

# Taming the *flexibility versus safety* challenge in Distributed Embedded Systems

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## Definitions

### ❖ Safety

❖ Property of a computer system for which reliance can justifiably be placed in the **absence of catastrophic accidents** (mishaps).

(Laprie, 1994)

❖ Well defined, e.g., IEC 61508 and its 4 SIL levels

### ❖ Flexibility ??

❖ “a **ready capability to adapt** to new, different, or changing requirements”

(Merriam-Webster on-line)

❖ Within embedded systems, **flexibility** can apply to both **design-time** and **run-time** features

## Motivation

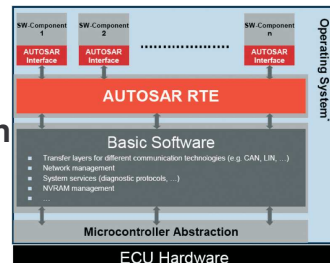
### • Typical design of safety-critical systems

- ❖ Adoption of **static (design-time) techniques**
  - Maximize *a priori* knowledge
  - Single well-defined operational mode
  - Simplify error detection and fault tolerance
  - Facilitate replica determinism (for modular redundancy)
- ❖ Downside:
  - **Modular redundancy** and hooks for **future changes** imply **over-provisioning** of resources (**exclusive use**)
    - Low efficiency in resource usage → High costs
  - On-line resource reallocation typically impossible
- ❖ **Could we apply *resource sharing* to increase flexibility and efficiency thus decreasing costs?**

## Motivation

### • What can we do with *resource sharing*?

- ❖ **Share ECUs / networks across subsystems**
  - Reduce active components
  - Requires SW components that are unaware of HW
  - Increase design flexibility (more options for OEMs ...)
  - Purpose of **Autosar, IMA, IEC61499, ...**
- ❖ **Multiple operational modes**
  - Modes change on-line using overlapping resources
- ❖ **Promote functional integration**
  - Create new functionality by integrating subsystems
  - Centralized control



# Motivation

## • And also

### ❖ Post-deployment reconfiguration

- Adding/changing subsystems at anytime
  - e.g., adding SW components using resources that are free

### ❖ Graceful degradation

- Replicas with different QoS using remaining resources
  - Unintentional redundancy
  - Low-cost fault tolerance

### ❖ Dynamic QoS management

- Allocating resources bandwidth dynamically as needed
- Applications must be able to execute with different QoS levels provided by the infrastructure
  - Different **modes per subsystem** (including *standby*)

All run-time

# Problem

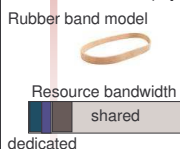
## • Can we reconcile flexibility and safety?

### ❖ Unforeseen and unsafe changes can occur because of

- Insufficient resources
- Coupling between functionality and timing
- Mutual interference across subsystems

### ❖ Hypothesis: **Yes, if we...**

- **Bound flexibility**, i.e., keeping the system within a well defined operational region → strict policing
- Guarantee the **minimum resources needed** for critical tasks
  - **And manage (share) the remaining resources**
- Use **components (networks...)** providing needed properties
  - Partitioning (virtualization), policing, QoS management, error confinement, atomic broadcast,...



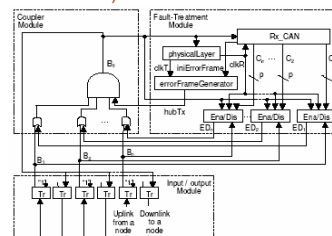
## Our contributions in this topic

- **Boosting error confinement in CAN**
  - ❖ **CANcentrate and ReCANcentrate**
    - The first CAN hubs designed to improve dependability
- **Towards flexible partitions**
  - ❖ **FTT-CAN, FTT-SE** (software solutions)
    - Fully configurable ET/TT partition
    - Support for further partitioning: **Server-CAN, Server-SE**
      - Confine errors in the temporal domain
    - Support for flexible modes
      - Modes defined at subsystem level, independently
  - ❖ **FTT-switches** (hardware solutions)
    - HaRTES (Hard Real-Time Ethernet Switch)
    - Partitions enforced by the switch (a boost on robustness)

## CANcentrate

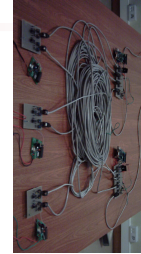
- ❖ **First CAN-hub designed for error-confinement**
  - **Wired-AND of CAN bus replaced by logical AND**
  - **Uplinks separated from downlinks**
  - **Allows fast detection of several types of errors**
    - **Link isolation when error threshold crossed**  
(latency to isolate stuck-at or bit-flipping faults: 73µs, 150...600µs)
    - **Automatic reintegration after error-free period**  
(latency to reintegrate isolated links: 5.2ms)

- ❖ **Works with COTS CAN controllers and any existing application**
  - It is just a wiring replacement

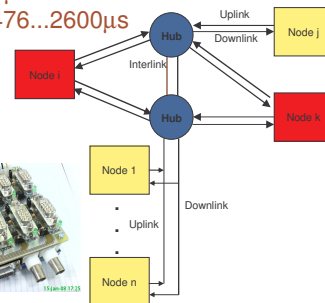
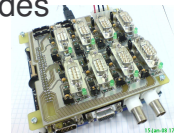


## ReCANcentrate

- ❖ **First replicated CAN-hub architecture**
  - Targets very demanding safety requirements
  - Replicated hubs are synchronized bit-by-bit
  - Made by two interconnected CANcentrate hubs
    - *Hubs can isolate / reintegrate one another*
    - Latency to isolate stuck-at errors: 216µs
    - Latency to isolate bit-flipping errors: 476...2600µs
    - Reintegration latency: 4.2ms



- ❖ Supports **mixed architectures** with critical / non-critical nodes as well as bus segments

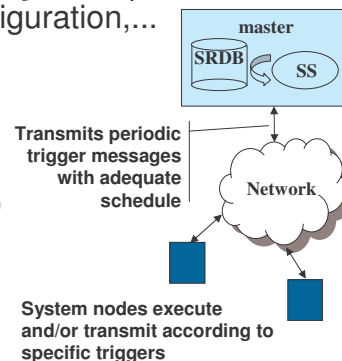


## FTT-xx (software)

The **FTT** protocols provide a **synchronous framework** for distributed applications that, nevertheless, is amenable to **operational flexibility** as required for on-line adaptation, on-line reconfiguration,... but **under adequate control**

Examples:

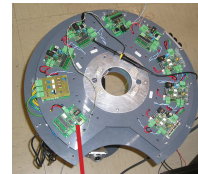
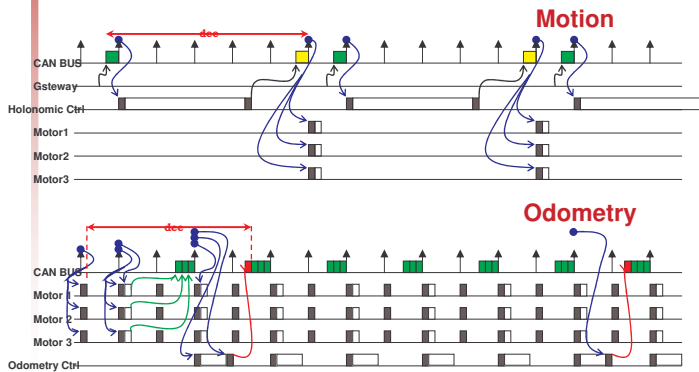
- ✓ Supporting **adaptation in control applications** (reacting to overloads)
- ✓ Creating **virtual channels** with guaranteed QoS, even with varying load (partitioning)



Flexible Time-Triggered architecture

# FTT-CAN – timing control

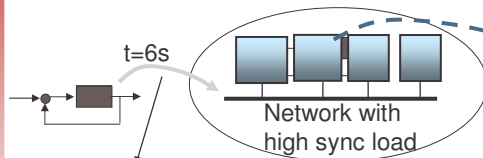
- Flexibility to
  - ❖ Separate functionality from timing
  - ❖ Update timing easily (phases)



3rd of July, V. Silva, R. Marau, L. Almeida, J. Ferreira, M. Calha, P. Pedreiras, J. Fonseca. Implementing a distributed sensing and actuation system: The CAMBADA robots case study. ETFA 2005, Catania Italy, September 2005.

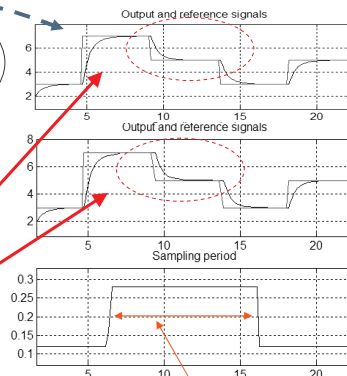
# FTT-CAN – rate adaptation

- Flexibility to
  - ❖ Adapt rates of control loops dynamically (e.g., for overloads)



- Rejected with fixed h
- Accepted with variable h

h	ISE	Degrad	BW
0.12s	48.2	0	1
0.12s / 0.28s	50,3	4,4%	0,66



Period adaptation interval from t=6s to t=16s

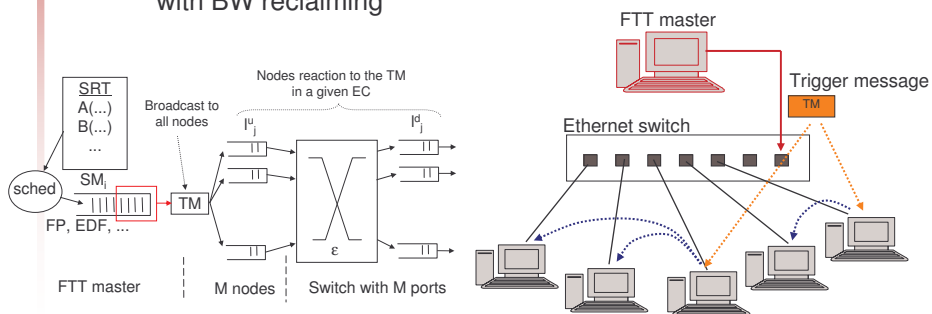
3rd of July, 2009 A. Antunes, P. Pedreiras, A. Mota, L. Almeida, ETFA 2005, ANIPLA 2006, INDIN 2007

## FTT-SE – load control

### ✓ Flexibility to

- ✓ Support **any scheduling policy**
- ✓ Support **dynamic QoS management**
- ✓ Create **partitions** (virtual channels)
- ✓ **Seamlessly integrate sync/async/NRT traffic** with BW reclaiming

Traffic is **scheduled per cycles** considering links capacity.  
**Memory overloads are eliminated**



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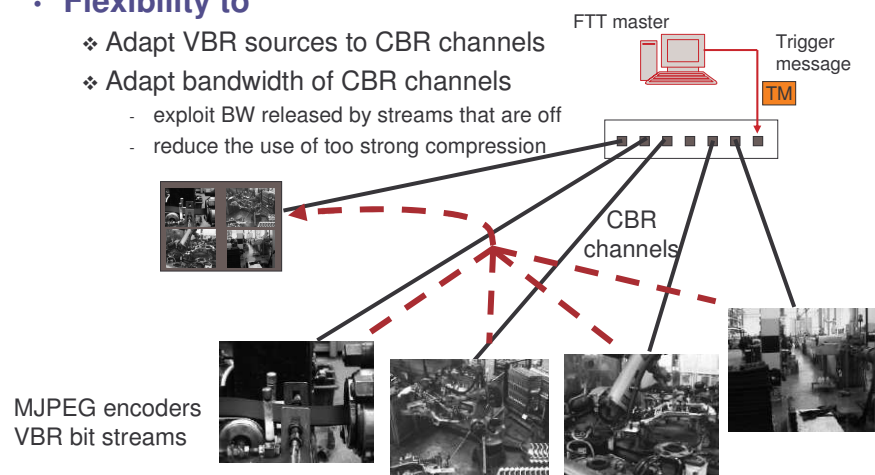
R. Marau, L. Almeida, P. Pedreiras. Enhancing Real-Time Communication over COTS Ethernet switches. WFCS 2006, IEEE 6th Workshop on Factory Communication Systems, Turin, Italy, June 2006.

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## FTT-SE – QoS adaptation

### • Flexibility to

- ❖ Adapt VBR sources to CBR channels
- ❖ Adapt bandwidth of CBR channels
  - exploit BW released by streams that are off
  - reduce the use of too strong compression



3rd of July

J. Silvestre, L. Almeida, R. Marau, P. Pedreiras. Dynamic QoS Management for Multimedia Real-Time Transmission in Industrial Environments. ETFA 2007, Patras, Greece, September 2007.

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## FTT-SE – QoS adaptation



5 Mbit/s  
V  
1 Mbit/s

after 20s

1 Mbit/s  
V  
5 Mbit/s

**6Mbit/s at all times**

## FTT(Server)-SE – partitions

### • Flexibility to

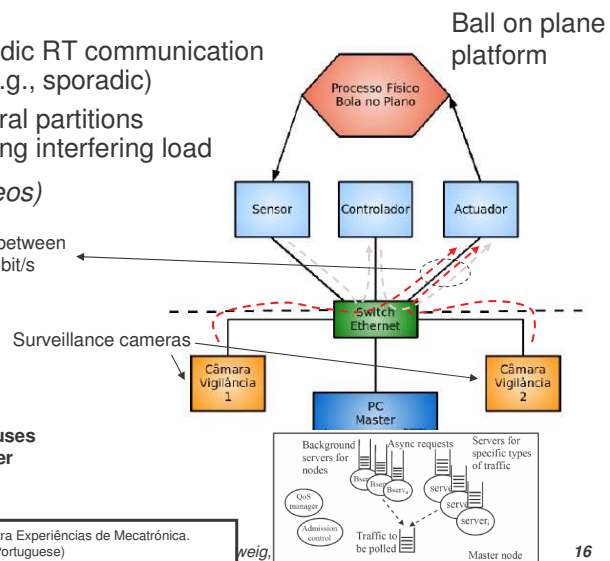
- ❖ Support aperiodic RT communication with servers (e.g., sporadic)
- ❖ Enforce temporal partitions resilient to strong interfering load

(see videos)

Total load varied between 18Mbit/s and 91Mbit/s

1 short control packet together with 2 long video frames

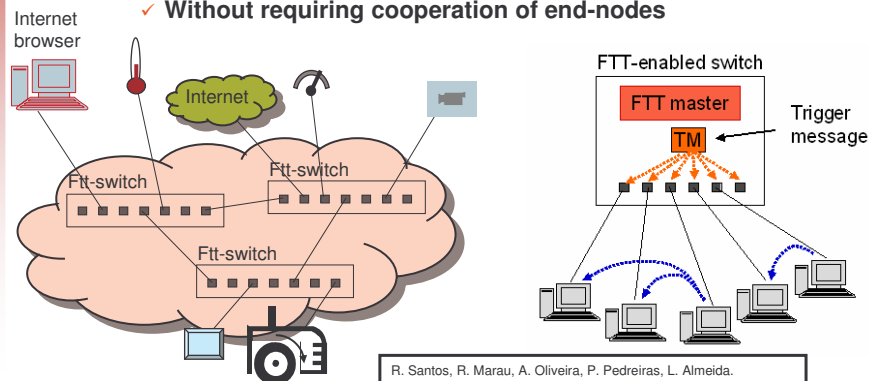
Without isolation, too high load causes control packet losses and large jitter



## FTT-xx (hardware)

Integrating FTT features in the switches results in **timeliness, flexibility and high robustness**

- ✓ Enforce negotiated channel characteristics (**Policing**)
- ✓ Confine specific traffic to separate windows (**Partitioning**)
- ✓ **Without requiring cooperation of end-nodes**



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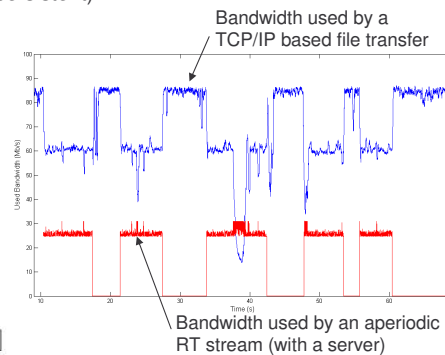
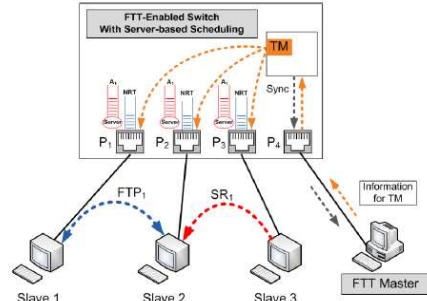
R. Santos, R. Marau, A. Oliveira, P. Pedreiras, L. Almeida.  
Designing a Customized Ethernet Switch for Safe Hard Real-Time Communication. WFCSS 2008. Dresden, Germany, 21-23 May 2008.

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## HaRTES – server-based scheduling

### • Flexibility to

- ❖ Enforce temporal partitions with efficient servers (e.g., sporadic)
  - That support BW reclaiming by NRT traffic (background).
- ✓ 1 **real-time aperiodic** stream (switched on and off, handled with a server)
- ✓ 1 **non-real-time stream** (bursty and persistent)
  - ✓ **Large file transfer**
- ✓ Sent to the same output link



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## Conclusions

- **Resource sharing** is a key to build more **resource efficient** distributed embedded systems
- Further benefits may arise from **Dynamic reconfiguration** and **QoS management** to always adjust resource usage to current needs
- Using these techniques in **safety-critical** systems requires
  - ❖ Strong protection mechanisms across resource partitions
  - ❖ Prompt temporal control for fast adaptations of resource allocations
- We have shown **several networking components** that provide such features and can potentially open the way to **flexible safety-critical** applications (thus less expensive)
  - ❖ **(Re)CANcentrate, FTT(Server)-xx**