictor

Branch prediction - One- and Two bit predictors



#### One-bit predictor

- Most simple dynamic branch prediction technique
- Repeat last branching command
- Every branch inside an iterated loop is correctly predicted
- In nested loops: First and last branch incorrectly predicted
  - Two bit predictors more accurate in this case

Branch prediction - One- and Two bit predictors



Two-bit predictor

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Branch prediction - One- and Two bit predictors

- Two-bit predictor
  - Two wrong predictions in turn required to change prediction behaviour
  - In nested loops, only one wrong prediction: After the last iteration
- Two implementations of Two-bit predictors:
  - Saturation counter
  - Hysteresis counter
- Extension to *n*-bit predictors possible but nearly no improvement in predictoin accuracy

Branch prediction – Correlation predictors

- Two bit predictors only consider the branch itself
- Intercorrelation between distinct branches are not considered
- However, intercorrelations do matter <sup>2</sup>

<sup>&</sup>lt;sup>2</sup>Shien-Tai Pan, Kimming so and Joseph T. Rahmeh, *Improving the acuracy of dynamic branch prediction using branch correlation*, In: Proceedings of the fifth international conference on Architectural support for programming languages and operating systems (ASPLOS V), pp 76-84, 1992.

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Branch prediction – Two-stage adaptive predictors

- Two dimensional prediction tables<sup>3</sup>
  - First table selects prediction bits of second table
  - First table: Branch history (Shift register)
  - Second table: Pattern history

<sup>&</sup>lt;sup>3</sup>Tse-Yu Yeh and Yale N. Patt, *Alternative implementation of two-level adaptive branch prediction*, In: Proceedings of the 19th annual symposium on computer architecture (ISCA-19), pp 124-134, 1992.

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Branch prediction – Two-stage adaptive predictors

- Classes of two-stage adaptive predictors<sup>4</sup>
  - First letter: [G,P,S] Mechanism in first table
  - Last letter: [g,p,s] Mechanism in second table
- BHR = Branch Histroy Register of k Bits length
- PHT = Pattern History Table

	global PHT	per-set PHTs	per-address PHTs
global BHR	GAg	GAs	GAp
per-address BHT	PAg	PAs	PAp

<sup>&</sup>lt;sup>4</sup>Tse-Yu Yeh and Yale N. Patt, A Comparison of Dynamic Branch predictors that use two levels of branch hostory, In: Proceedings of the 20th annual symposium on computer architecture (ISCA-20), pp 257-266, 1993.

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Branch prediction - Two-stage adaptive predictors

Example: GAg(k) predictor



Branch prediction - Two-stage adaptive predictors

• Example: GAg(4) predictor



Pattern History Table (PHT)

Branch prediction – Two-stage adaptive predictors

- Problems of two-stage adaptive predictors
  - The same bit pattern in BHR can be referenced to different parts of a program
  - This leads to possible PHT-interferences
    - Unrelated branching commands impact the prediction of each other
    - However, solutions to this problem exist

Prediction by partial matching

- PPM-Algorithm of order n is composed of n + 1 Markov Predictors of oder 0 to n
  - Algorithm tries to find a pattern that matches last *n* states with Markov predictor of order *n*
  - If this is not successful, Markov predictor with decreased order is instantiated.

Prediction by partial matching

Prediction by partial matching				
Input: $z_1, \ldots, z_{t-1}$	(Sequence of States)			
Output: z <sub>t</sub>	(next state)			
1 REPEAT 2 Choose Markov predi	ctor of order n			
3 If (pattern $z_{t-n}, \ldots, z_{t-1}$ is f	ound)			
4 Calculate prediction				
5 BREAK				
6 Else				
7 $n=n-1$				
8 UNTIL $n = 0$				
9 If $(n = 0)$				
10 Calculate prediction with p	redictor of order O			

#### Prediction by partial matching



### Outline

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- 4 Estimation of Reliability
- 6 Hybrid predictors
- 6 Properties of the state predictor method

Introduction

- The state predictor method
  - Utilise branch prediction techniques implemented on microprocessors
  - Distinction between local and global prediction
    - Example sequence: ACBCACBCABCAB
    - Global sub-sequence:
       BCAB
    - Local sequence with respect to A: CCBB
  - Local sequence: Context observes only neighbouring contexts that succeed its own occurrence

Introduction

• Various state classes of state prediction methods



One-stage state predictors



- One-state predictor
  - For every context a prediction graph of succeeding contexts is maintained
  - Prediction is always adapted to the observed succeeding context.
  - When prediction is incorrect, predicted context is adapted to the observed context

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One-stage state predictors

#### • Example: Location prediction



One-stage state predictors

- Benefits
  - Small memory requirements
  - Fast 'adaptation' (lerning)
- Drawbacks
  - Learned behaviour is rapidly lost/forgotten

One-stage state predictors

$$A\left(\left(A1\right) \xrightarrow{B} \left(A0\right) \xrightarrow{B} \left(B0\right) \xrightarrow{B} \left(B1\right)\right) B$$

- Two-state predictor
  - For every context a prediction graph of succeeding contexts is maintained
  - Two possible states for every possible succeding context
    - Weak state
    - Secure state
One-stage state predictors



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One-stage state predictors

- Benefits
  - Small memory requirements
  - Fast 'adaptation' (lerning)
- Drawbacks
  - Learned behaviour is easily lost/forgotten

One-stage state predictors

- k-state predictor
  - For every context a prediction graph of succeeding contexts is maintained
  - k possible states are associated with every possible succeeding context
    - k − 1 weak states
    - One secure state

One-stage state predictors



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One-stage state predictors

- k-state predictor Aspects
  - Benefits
    - Small memory requirements
  - Drawbacks
    - Learning of behaviour only restricted to local context view
  - Further aspects
    - Dimension k can also be learned

Two-stage state predictors

- Global two-state predictors
  - Prediction on observed global sequence
  - Prediction stated by Two-state predictor
  - Also: Prediction by arbitrary k-State predictor possible

Two-stage state predictors

#### Example

• Observed context pattern:

#### ABCACBABACBACB

Schieberegister		Musterverlaufstabelle		
$A C B \cdots$		Muster	Two-State-Prädiktor	
		A B A	C0	
		A B C	AO	
		A C A	-	
	$x \leftrightarrow x \leftrightarrow$	A C B	A1	
		B A B	A0	
		B A C	B1	
		B C A	C0	
		B C B	-	
		C A B	-	
		C A C	<i>B</i> 0	
		C B A	C0	
		C B C	-	

Two-stage state predictors

#### • Local two-state predictors

- Prediction of local context sequences
- Otherwise identical to global two-state predictors

#### State predictor methods

Prediction by partial matching

- Extension of two-stage state predictor method by PPM
- Instead of fixed order of the prediction method: Variable order dependent on the maximum matching pattern length found in the obseved context sequence.

#### State predictor methods

Prediction by partial matching – PPM and simple PPM



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Augsburg benchmarks



#### Augsburg benchmarks











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- Prediction is naturally error prone
- Sometimes, no prediction might be better than an erroneous prediction

- Static reliability
  - Assumption:
    - For some patterns or contexts a prediction is not taken
    - Patterns are not classified as typical
    - Patterns frequently change Low prediction accuracy
  - Contexts divided into reliable and non-reliable contexts

- Consideration of secure states
  - Only applicable to two-state predictors
  - Two state predictor has for every context two states
    - Weak state
    - Secure state
- When context is observed often, it becomes secure
- Predictions are provided exclusively in secure states

Methods for reliability estimation

• Consideration of secure states



Methods for reliability estimation

- Threshold method
  - Compare recent prediction accuracy with predefined threshold.
    - Secure state: Accuracy above threshold
    - Insecure state: Accuracy below threshold
  - Correct predictions: c
  - Incorrect prediction: *i*

$$\frac{c}{c+i} \ge \alpha : \text{Secure state}$$
(3)

Schieberegister

$$p_1 \dots p_r$$
 ....

Musterverlaufstabelle

	Muster	Two-State-P.	С	i
·	$p_1 \dots p_r$	C1	x	y



- Reliability counter
  - Accuracy of predictions 'counted' by reliability counter
  - Initial position of counter arbitrary
  - When prediction is correct, counter is increased
  - When prediction is incorrect, counter is decreased
  - When counter exeeds threshold, prediction is provided



Evaluation

• Static reliability – Augsburg Benchmarks



Evaluation

• Secure states – Augsburg Benchmarks



Person A				
Ordnung	Steigerung			
1	48,80%			
2	47,66%			
3	39,41%			
4	$50,\!65\%$			
5	29,57%			
PPM(5)	56,35%			
SPPM(5)	$54,\!39\%$			

Evaluation

• Secure state - Nokia context data



Zell-ID				
Ordnung	Steigerung			
1	35,79%			
2	41,29%			
3	41,82%			
4	39,23%			
5	38,04%			
PPM(5)	47,96%			
SPPM(5)	40,85%			

Evaluation

• Threshold method – Augsburg Benchmarks



Evaluation

• Threshold method - Nokia context data



Evaluation

• Reliability counter - Augsburg Benchmarks



Evaluation

• Reliability counter - Nokia context data



Conclusion

• Is the state prediction method a Markov prediction class algorithm?

Conclusion

 Is the state prediction method a Markov prediction class algorithm?



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#### Conclusion

- Is the state prediction method a Markov prediction class algorithm?
  - The state prediction approach defines the mechanism to adapt transition probabilities
  - Only transition probabilities 1 and 0 possible
  - Markov prediction more powerful



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### Hybrid predictors

Introduction

- Can the combination of multiple prediction approaches improve the prediction accuracy?
  - Warm-up predictor
  - Majority predictor
  - Reliability predictor
Warm-up-predictor

- Prediction approaches with low-order often provide quick szenario adaptation
- Complex patterns not possible with low-order models



Majority prediction

- Compute prediction by various prediction approaches
- Majority of prediction results determines actual prediction
  - Relative majority
    - Prediction that was stated most often by all approaches
  - Bare majority
    - More than half of the stated predictions are identical
  - Absolute majority
    - More than half of the possible predictions identical

Majority prediction

#### • Relative majority



Majority prediction

#### • Bare majority



Majority prediction

• Absolute majority



Reliability predictor

• Choose the prediction with highest prediction accuracy



Reliability predictor

#### • Choose the prediction with highest prediction accuracy

- Several selection criteria possible
  - Primary selection criterium
  - Secondary selection criterium
  - Tertiary selection criterium

Reliability predictor

#### Selection criteria

primär	sekundär	tertiär
Sicherer Zustand	relativ	ohne Barriere
		mit Barriere
	einfach	ohne Barriere
		mit Barriere
Schwellenwert-Verfahren	relativ	ohne Barriere
		mit Barriere
	einfach	ohne Barriere
		mit Barriere
Zuverlässigkeitszähler	relativ	ohne Barriere
		mit Barriere
	einfach	ohne Barriere
		mit Barriere

Conclusion

- It was shown, that the warm-up predictors achieve low accuracy
- With Majority predictors and reliability predictors the prediction accuracy can be improved

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Processing load

Runtime for computing a prediction: (O(1))
 Current state directly prediction

Memory requirements

- Memory requirements
  - Dependent on the number of contexts observed size of the transition matrix
  - Order 1:  $O(|C|^2)$
  - Order k:  $O(|C|^{k+1})$

Prediction horizon

- Prediction horizon can be extended by iterative prediction
  - Utilise predicted contexts as input
- Problem: Less accurate
  - Predicted contexts more error prone than measured values

#### Properties of the state predictor approach Adaptability

- The state prediction approach is able to adapt to changing environments
  - Adaptation only to simple patterns

Multi-dimensional time series

- The state prediction algorithm is not suited for multi-dimensional time series
  - Designed for one-dimensional Input
  - Possible: Aggregation of multi-dimensional time series to one-dimensional time series.

Iterative prediction

#### • Iterative Prediction possible

- Steep decrease in prediction accuracy expected since prediction horizon is only 1
- Increase of prediction horizon possible by Aggregation of context sequence of fixed length in one state
  - Prediction horizon fixed
  - Increase in Memory consumption and processing time
  - When I contexts are aggregated:  $I^C$  states
  - Runtime:

$$O(n \cdot I^{C^2}).$$

Memory consumption:

```
O(I^{C^2}) (order one)
O(I^{C^{k+1}}) (order k)
```

Prediction of context durations

- Prediction of context duration not possible
  - Only simple sequence of occurring contexts possible

Approximate matching of patterns

#### • Exact pattern matching

• The state prediction algorithm utilises exact pattern matching

Context data types

- All context data types supported
  - Every distinct context type one state
  - Probably drastic increase in runtime and memory consumption for numeric context types
  - Possible: Assign intervals to states

Pre-processing

- Pre-processing required to construct context transition probabilities
- On-line approach feasible learning
- Runtime: O(k)
  - Count frequency of specific context transitions in training time series of length  $\boldsymbol{k}$

## Aspects of prediction algorithms

Summary

	IPAM	ONISI	Markov	CRF
Numeric Contexts	yes	no yes		
Non-numeric Contexts	yes	yes	yes	
Complexity	O(k)	()	$O(C^2)$	
Learning ability	(no)	yes	yes	
Approximate matching	no	no	no	
Multi-dim. TS	(no)	(no)	(no)	
Discrete data	yes	yes	yes	
Variable length patterns	no	yes	no	
Multi-type TS	yes	no	(no)	
Continuous data	no	no	no	
Pre-processing	O(k)	_	O(k)	
Context durations	no	no	no	
Continuous time	no	no	yes	

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## Aspects of prediction algorithms

Summary

	SPM	Align	SOM	PCA
Numeric Contexts	yes			
Non-numeric Contexts	yes			
Complexity	O(1)			
Learning ability	(yes)			
Approximate matching	no			
Multi-dim. TS	(no)			
Discrete data	yes			
Variable length patterns	yes			
Multi-type TS	no			
Continuous data	no			
Pre-processing	O(k)			
Context durations	no			
Continuous time	no			

Conclusion

- Simple, straightforward prediction approach
- Model can be applied to numerical and non-numerical data alike.
- Special case of a Markov predictor
- Less powerful than Markov prediction
- Not suited for complex prediction scenarios
- Prediction that reaches farther into future implicitly utilises already predicted data which might consequently decrease the prediction accuracy.